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13. ABSTRACT

Pseudo random number generating programs coded in machine language for the IBM 360/65 computer have been prepared for producing random values from the normal, exponential, and gamma distributions as well as from any discrete probability distribution. The schemes are based largely on combinations of sophisticated techniques first suggested by Marsaglia. The distributions simulated are exact within the word size of the computer and average production time per number is very short, e.g. 10 to 40 microseconds depending upon the particular distribution.

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PSEUDO RANDOM NUMBER GENERATORS FOR
STATISTICAL APPLICATIONS

LOVICK EDWARD CANNON, III

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CHAPTER 1

INTRODUCTION

With the increasing use of electronic computers, random sampling methods have become a very useful tool for providing solutions to problems involving probability as well as for problems of a deterministic nature. The availability of sequences of numbers which appear to be drawn at random from particular probability distributions is a vital ingredient in the random sampling process. Such numbers are referred to as pseudo-random numbers or, for convenience, simply as random numbers. This paper is concerned with the rapid and accurate generation of such numbers in a stored computer program.

The use of random sampling to estimate distribution functions originated with Student [17] in 1908. At that time and for some time to follow, necessary random numbers were obtained by drawing cards from a deck, counters from an urn, or by rolling dice. Such processes are very slow and make it quite difficult to insure randomness.

To facilitate the use of random numbers large tables of random digits were compiled. The first such table, published by Tippett [19] in 1927, consisted of 41,600 digits taken at random from census reports and combined into 10,400 four digit numbers. The requirement for a larger set of numbers led to the publication of 100,000 random digits in 1939 by Kendall and Babington-Smith [7]. These digits were generated by means of

a mechanical process and were the first to be so produced. Kendall and Babington-Smith are also responsible for developing many of the tests frequently applied to sequences of random numbers (5,6). Other such tables have since been constructed, primarily by the use of physical devices, culminating in the most extensive, "A Million Random Digits with 100,000 Normal Deviates" published by the RAND Corporation [15] in 1955.

The use of a physical device for the generation of random numbers on line with a computer is both expensive and difficult to maintain. The storing of a table of random numbers on magnetic tape or on cards is also an unsatisfactory method of generating numbers for use in a computer. This would necessitate the use of an input device; also the time required to read in the numbers would be excessive.

The development of arithmetic procedures for the generation of random numbers began in the 1940's with the introduction of computers. The first such procedure was suggested by von Neumann and Metropolis in 1946 and is described in [4]. It was also during this time that calculations involving random numbers received the picturesque name "Monte Carlo." The "middle-square" method proposed by von Neumann and Metropolis is simple, fast, and requires only an initial starting value. In this procedure each new number is produced by taking the middle n digits of the square of the previous n -digit number. As pointed out in [4] and in [16] the sequence of numbers produced using this method sooner or later degenerates to a cycle which often is very small, and at worst consists of only a single number. In addition some of the statistical tests performed on samples generated by the middle-square procedure have resulted in failures.

Improvements on the von Neumann method are plagued by similar difficulties

[2] and [20].

Of the various arithmetic procedures sequences of random numbers with the best statistical properties and longest periods are generated by means of the congruence relation

$$x_{n+1} = a \cdot x_n + c \cdot (\text{mod } m). \quad (1.1)$$

The representation (1.1) is termed the mixed congruential method. The multiplicative congruential method is defined by taking $c = 0$ in (1.1). Lehmer [8] is credited with the invention of the multiplicative method. In their very informative survey Hull and Dobell [3] prescribe conditions for a , m , c , and x_0 which insure maximum period. Other procedures for the generation of random numbers based upon reduced Fibonacci series [21] and upon transcendental numbers [3] are inferior to those based on the congruence relation (1.1) (See [4]). In [9] MacLaren and Marsaglia present a table look-up scheme which seems to offer some promise.

Random numbers generated by arithmetic procedures are not "truly random" in that the entire sequence is determined in advance and can be reproduced simply by using the same starting value, x_0 . The statistical behavior of sequences generated by both the mixed and multiplicative congruential methods is quite good, and, in fact, numbers generated in this fashion do appear to be drawn at random from the uniform (0,1) distribution

$$f(u) = 1, \quad 0 < u < 1. \quad (1.2)$$

The reproducibility of sequences is an advantage in debugging and in certain calculations it may be desirable to reproduce a given sequence.

Bargmann [1] has provided a procedure for the generation of independent uniform (0,1) random numbers on a binary computer. This procedure, which is a form of the multiplicative congruential method, is defined by

$$u_{n+1} = a \cdot u_n \pmod{2^{32}} \quad (1.3)$$

where a is chosen such that $a + 1$ and $a - 1$ end in as few zero bits as possible. This is insured by requiring that neither $a - 1$ nor $a + 1$ be divisible by 2^k for $k = 3, 4, \dots$. Choosing a as an odd power of 5 determines that the low order 3 digits will be 125, and neither 124 nor 126 is divisible by 2^k for $k = 3, 4, \dots$. Hence, $a = 5^{13}$ is a reasonable choice. This procedure requires that u_0 , the starting value, be an odd positive integer.

The use of this procedure results in a very fast computer program requiring only 3 operations: load, multiply, and store. A description of the usage of this program along with time and storage requirements is given in Chapter 6 of this paper. Results of statistical tests performed on this procedure are presented in Chapter 5, and a listing of the program is given in Appendix A.

The preceding discussion has dealt only with the generation of random numbers which appear to be uniformly distributed. In principle it should be very easy to obtain any other distribution from the uniform distribution, requiring only a solution to the equation

$$u = F(y) \quad (1.4)$$

for y , where u is uniformly distributed on the interval (0,1) and F is the

required cumulative distribution function. When the inverse of F is known, as for the exponential distribution, this is a simple matter:

$$y = -\ln(u). \quad (1.5)$$

However, the evaluation of the \ln function may be somewhat time consuming. An alternative approach might be to store a value of y for each possible value of u based on the relation

$$y = F^{-1}(u). \quad (1.6)$$

To generate y , simply generate a uniform $(0,1)$ random number u which will determine the location of a stored value of y . Let u be given to 8 digits and let β be the base of the number system in which the digits are represented, then for this procedure a total of β^8 storage locations are required.

Various approximations have been proposed for the normal distribution. Most of these involve taking the sum of a fixed number of uniform $(0,1)$ deviates or the use of Chebyshev approximations [14] or a combination of the two [18]. Such approximations frequently lack accuracy and are either slow or space consuming.

The procedures suggested by Marsaglia et al. in [11], [12], and [13] for transforming uniform $(0,1)$ random numbers to random numbers having other distributions are superior to any encountered. They are fast, require minimal space, and are simple to program. In addition these procedures are completely accurate; the precision of the result is dependent only on the word size of the computer. These procedures along with programs and results of statistical tests are presented in detail in this paper.

CHAPTER 2
GENERATION OF DISCRETE RANDOM VARIABLES
IN A COMPUTER

Let Y be a discrete random variable with point probabilities $p_i = \Pr(Y = y_i)$ for $i = 1, 2, \dots$. The direct way to generate Y in a computer is to generate a uniform $(0, 1)$ random number u and put $Y = y_i$ if $p_1 + p_2 + \dots + p_{i-1} < u \leq p_1 + p_2 + \dots + p_i$. However, techniques based on this method lead to complicated programs that are excessively time consuming.

An alternative method proposed by Marsaglia [11] is simple to program and requires minimal time and storage. Let p_i for $i = 1, 2, \dots, n$ be expressed by k digits as $p_i = .\delta_{i1}\delta_{i2} \dots \delta_{ik}$. If the domain of the random variable is infinite, the probability distribution must be truncated at some p_n . The fastest method for generating Y is as follows.

Let β be the base of the number system in which the δ_{ji} 's are represented. In memory locations 0 to $(\beta^k - 1)$ store $\delta_{11}\delta_{21}\delta_{31} \dots \delta_{k1} y_1$'s, $\delta_{12}\delta_{22} \dots \delta_{k2} y_2$'s, \dots , $\delta_{1n}\delta_{2n} \dots \delta_{kn} y_n$'s. If u is a uniform $(0, 1)$ random number, $u = .d_1d_2 \dots d_k$, look up the number in location $d_1d_2 \dots d_k$ and let that be Y .

Though this may be the fastest method, it clearly requires an excessive amount of storage space. Even if the p_i 's are truncated to four digits, β^4 memory locations will be required.

Marsaglia [11] suggests a technique that offers a considerable reduction in storage space with very little sacrifice in execution time. For convenience assume that the p_i 's are truncated to four digits. We have the following situation:

Value of Y	Probability
y_1	$\cdot \delta_{11} \delta_{21} \delta_{31} \delta_{41}$
y_2	$\cdot \delta_{12} \delta_{22} \delta_{32} \delta_{42}$
y_n	$\cdot \delta_{1n} \delta_{2n} \delta_{3n} \delta_{4n}$

Define

$$P_0 = 0, \quad P_r = \beta^{-r} \sum_{i=1}^n \delta_{ri} \quad (2.0.1)$$

for $r = 1, 2, 3, 4$, and

$$\Pi_s = \sum_{j=1}^s \sum_{i=1}^n \delta_{ji} \quad (2.0.2)$$

for $s = 1, 2, 3, 4$ and $\Pi_0 = 0$. Segment memory locations 0 to ($\Pi_4 - 1$) into four mutually exclusive sets such that set 1 consists of locations 0 to ($\Pi_1 - 1$), set 2 comprises locations Π_1 to ($\Pi_2 - 1$), and in general the s^{th} set occupies memory locations Π_{s-1} to ($\Pi_s - 1$) for $s = 1, 2, 3, 4$. In each set s for $s = 1, 2, 3, 4$ each y_i is stored in δ_{si} locations for $i = 1, 2, \dots, n$, a total of $\Pi_s - \Pi_{s-1}$ locations. The total memory requirement is then $\Pi_4 = \sum_{j=1}^4 \sum_{i=1}^n \delta_{ji}$ locations. Now choose set 1 with probability P_1 , set 2 with probability P_2 , set 3 with probability P_3 , or

set 4 with probability P_4 . Having chosen a set, select at random a location within that set; the number occupying that location is the desired random variable.

This procedure gives the required discrete distribution, as can be seen by defining:

$$q_{ri} = \delta_{ri} / \sum_{i=1}^n \delta_{ri}$$

for $i = 1, 2, \dots, n$ and $r = 1, 2, 3, 4$. Thus q_{ri} is simply the probability of choosing y_i from set r . Then the probability of generating $Y = y_i$ is

$$\sum_{r=1}^4 p_r \cdot q_{ri} = \sum_{r=1}^4 \beta^{-r} \delta_{ri} = \delta_{1i} \delta_{2i} \delta_{3i} \delta_{4i},$$

which is the probability $P(Y = y_i)$.

In order to select the proper set and location within that set, generate a uniform $(0, 1)$ random variable $u = .d_1 d_2 \dots$, and let $A\{j\}$ denote the contents of memory location j . Then if

$$\sum_{r=0}^{s-1} p_r \leq u \leq \sum_{r=0}^s p_r$$

put

$$Y = A\{d_1 d_2 \dots d_s + \pi_{s-1} - \beta^s \sum_{r=0}^{s-1} p_r\}.$$

The following example illustrates this procedure. The hexadecimal ($\beta=16$) number system is used as it is the representation of the IBM 360/65,

i	y_i	p_i
1	0	0.1000
2	1	0.3800
3	2	0.14A2
4	3	0.0F50
5	4	0.0202
6	5	0.020C

Defining P_r and Π_s as before gives

$$P_0 = 0$$

$$P_1 = 16^{-1} \sum_{i=1}^6 \delta_{1i} = 0.E$$

$$P_2 = 16^{-2} \sum_{i=1}^6 \delta_{2i} = 0.1F$$

$$P_3 = 16^{-3} \sum_{i=1}^6 \delta_{3i} = 0.00F$$

$$P_4 = 16^{-4} \sum_{i=1}^6 \delta_{4i} = 0.0010$$

and

$$\Pi_0 = 0$$

$$\Pi_1 = \sum_{j=1}^1 \sum_{i=1}^6 \delta_{ji} = E$$

$$\Pi_2 = \sum_{j=1}^2 \sum_{i=1}^6 \delta_{ji} = 2D$$

$$\Pi_3 = \sum_{j=1}^3 \sum_{i=1}^6 \delta_{ji} = 3C$$

$$\Pi_4 = \sum_{j=1}^4 \sum_{i=1}^6 \delta_{ji} = 4C.$$

The four sets are stored as follows:

A - TABLE

SET 1		SET 2				SET 3		SET 4	
Loc.	Con.	Loc.	Con.	Loc.	Con.	Loc.	Con.	Loc.	Con.
0	0	E	1	1E	3	2D	2	3C	2
1	0	F	1	1F	3	2E	2	3D	2
2	0	10	1	20	3	2F	2	3E	4
3	0	11	1	21	3	30	2	3F	4
4	0	12	1	22	3	31	2	40	5
5	0	13	1	23	3	32	2	41	5
6	0	14	1	24	3	33	2	42	5
7	0	15	1	25	3	34	2	43	5
8	0	16	2	26	3	35	2	44	5
9	0	17	2	27	3	36	2	45	5
A	1	18	2	28	3	37	3	46	5
B	1	19	2	29	4	38	3	47	5
C	1	1A	3	2A	4	39	3	48	5
D	2	1B	3	2B	5	3A	3	49	5
		1C	3	2C	5	3B	3	4A	5
		1D	3					4B	5

TABLE 2,1

To generate the desired random variable, let $u = .d_1d_2 \dots$ be a uniform $(0,1)$ random number and let $A(j)$ denote the contents of memory location j and proceed as follows:

- 1) If $0 \leq u < 0.E$ put $Y = A\{d_1\}$
- 2) If $0.E0 \leq u < 0.FF$ put $Y = A\{d_1d_2 - D2\}$
- 3) If $0.FF0 \leq u < 0.FFF$ put $Y = A\{d_1d_2d_3 - FC3\}$
- 4) If $0.FFF0 \leq u$ put $Y = A\{d_1d_2d_3d_4 - FFB4\}.$

Examples:

$$u = 0.2170 \dots Y = A\{2\} = 0$$

$$u = 0.EF10 \dots Y = A\{EF-D2\} = A\{1D\} = 3$$

$$u = 0.FFFE \dots Y = A\{FFFE - FFB4\} = A\{4A\} = 5.$$

In this example $4C_{(16)}$ memory locations are required as compared to 16^4 memory locations for the fastest method. Suppose the times for certain operations in the computer are:

<u>Operation</u>	<u>Time</u>
Compare two integers	P
Subtract two integers	S
Look up an addressed location	L

Then the fastest method for generating Y requires a total time of $P + L$ and 16^4 memory locations. In the example presented only $4C_{(16)}$ memory locations are required and an average generation time of

$$0.E + (P + L) + 0.F + (2P + S + L) + 0.00F + (3P + S + L) + 0.001$$

$$\cdot (3P + S + L) \approx 1.21_{(16)} + P + 0.20_{(16)} + S + L$$

A schematic for storing the four sets of discrete variables and the subsequent generation of the desired random variables is presented in figure 2.1.

FIGURE 2.1

Flowchart for the Generation of Random
Numbers from a Discrete Distribution

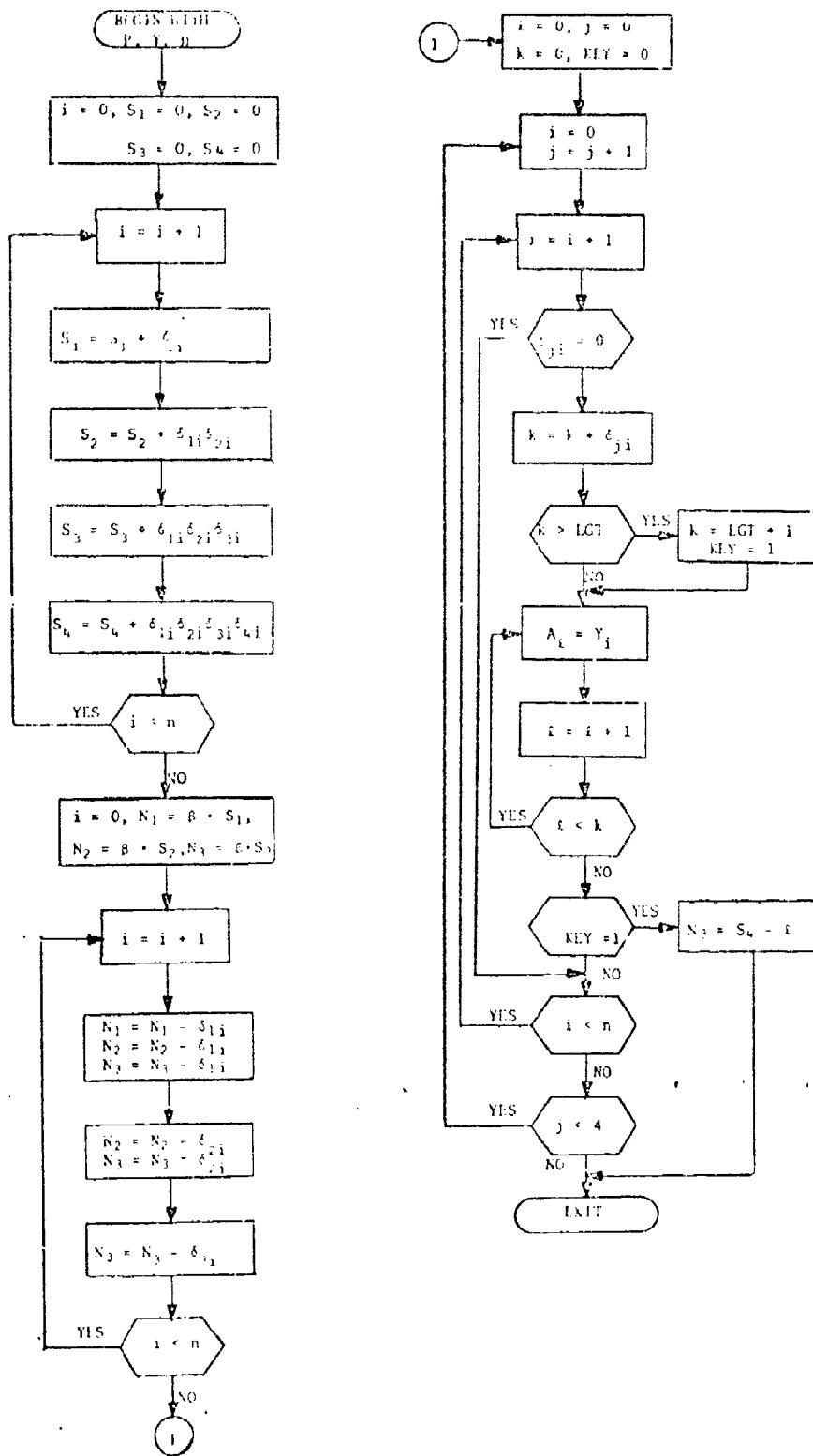
P is a vector of probabilities ($p_i = \delta_{1i} \delta_{2i} \delta_{3i} \dots$)

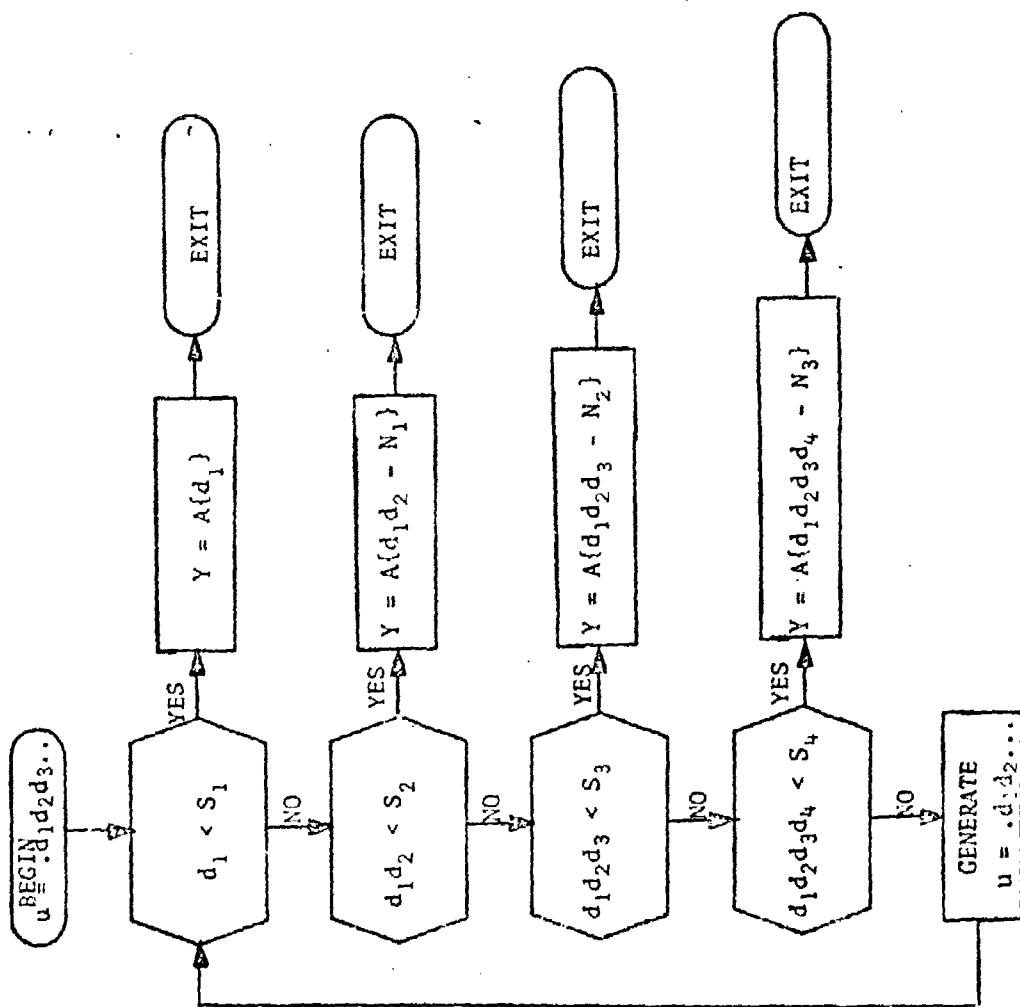
Y is a vector of the corresponding discrete variables (integer mode)

n is the number of elements in P and in Y.

β is the base of the number system in which the δ_{ij} 's are represented
(for the 360/65 $\beta = 16$, for a decimal machine $\beta = 10$).

LGT is the maximum length of the A-Table. If LGT is exceeded, the table
loading process ceases and N_3 is adjusted (since N_3 was calculated on the
basis that loading would be complete).





CHAPTER 3
GENERAL TECHNIQUES FOR THE GENERATION
OF PSEUDO-RANDOM NUMBERS HAVING A
CONTINUOUS DISTRIBUTION FUNCTION

Certain techniques of a rather general nature are used in some of the routines described in this paper. They are especially well suited for use in very fast computer programs. The principle of each technique is described below and reference will be made to these techniques as they are used in the specific generator programs.

3.1 The Composition Technique

The fundamental procedure used for generating pseudo-random numbers from a continuous probability distribution, as suggested by Marsaglia [10], is based upon a decomposition of the density function f into a mixture of 3 densities:

$$f(t) = p_1g_1(t) + p_2g_2(t) + p_3g_3(t) \quad (3.1.1)$$

where $p_1 > p_2 > p_3$, $p_1 + p_2 + p_3 = 1$ and p_1 is very close to 1.

A random number from $f(t)$ is obtained as either a number from $g_1(t)$, from $g_2(t)$ or from $g_3(t)$ with respective probabilities p_1 , p_2 , p_3 . Consider a density $f(t)$ defined from all $t \geq a$ (figure 3.1).

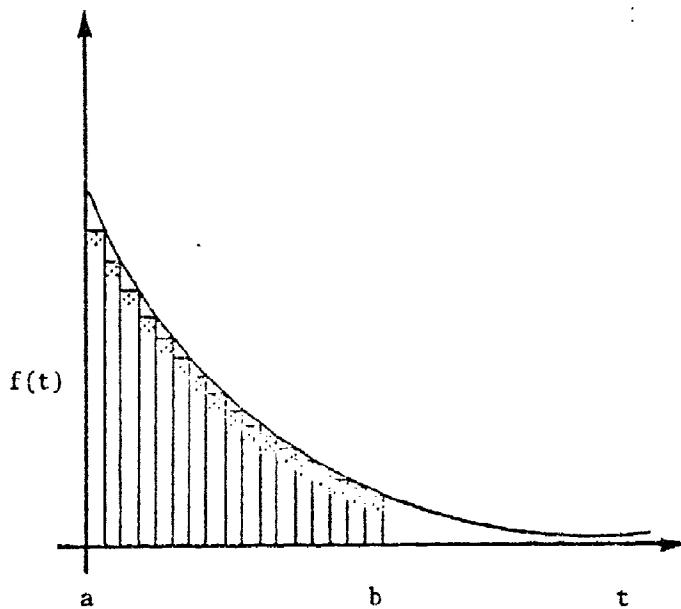


FIGURE 3.1

The density $g_1(t)$ is represented by the rectangles, appropriately standardized, including the shaded upper portions, and is defined for $a \leq t < b$. The width of the rectangles is Δ , a quantity whose value is dictated by the number system used within the computer. For a binary machine using hexadecimal arithmetic, $\Delta = 0.1_{(16)} = 0.0625_{(10)}$. The total area represented by the rectangles is p_1 which clearly is very close to 1. The density $g_2(t)$ is represented by the "triangular" regions, appropriately standardized, lying below $f(t)$ and above $g_1(t)$, and is also defined for $a \leq t < b$. The total area occupied by the "triangles" is p_2 . The density $g_3(t)$ is defined for $t \geq b$ and represents

the tail of $f(t)$. Its area is p_3 .

Random variables having density $g_1(t)$ may be rapidly generated as follows:

Define

$$N = (b-a)/\Delta \quad (3.1.2)$$

and

$$t_i = a + (i-1) \cdot \Delta \quad \text{for } i = 1, 2, 3, \dots, N. \quad (3.1.3)$$

Now assign each discrete t_i a probability

$$P_i = \Delta \cdot f(t_{i+1}) = \delta_{1i} \delta_{2i} \dots \quad (3.1.4)$$

The P_i 's are simply the areas of the individual rectangles comprising $g_1(t)$. Based on (say) the high order 4 digits of the P_i 's store the discrete t_i 's in a table according to the technique described in Chapter 2 of this paper. Thus to generate a random variable Y from this portion of g_1 (unshaded in figure 3.1), generate a uniform $(0,1)$ random number $u = .d_1 d_2 \dots$. If

$$u < \sum_{i=1}^N \delta_{1i} \delta_{2i} \delta_{3i} \delta_{4i}$$

then allow the 4 high order digits of u , $d_1 d_2 d_3 d_4$, to locate a particular t_i from the discrete table, as described in Chapter 2, and set $Y = t_i + \Delta \cdot (.d_5 d_6 \dots)$.

Random variables from the residual of g_1 , corresponding to the remaining digits of the P_i 's ($.0000 \delta_{5i} \delta_{6i} \dots$) and represented by the shaded region in figure 3.1, are generated according to the following scheme:

Let N , t_i , and P_i be defined as before (3.1.2), (3.1.3), and (3.1.4); and define:

$$H_i = \delta_{1i} \delta_{2i} \delta_{3i} \delta_{4i} \quad (3.1.5)$$

(area of the unshaded portion of the i^{th} rectangle),

$$R_i = P_i - H_i \quad (3.1.6)$$

(area of the shaded portion of the i^{th} rectangle),

$$T_i = F(t_{i+1}) - F(t_i) - P_i \quad (3.1.7)$$

where

$$F(t_i) = \int_a^{t_i} f(t) dt.$$

(T_i is the area of the "triangular" region above the i^{th} rectangle), and

$$P'_i = \sum_{i=1}^N H_i \quad (3.1.8)$$

(total area occupied by the unshaded rectangles). It is obvious from the above definitions that

$$P_2 = \sum_{i=1}^N T_i, \quad (3.1.9)$$

$$P_1 = \sum_{i=1}^N P_i, \text{ and} \quad (3.1.10)$$

$$P_1 - P'_1 = \sum_{i=1}^N R_i \quad (3.1.11)$$

(See (3.1.1) for the significance of P_1 and P_2). Now store the t_i 's in N consecutive locations. The order in which the t_i 's are stored is not important*. However, for convenience let $D(1)$ denote the t -value occupying

* They may be overlapped with the previously stored t -values in order to save storage space.

the first location, $D(2)$ denote the t -value occupying the second location, etc. For each t_i define a pair of values $C(k)$ and $B(k)$ as follows:

$$C(1) = p_1' + R[D(1)], \quad (3.1.12)$$

$$B(k) = C(k) + T[D(k)], \text{ and} \quad (3.1.13)$$

$$C(k+1) = C(k) + R[D(k+1)] \quad (3.1.14)$$

for $k = 1, 2, \dots, N$. $R[D(k)]$ and $T[D(k)]$ denote respectively the area of the rectangle residue (shaded in figure 3.1) and the area of the "triangular" region corresponding to the t -value occupying location $D(k)$. It is apparent that

$$B(N) = p_1 + p_2. \quad (3.1.15)$$

If u is a uniform $(0,1)$ random number such that $p_1' \leq u < p_1 + p_2$, a random variable Y is generated from either the rectangular residues of g_1 or the "triangular" regions of g_2 as follows:

If $B(k-1) \leq u < C(k)$ then generate a new uniform $(0,1)$ random number v and set $Y = D(k) + \Delta \cdot v$. A random variable generated in this fashion will have the density represented by the shaded rectangles in figure 3.1. If $C(k) \leq u < B(k)$, then the random variable Y must be generated from the "triangular" density of g_2 corresponding to the t -value occupying location $D(k)$. The method used depends upon the particular function $f(t)$. The same applies for random variables from g_3 for $t > b$.

The techniques used for the exponential and normal densities will be described subsequently.

Since p_1 is close to 1, most of the time a random variable from g_1 is generated. Though g_2 and g_3 may be complicated and require longer running time, the average time per generation is small since they must be handled so rarely.

3.2 The Acceptance-Rejection Method

Consider the density $h(t)$ defined on the interval (ξ_0, ξ_1) . See figure 3.2. Now if u, v is a pair

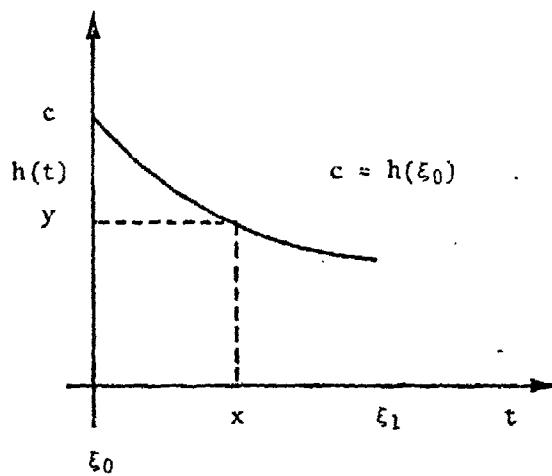


FIGURE 3.2

of independent uniform $(0,1)$ random numbers, $x = \xi_0 + u \cdot (\xi_1 - \xi_0)$ and $y = v \cdot c$ will have uniform densities on (ξ_0, ξ_1) and $(0, c)$ respectively. Suppose that $y \leq h(x)$. Then the conditional distribution function of x is given by

$$P[X \leq x \mid Y \leq h(x)] = \int_{\xi_0}^x \int_0^{h(t)} dy dt = \int_{\xi_0}^x h(t) dt. \quad (3.2.1)$$

Consequently the conditional density of x given $y \leq h(x)$ is $h(x)$.

When a pair of independent uniform $(0,1)$ numbers, u and v , are generated and the resulting x and y satisfy the inequality $y \leq h(x)$, x is then a required random variate from $h(x)$. Otherwise the pair is rejected and a new pair is generated and checked. It is clear that the inequality will be satisfied with a high probability only if the value $c \cdot (\xi_1 - \xi_2)$ is close to one. If not, the procedure will produce a large proportion of inadmissible (u,v) pairs and reduce the efficiency of the scheme. Even when this proportion is small, the time required to check $y \leq h(x)$ is generally rather long compared with that of the table look-up procedures. Consequently, the acceptance-rejection method is best used for generating those infrequent values from the "triangular" and tail regions described in section 3.1.

3.3 Generation of Numbers with a Triangular Density

Let u and v be independent uniform $(0,1)$ random numbers. Then $T = \min(u, v)$ has the distribution function

$$G(t) = 1 - P[T > t] = 1 - P[u > t, v > t] = 1 - (1-t)^2, \quad 0 < t < 1, \quad (3.3.1)$$

and the density function

$$g(t) = 2(1-t), \quad 0 < t < 1. \quad (3.3.2)$$

See figure 3.3a. A linear transformation $T' = a + bT$ produces a random variable with density

$$g(t') = \frac{2}{b} - \frac{2}{b^2} (t' - a), \quad a < t' < a+b \quad (3.3.3)$$

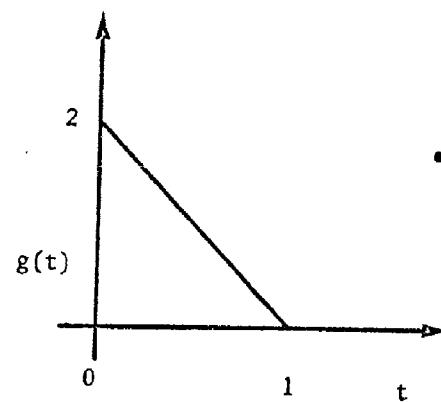


FIGURE 3.3(a)

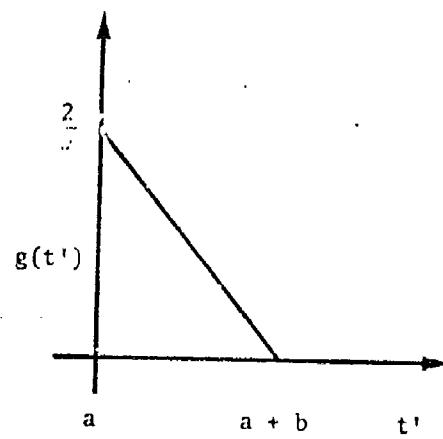


FIGURE 3.3(b)

3.4 The Distribution of the Minimum of a Random Number of Uniform (0,1) Random Variates: Some Special Results for the Exponential Distribution

Let $x = c \cdot \min(u_1, u_2, \dots, u_n)$ where the u_i are random independent uniform (0,1) variates, c is a positive constant, and N is a discrete valued random variable with probability function $P(N=n) = q(n)$ for $n = 1, 2, \dots$. Then the distribution function of x can be expressed as follows:

$$F(x) = \sum_{n=1}^{\infty} P[X \leq x \mid N = n] \cdot P[N = n]$$

$$= \sum_{n=1}^{\infty} [1 - (1 - x/c)^n] \cdot q(n)$$

$$F(x) = 1 - \sum_{n=1}^{\infty} (1 - x/c)^n \cdot q(n) \quad 0 < x < c \quad (3.4.1)$$

Now consider the special case

$$q(n) = c^n / [n! (e^c - 1)] \quad n = 1, 2, \dots$$

The distribution of x becomes

$$F(x) = 1 - \sum_{n=1}^{\infty} (c-x)^n / n! (e^c - 1) = 1 - (e^{c-x} - 1) / (e^c - 1)$$

$$F(x) = (1 - e^{-x}) / (1 - e^{-c}) \quad 0 < x < c \quad (3.4.3)$$

Thus x has the exponential distribution truncated on the right at the value c .

If, instead of the probability function given in (3.4.2), N has the distribution specified by

$$g(n) = \Delta^n / n! (e^\Delta - 1 - \Delta) \quad n = 2, 3, \dots, \quad (3.4.4)$$

where Δ is a constant as defined in section 3.1, then the distribution function of x becomes

$$F(x) = 1 - (e^{\Delta-x} - 1 - \Delta + x) / (e^\Delta - 1 - \Delta), \quad 0 < x < \Delta, \quad (3.4.5)$$

with density

$$f(x) = (e^{-x} - e^{-\Delta}) / [1 - e^{-\Delta}(1+\Delta)], \quad 0 < x < \Delta. \quad (3.4.6)$$

This is the distribution function of a random variable from the "tooth" of the exponential distribution between 0 and Δ . See figures 3.4(a) and 3.4(b).

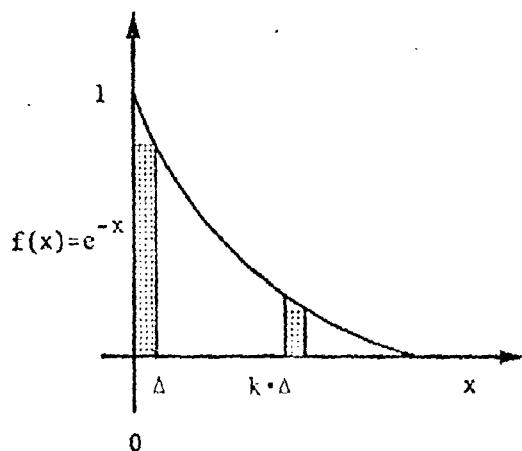


FIGURE 3.4(a)

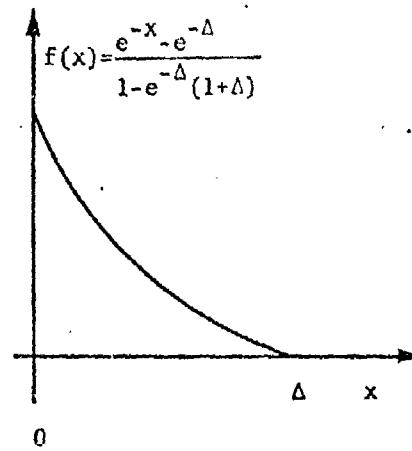


FIGURE 3.4(b)

Figure 3.4(c) shows the density of a random variable from any other such "tooth" defined by $k \cdot \Delta < z < (k+1) \cdot \Delta$.

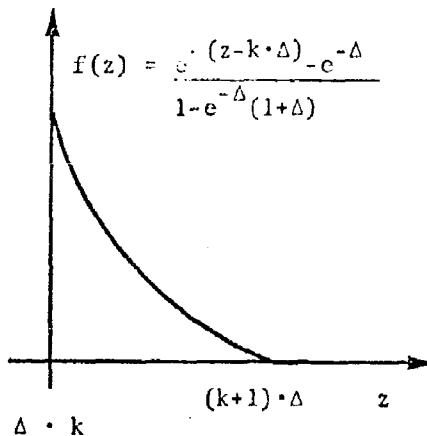


FIGURE 3.4(c)

It is obvious that all such densities are identical, a unique characteristic of the exponential distribution. Consequently, a random variable from the $(k+1)^{\text{st}}$ "tooth" is always produced by taking $z = k \cdot \Delta + x$ where the value $k \cdot \Delta$ is chosen by the technique described in section 3.1 and x is generated by taking Δ times the minimum of N uniform $(0,1)$ random numbers where N has the distribution (3.4.4). Its expected value when $\Delta = 0.1$ (16) is 2.02 or in general

$$E[N] = \Delta \cdot (e^\Delta - 1) / (e^\Delta - 1 - \Delta) \quad (3.4.7)$$

The same unique "no memory" characteristic of the exponential distribution used above is also used for obtaining random variates from g_3 , the tail of the exponential density where $t > c$.

Let $Y = W + X$ where W is discrete valued with probability function

$$P[W = w = c + k] = Ae^{-w} = Ae^{-c-k}, \quad k = 1, 2, 3, \dots, \quad (3.4.8)$$

$A = (1 - e^{-c})/e^{-c}$ and x is independent of w and has the distribution function (3.4.3). The distribution function of Y can be written as

$$G(y) = P[Y \leq y = w + x] = P[Y \leq w + c] + P[W = w, X \leq x] \quad (3.4.9)$$

Since the distribution function of W is

$$P[W \leq w = c + k] = A \sum_{j=1}^k e^{-j \cdot c} = 1 - e^{-c \cdot k}$$

the expression for $G(y)$ becomes

$$G(y) = 1 - e^{-(w-c)} + Ae^{-w}(1-e^{-x})/(1-e^{-c}) = 1 - e^{c-y}, \quad (3.4.11)$$

$$y > c.$$

This is the required distribution function for variates from the density e^{-y} truncated on the left at $y = c$. Thus values of t from g_3 in the exponential case are generated as the sum of a discrete variate taking values $c, 2c, 3c, \dots$ and a continuous variate from the truncated exponential density on $(0, c)$.

3.5 Generating Numbers from a 'Nearly Triangular' Density: Application to the Half Normal Density

The results of this section are applied to the half normal density

$$f(t) = \sqrt{2/\pi} e^{-\frac{1}{2}t^2} \quad 0 < t < \infty \quad (3.5.1)$$

for the generation of random normal deviates to be described in chapter 4. Consider the concave and convex triangular densities (members of g_2) depicted in figures 3.5(a) and 3.5(b).

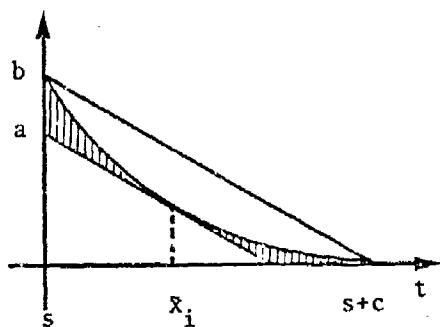


FIGURE 3.5(a)

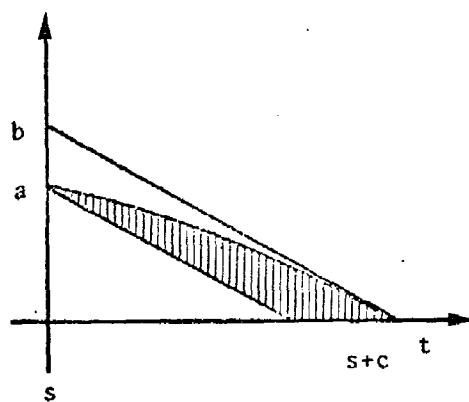


FIGURE 3.5(b)

The parallel chords and tangents enclosing $f(t)$ are determined by construction as indicated in the figures. In both cases, the inner right triangle represents most of the area under the density $F(t)$. This will be true in general as long as the ratio a/b is close to one. A random variate from the density represented by this triangle is generated as described in section 3.3. Infrequently, a random variate from the shaded area must be generated by the acceptance-rejection technique described in section 3.2. Marsaglia et al. [12] have combined these two procedures as follows.

- 1) Generate independent uniform $(0,1)$ random variables u and v .
- 2) If $\max(u,v) \leq a/b$, put $t = s + c \cdot \min(u,v)$
- 3) If not, test $b|u-v| \leq f(s + c \cdot \min(u,v))$.
 - If yes, put $t = s + c \cdot \min(u,v)$.
 - If no, go to step 1 and try again.

In order to show that this procedure produces a variate t with the required density, we conveniently take $s = 0$ without loss of generality and proceed as follows.

Define

$$m = \min(u,v)$$

$$M = \max(u,v)$$

$$x = c \cdot m$$

$$y = b(M-m) = b(u-v) \quad (3.5.2)$$

Then the pair (x,y) is uniformly distributed over the triangle in figure 3.5(c)

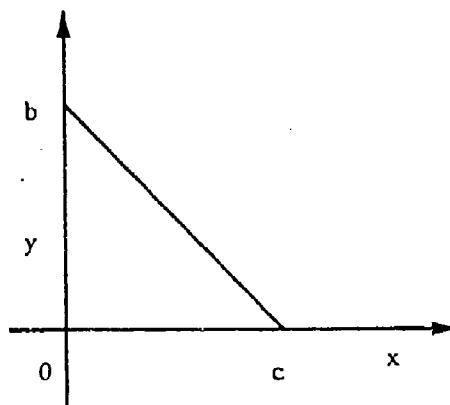


FIGURE 3.5(c)

To verify this, consider the joint distribution function of X and Y :

$$F(x,y) = P[X \leq x, Y \leq y] = P[m \leq x/c, M - m \leq y/b]$$

This is just the area of the cross hatched region in figure 3.5(d) which is seen to be

$$F(x,y) = 2 \cdot (y/b) \cdot (x/c) = 2xy/bc, \quad 0 < x < c, \quad (3.5.3)$$

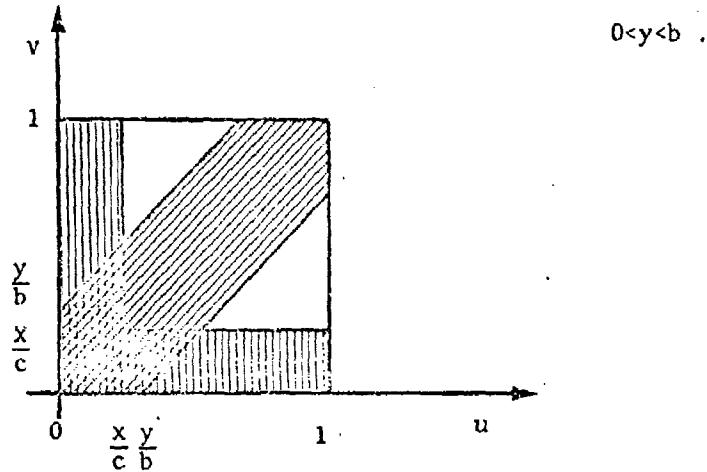


FIGURE 3.5(d)

Thus, (x,y) is a point randomly chosen from the larger triangle in figures 3.5(a) and 3.5(b). By the acceptance-rejection technique, if this point falls within the smaller triangle of figures 3.5(a) and 3.5(b), the value $x = t$ has the triangular density $-bx/c + a$. Now the condition $\max(u,v) \leq a/b$ is equivalent to:

$$bM \leq a$$

$$b(M-m) \leq a - bm = -bx/c + a$$

or

$$y \leq -bx/c + a \quad (3.5.4)$$

which means that (x,y) falls in the small triangle. If this condition is not satisfied, a second check is made to determine whether the point (x,y) falls within the shaded area between $f(t)$ and y_2 in figures 3.5(a) and 3.5(b). If so, then clearly

$$y = b(M - m) = b|u - v| \leq f(x) = f(c + \min(u,v)) \quad (3.5.5)$$

as required by the second test. The acceptance-rejection principle insures that every x value passing at least one of these two tests will have the density $f(t)$. The first test is rapidly made; the second test involves evaluation of $f(x)$ and is time consuming. The first test will be passed with a probability

$$P[M \leq a/b] = (a/b)^2$$

as given by the distribution function of $M = \max(u,v)$. It is therefore desirable to have $(a/b)^2$ close to 1 to achieve a short average execution time. This will be the case if $f(t)$ is "nearly triangular."

For the half normal density the a_i/b_i ratios are obtained as follows.

Case (a) "Concave Triangles" (See figure 3.5(a)).

Let $t_i = i \cdot c = s$

$$t_{i+1} = (i+1)c = s + c, \quad i = 0, 1, 2, \dots \quad (3.5.6)$$

The equation of the chord and the tangent are respectively:

$$\begin{aligned}y_1 &= -b_i(s)/c + b_i \\y_2 &= -b_i(x-s)/c + a_i\end{aligned}\quad (3.5.7)$$

The value of b_i is clearly

$$b_i = f(t_i) - f(t_{i+1}) = f(s) - f(s+c) \quad (3.5.8)$$

The slope of both lines, $-b_i/c$, must equal $f'(\bar{x}_i)$. Thus

$$-b_i/c = -\sqrt{2/\pi} x_i e^{-\frac{1}{2}x_i^2} = -\bar{x}_i f'(\bar{x}_i) \quad (3.5.9)$$

This equation must be solved iteratively for \bar{x}_i , the abscissa of the point of tangency. Finally, this value is substituted into the tangent's equation to give

$$f(\bar{x}_i) = f(t_{i+1}) = (-b_i/c)(\bar{x}_i - s) + a_i. \quad (3.5.10)$$

Substituting the expression for b_i/c and solving for a_i gives

$$a_i = f(\bar{x}_i)[1 + \bar{x}_i(\bar{x}_i - s)] - f(s+c). \quad (3.5.11)$$

Consequently for "concave triangles"

$$a_i/b_i = [f(\bar{x}_i)[1 + \bar{x}_i(\bar{x}_i - s)] - f(s+c)]/[f(s) - f(s+c)] \quad (3.5.12)$$

Case (b) "Convex Triangles" (See figure 3.5(b))

In this case the point of tangency occurs at the end of the interval where $x_i = t_{i+1} = s + c$. The slope of both lines is given by

$$-b_i/c = f'(s+c)$$

$$\text{or } b = -cf'(s+c). \quad (3.5.13)$$

The value of a_i is simply

$$f(t_i) - f(t_{i+1}) = f(s) - f(s+c).$$

Consequently, for "convex triangles"

$$a_i/b_i = [f(s) - f(s+c)]/[-cf'(s+c)] \quad (3.5.14)$$

3.6 Generation of Random Variates from the Tail of the Half Normal Distribution

The procedure suggested by Marsaglia et al. [12] for the generation of random variates from the tail of the half normal distribution is based on the rejection principle described in section 3.2. The procedure is to generate a pair of independent half normal variates x_1 and x_2 such that the point (x_1, x_2) lies outside the quarter circle (see figure 3.6(a)).

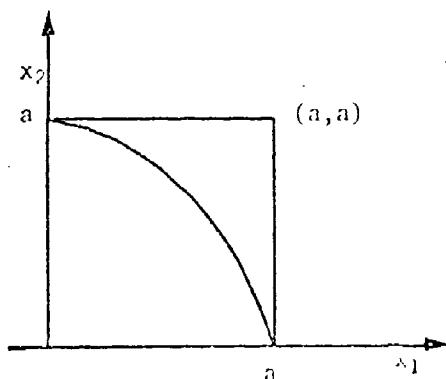


FIGURE 3.6(a)

This is done by generating pairs of uniform (0,1) random numbers u_1 and u_2 until $0 < u_1^2 + u_2^2 \leq 1$ and setting

$$x_1 = \sqrt{a^2 - 2\ln(u_1^2 + u_2^2)} \cdot \sqrt{u_1^2/(u_1^2 + u_2^2)} \quad (3.6.1)$$

$$x_2 = \sqrt{a^2 - 2\ln(u_1^2 + u_2^2)} \cdot \sqrt{u_2^2/(u_1^2 + u_2^2)} \quad (3.6.2)$$

To show that the point (x_1, x_2) lies outside the quarter-circle, let $w = -\ln(u_1^2 + u_2^2)$ where $0 < u_1^2 + u_2^2 \leq 1$. Then take the sum $x_1^2 + x_2^2 = a^2 + 2w$ where $0 \leq w < \infty$ so that $x_1^2 + x_2^2 \geq a^2$. If the point (x_1, x_2) , generated in this fashion, lies within the square (figure 3.6(a)), it is rejected and a new pair of uniform (0,1) random numbers is generated and tested. Otherwise the variable, x_1 or x_2 , whose value exceeds a is taken as the required random number.

In order to show that this procedure produces the desired result, let u_1 and u_2 be a pair of uniform (0,1) random numbers conditioned by $0 < u_1^2 + u_2^2 \leq 1$. Then the joint density of u_1 and u_2 is

$$f(u_1, u_2) = 4/\pi, \quad 0 < u_1 < 1, \quad 0 < u_2 < 1, \quad 0 < u_1^2 + u_2^2 \leq 1. \quad (3.6.3)$$

Now define:

$$X = u_1^2 + u_2^2, \quad 0 < X < 1, \text{ and} \quad (3.6.4)$$

$$Y = u_2/u_1, \quad 0 < Y < \infty. \quad (3.6.5)$$

Then the probability distribution of X is

$$F(x) = P(X \leq x) = (\pi + x^2/4)/(\pi + 1^2/4) = x \quad (3.6.6)$$

and the density of x is

$$f(x) = 1, \quad 0 < x < 1. \quad (3.6.7)$$

Consequently X is uniform on $(0,1)$. The joint density of X and Y is:

$$g(x,y) = f(u,v) \cdot |J| = (4/\pi)[1/2(1+y^2)] = 2/\pi(1+y^2) = g_1(x) \cdot g_2(y). \quad (3.6.8)$$

Hence X and Y are stochastically independent. Substitution of (3.6.4) and (3.6.5) into (3.6.1) and (3.6.2) yields:

$$x_1 = \sqrt{a^2 - 2\ln X} + \sqrt{1/(1+Y^2)} \quad (3.6.9)$$

$$x_2 = \sqrt{a^2 - 2\ln X} - \sqrt{Y^2/(1+Y^2)}. \quad (3.6.10)$$

The inverse transformation is:

$$X = e^{-\frac{1}{2}(x_1^2+x_2^2-a^2)} \quad (3.6.11)$$

$$Y = x_2/x_1 \quad (3.6.12)$$

The joint density of x_1 and x_2 is given by:

$$\begin{aligned} f(x_1, x_2) &= g(x, y) \cdot |J| \\ &= [2/\pi(1+y^2)](1+y^2) \cdot e^{-\frac{1}{2}(x_1^2+x_2^2-a^2)} \\ &= (\sqrt{2/\pi} e^{-\frac{1}{2}x_1^2})(\sqrt{2/\pi} e^{-\frac{1}{2}x_2^2}) e^{a^2/2} \quad (3.6.13) \\ &\quad a^2 \leq x_1^2 + x_2^2 < \infty. \end{aligned}$$

Clearly x_1 and x_2 are independent random variables from the tail of the half normal distribution ($x_1, x_2 > a$).

CHAPTER 4
GENERATION OF PSEUDO-RANDOM NUMBERS FROM THE
EXPONENTIAL, GAMMA, AND NORMAL DISTRIBUTIONS

Before presenting the procedures for the generation of random numbers from the exponential, gamma, and normal distributions, a simple example is given demonstrating the use of the techniques discussed in Chapter 3 for obtaining random variates from any continuous distribution defined on a finite interval. In the interest of clarity all numbers are expressed to the base 10. Consider the distribution defined by

$$F(y) = (y/2)(3-y^2) \quad 0 < y < 1 \quad (4.0.1)$$

with density

$$f(y) = (3/2)(1-y^2) \quad 0 < y < 1. \quad (4.0.2)$$

Using the composition technique presented in 3.1, $f(y)$ is represented as a mixture of 2 densities (since there is no tail):

$$f(y) = p_1 g_1(y) + p_2 g_2(y). \quad (4.0.3)$$

As before, g_1 is a series of rectangles and g_2 represents the nearly triangular regions between g_1 and f . See figures 4.0(a), (b), and (c).

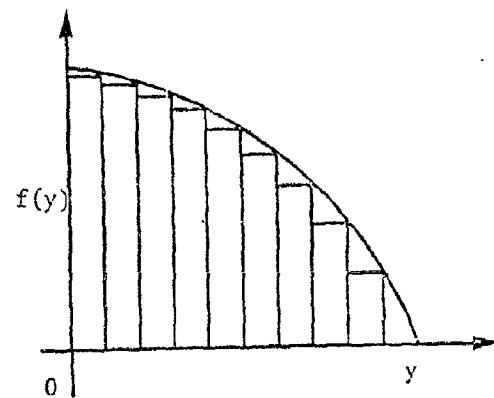


FIGURE 4.0(a)

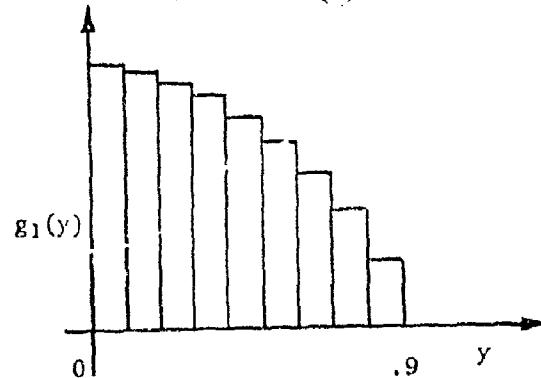


FIGURE 4.0(b)

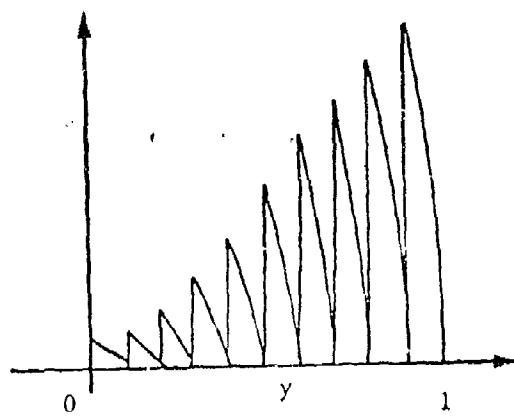


FIGURE 4.0(c)

As described in 3.1, define discrete variables

$$x_i = (i-1) \cdot \Delta \text{ for } i=1, 2, \dots, 10 \text{ and in this example } \Delta=0.1_{(10)}. \quad (4.0.4)$$

Now assign each x_i a probability P_i as follows:

$$P_i = \Delta \cdot f(x_{i+1}). \quad (4.0.5)$$

Now define a value T_i as given in (3.1.7), the area of the triangular region above the i^{th} rectangle:

$$T_i = F(x_{i+1}) - F(x_i) - P_i. \quad (4.0.6)$$

These values are recorded in table 4.0.1. At this point it is obvious that the values of p_1 and p_2 , defined in (4.0.3), are $p_1 = 0.9225$ and $p_2 = 0.0775$. Next, the four sets of the discrete values are stored, based on the high order of digits of the P_i 's.

x_i	$p_i = 0.1 \cdot f(x_{i+1})$	$F(x_{i+1}) - F(x_i)$	$T_i = F(x_{i+1}) - F(x_i) - p_i$
0.0	0.1485	0.1495	0.0010
0.1	0.1440	0.1465	0.0025
0.2	0.1365	0.1405	0.0040
0.3	0.1260	0.1315	0.0055
0.4	0.1125	0.1195	0.0070
0.5	0.0960	0.1045	0.0085
0.6	0.0765	0.0865	0.0100
0.7	0.0540	0.0655	0.0115
0.8	0.0285	0.0415	0.0130
0.9	0.0000	0.0145	0.0145
Total	0.9225	1.0	0.0775

TABLE 4.0.1

Since the P_i 's have only 4 digits, application of this technique will account for all of g_1 and there is no rectangular residue in this simple example. Since the composition of these 4 sets is constant (the distribution is parameter free), the sets may be overlapped and considerable storage space saved. Looking at table 4.0.2, set 1 occupies locations 2 through 6, set 2 occupies locations 4 through 40, set 3 occupies locations 18 through 67, and set 4 occupies locations 46 through 70. Locations 0 through 9 are also used for the generation of variates from g_2 , as described subsequently.

TABLE 4.0.2

A-TABLE

As presented in Chapter 2, define:

$$Q_0 = 0, \quad Q_r = 10^{-r} \cdot \sum_{i=1}^{10} \delta_{ri} \quad (4.0.7)$$

and

$$\Pi_0 = 0, \quad \Pi_r = \sum_{j=1}^r \sum_{i=1}^{10} \delta_{ji} \quad \text{for } r=1,2,3,4. \quad (4.0.8)$$

Using the above definitions define:

$$S_0 = 0, \quad S_k = 10^k \sum_{r=1}^k Q_r \quad (4.0.9)$$

and

$$N_k = \Pi_{k-1} - 10 \cdot S_{k-1} \quad \text{for } k=1,2,3,4. \quad (4.0.10)$$

The N_k 's are identical to the quantity

$$\Pi_{s-1} + \beta^s \cdot \sum_{r=0}^s P_r$$

in (2.0.3). Because of overlapping between the four sets, the N_k 's must be adjusted. Their values and the values of the S_k 's are recorded in table 4.0.3.

S-TABLE		N-TABLE	
Loc	Con	Loc	Con
0	5	0	2
1	87	1	-46
2	920	2	-852
3	9225	3	-9154

Table 4.0.3

The procedure for generating random numbers from g_1 is:

- 1) Generate a uniform $(0,1)$ random number $u = 0.d_1d_2\dots$
- 2) If $d_1 < S(1)$, set $Y = A\{d_1 + N(1)\} + 0.1 \cdot (.d_5d_6\dots)$;
otherwise go to (3).
- 3) If $d_1d_2 < S(2)$, set $Y = A\{d_1d_2 + N(2)\} + 0.1 \cdot (.d_5d_6\dots)$;
otherwise go to (4).
- 4) If $d_1d_2d_3 < S(3)$, set $Y = A\{d_1d_2d_3 + N(3)\} + 0.1 \cdot (.d_5d_6\dots)$;
otherwise go to (5).
- 5) If $d_1d_2d_3d_4 < S(4)$, set $Y = A\{d_1d_2d_3d_4 + N(4)\} + 0.1 \cdot (.d_5d_6\dots)$;
otherwise generate a random variate from g_2 .

In order to generate random variates from g_2 , define:

$$C(1) = S_4 + T[A(1)] \text{ and} \quad (4.0.11)$$

$$C(k+1) = C(k) + T[A(k)] \quad \text{for } k=1,2,\dots,9. \quad (4.0.12)$$

This is identical to the definition given in (3.1.13), except that in this example there is no rectangular residue. The notation $T[A(k)]$ denotes the area of the triangular region (the T_i 's in table 4.0.1) corresponding to the discrete value stored in the k^{th} location of the A-table (Table 4.0.2). The values of the $C(k)$'s are given in table 4.0.4. If u is a uniform $(0,1)$ random number such that $u \geq S_4$, the proper triangle of g_2 may be chosen simply by testing:

$$u < C(k) \text{ until the condition is satisfied (note that } C_{(10)} = 1.0)$$

The discrete value occupying location k of the A-table denotes the correct triangle.

Having selected the proper triangle of g_2 , a random variate having the density of that triangle is to be generated. This is accomplished

through the use of the acceptance-rejection principle discussed in 3.2 and the technique given in 3.5. Consider a particular triangle or tooth from g_2 (figure 4.0(d)). The tangent and the chord are constructed as indicated in 3.5

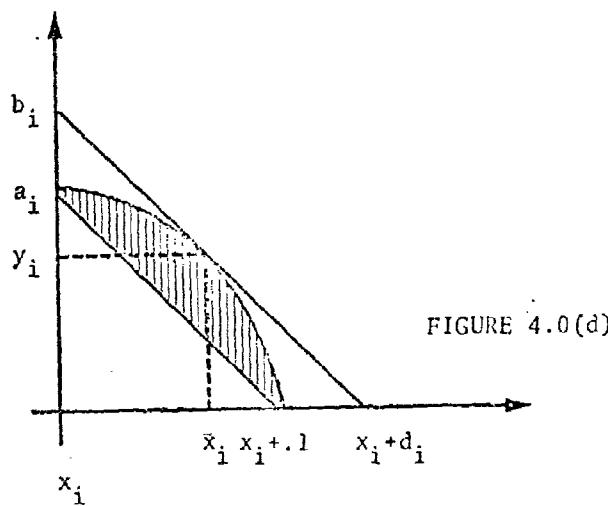


FIGURE 4.0(d)

The inner right triangle encloses most of the area under $f(x)$. A random variate from the density of this triangle is generated as described in 3.3. Occasionally, a random variate from the shaded region must be generated by the acceptance-rejection technique. The values of a_i , b_i , and d_i are obtained as follows:

$$(a_i/0.1) = f'(x_i) = 3x_i \quad (4.0.13)$$

where x_i is the abscissa of the point of tangency and

$$a_i = f(x_i) - f(x_{i+1}). \quad (4.0.14)$$

Substituting (4.0.14) into (4.0.13) and solving for \tilde{x}_i yields:

$$\tilde{x}_i = [f(x_i) - f(x_{i+1})]/0.3 \quad (4.0.15)$$

The ordinate of the point of tangency is:

$$y_i = f(\tilde{x}_i) - f(x_{i+1}). \quad (4.0.16)$$

At $x = \tilde{x}_i$, $y_1 = y_i$ and

$$b_i = y_i + (a_i/0.1)(\tilde{x}_i - x_i) \quad (4.0.17)$$

The value of d_i may be obtained by noting that:

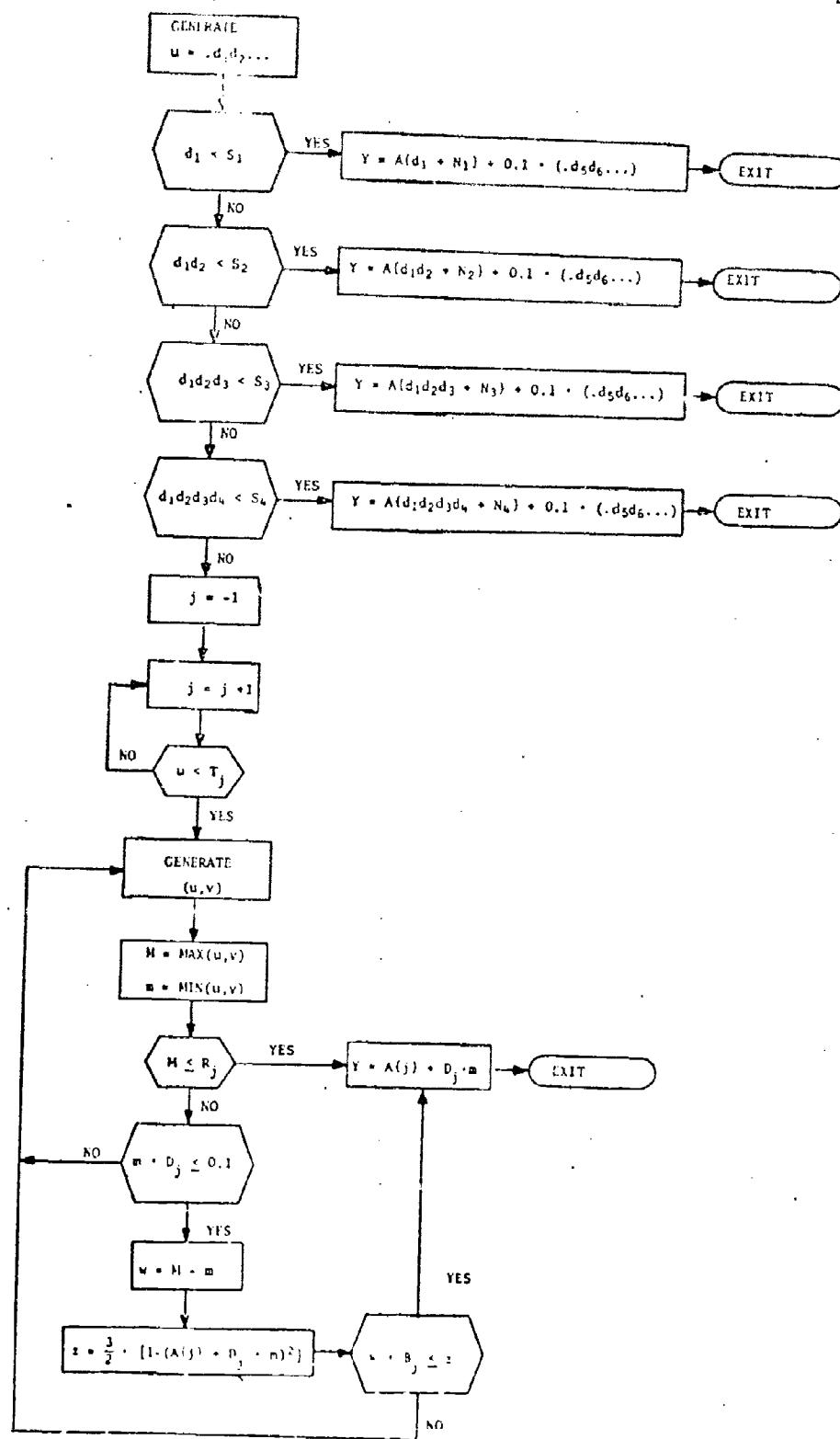
$$(a_i/0.1) = (b_i/d_i) \text{ and hence } d_i = (0.1 \cdot b_i/a_i). \quad (4.0.18)$$

The values of a_i/b_i , d_i , and b_i are recorded in table 4.0.4 in the R-table, D-table, and B-table respectively.

C-Table		R-Table		P-Table		B-Table	
Loc	Con	Loc	Con	Loc	Con	Loc	Con
0	.9370	0	.9870	0	.1013	0	.28875
1	.9500	1	.9855	1	.1015	1	.25875
2	.9555	2	.9655	2	.1036	2	.10875
3	.9580	3	.9231	3	.1083	3	.04875
4	.9650	4	.9730	4	.1028	4	.13875
5	.9690	5	.9524	5	.1050	5	.07875
6	.9700	6	.8000	6	.1250	6	.01875
7	.9815	7	.9836	7	.1017	7	.22875
8	.9915	8	.9811	8	.1019	8	.19875
9	1.0000	9	.9778	9	.1023	9	.16875

Table 4.0.4

FIGURE 4.0(e)
Flowchart for the Generation of Random
Variates from
 $f(y) = (3/2)(1 - y^2)$, $0 < y < 1$



The procedure for obtaining random variates from g_2 is:

- 1) Let u be a uniform $(0,1)$ random number such that $u \geq s_4$.
- 2) Test $u < C(k)$ for $k = 0, 1, \dots, 9$ until the inequality is satisfied.
- 3) Generate independent uniform $(0,1)$ random numbers u_1 and u_2 .
- 4) If $\max(u_1, u_2) \leq R(k)$, then set $Y = A(k) + D(k) + \min(u_1, u_2)$; otherwise go to (5).
- 5) If $B(k) + |u_1 - u_2| \leq f[A(k) + D(k) + \min(u_1, u_2)]$, set $Y = A(k) + D(k) + \min(u_1, u_2)$; otherwise go to step (3) and try again.

A schematic of the entire procedure for generating random variates from the density $f(y)$, given in (4.0.2), is presented in figure 4.0(e).

4.1 Generation of Exponentially Distributed Random Numbers

Applying the composition principle presented in 3.1, the exponential density function,

$$f(t) = e^{-t}, \quad 0 \leq t \leq \infty, \quad (4.1.1)$$

is represented as a mixture of three densities

$$f(t) = p_1 g_1(t) + p_2 g_2(t) + p_3 g_3(t). \quad (4.1.2)$$

As before, g_1 is a series of rectangles defined for $t \leq 4$, g_2 represents the toothlike region between g_1 and f defined for $t \leq 4$, and g_3 is the tail of f for $t \geq 4$. The probabilities p_1 , p_2 , and p_3 are respec-

tively the areas represented by g_1 , g_2 , and g_3 .

As described in 3.1 a random variable Y having density g_1 is generated as follows:

Define

$$M = 4/\Delta, \text{ where } \Delta = 10^{-1} \text{ for a decimal machine or } \Delta = 16^{-1} \text{ for a binary machine; and discrete variates} \quad (4.1.3)$$

$$t_i = (i-1) \cdot \Delta \text{ for } i = 1, 2, \dots, M. \quad (4.1.4)$$

Each t_i is assigned a probability

$$P_i = \Delta \cdot f(t_{i+1}) = \Delta \cdot e^{-t_{i+1}} = \delta_{1i} \delta_{2i} \delta_{3i} \dots \quad (4.1.5)$$

Now, using the high order 4 octal digits (binary machine) or the high order 3 decimal digits (decimal machine), store the t_i 's according to the technique described in Chapter 2. As in the case of the example given in the beginning of this chapter, the sets of discrete variates are stored in the manner that permits maximum overlap. See the D-table in the program listing of GEN1 in the appendix; set 1 occupies locations 49A thru 4BE, set 2 occupies locations 40C through 4AC, and set 3 occupies locations 37D through 453. Locations 374 through 3B3 are used for the generation of variates from g_2 and the residual of g_1 as described in 3.1.

At this point define:

$$Q_0 = 0, \quad Q_r = \beta^{-r} \cdot \sum_{i=1}^M \delta_{ri} \quad (4.1.6)$$

$$\Pi_0 = 0, \quad \Pi = \sum_{j=1}^r \sum_{i=1}^M \delta_{ji} \quad \text{for } r = 1, 2, 3, 4 \quad (4.1.7)$$

for binary machines where the δ 's are octal digits and $\beta = 8$ or $r = 1, 2, 3$

for decimal machines where the δ 's are decimal digits and $\beta = 10$. Also define:

$$S_0 = 0, S_k = \beta^k \cdot \sum_{r=1}^k \delta_r \quad (4.1.8)$$

and

$$N_k = \pi_{k-1} - \beta \cdot S_{k-1} \text{ for } k = 1, 2, 3, 4 \quad (4.1.9)$$

in the binary case ($\beta = 2$) or $k = 1, 2, 3$ in the decimal case ($\beta = 10$).

The S -values and the N -values are recorded in the program listing of GEN1 in the appendix. Note that the N -values have been adjusted to allow for overlap between the sets of the D-table.

The procedure for the generation of random variates from this truncated portion of g_1 is: (binary case)

- 1) Generate a uniform (0,1) random number $u = .d_1 d_2 \dots$
- 2) If $d_1 d_2 < S_2$, set $Y = D(d_1 d_2 + N_2) + \Delta \cdot (.d_5 d_6 \dots)$;
otherwise go to (3).
- 3) If $d_1 d_2 d_3 < S_3$, set $Y = D(d_1 d_2 d_3 + N_3) + \Delta \cdot (.d_5 d_6 \dots)$;
otherwise go to (4).
- 4) If $d_1 d_2 d_3 d_4 < S_4$, set $Y = D(d_1 d_2 d_3 d_4 + N_4) + \Delta \cdot (.d_5 d_6 \dots)$;
otherwise generate a random variate from the residual of g_1 , g_2 ,
or g_3 .

Note that in the above procedure the testing begins with S_2 . This is because $\delta_{1i} = 0$ for all i since $0 < f(x) < 1$ for $0 < x < 4$.

Random variates from the residual of g_1 and from g_2 are generated by first selecting one of the rectangles of g_1 with appropriate probability and then testing to determine whether the required variate should

come from the tooth or from the residual corresponding to that rectangle.

This is accomplished as follows:

In M consecutive locations store each of the t_i 's once (locations 374 through 3B3 of the D-table in the listing of GEN1 in the appendix). For each t_i define a pair of values as given in (3.1.12), (3.1.13), and (3.1.14)

$$C(1) = S_4 + R[D(1)], \quad (4.1.10)$$

$$B(k) = C(k) + T[D(k)] \text{ and} \quad (4.1.11)$$

$$C(k+1) = C(k) + R[D(k+1)]. \quad (4.1.12)$$

In the exponential case

$$T[D(k)] = e^{-D(k)} - (1 + \Delta) - e^{-(D(k)+\Delta)} \quad (4.1.13)$$

and

$$R[D(k)] = P[D(k)] - .\delta[1,D(k)]\delta[2,D(k)]\delta[3,D(k)]\delta[4,D(k)], \quad (4.1.14)$$

where $P[D(k)]$ is as defined in (4.1.5). The values $T[D(k)]$ and $R[D(k)]$ denote respectively the areas of the triangular region and rectangular residue lost upon truncation of $P[D(k)]$ corresponding to the t -value occupying location $D(k)$. The values of the $B(k)$'s and the $C(k)$'s are also recorded in the listing of GEN1 in the appendix. A random variate with a rectangular distribution is generated as described in 3.1, and in the exponential case a variate having the density of one of the teeth of g_2 is generated by the technique given in 3.4. The procedure for generating random variates from either the residual of g_1 or from g_2 is summarized as follows:

- 1) Let u be a uniform $(0,1)$ random number such that
 $S_4 \leq u < p_1 + p_2 = 1 - e^{-4}$.
- 2) Test $u < B(J)$ for $J = 1, 2, \dots, M$, until the inequality is satisfied (note $B(M) = 1 - e^{-4}$) for some $J = J'$.
- 3) If $u < C(J')$, generate a new uniform $(0,1)$ random number v and set $Y = D(J') + \Delta \cdot v$ (this produces a number from g_1); otherwise go to (4).
- 4) Generate a new uniform $(0,1)$ random number v and test
 $v < Q(n) = \sum_{j=2}^n \Delta^j / [j!(e^\Delta - 1 - \Delta)]$ for $n = 2, 3, \dots$
until the inequality is satisfied for $n = n'$.
- 5) Generate n' independent uniform $(0,1)$ random numbers and set
 $Y = D(J') + \Delta \cdot \min(u_1, u_2, \dots, u_{n'})$. This produces a number from g_2 .

A random variate Y having density g_3 is generated as a sum

$$Y = W + X \quad (4.1.15)$$

where X is exponential on $(0,4)$ and W has the distribution

$$P[W = w = 4 \cdot k] = e^{-4} e^{4k} (e^4 - 1), \quad k = 1, 2, \dots \quad (4.1.16)$$

The procedure for generating a random variate from g_3 is summarized as follows:

- 1) Let u be a uniform $(0,1)$ random number such that $u \geq p_1 + p_2 = 1 - e^{-4}$.
- 2) Test $u < 1 - e^{-4(k+1)}$ for $k = 1, 2, \dots$ until the inequality is satisfied for some $k = k'$.

- 3) Set $W = 4 + k'$ and generate a new uniform (0,1) random number v .
- 4) Set $u = (p_1 + p_2) \cdot v = (1 - e^{-4}) \cdot v$ and enter the procedure for generating exponentially distributed random variates on the interval (0,4) with $u = (1 - e^{-4}) \cdot v$ as the initial uniform number and generate X .
- 5) Set $Y = W + X$.

In [13] Marsaglia presents a schematic and the necessary constants for the generation of random variates from the exponential distribution. Application of this method results in a very fast generator program. As is evident from the discussion herein the method is also exact; the accuracy of the result is dependent only on the word size of the computer.

4.2 Generation of Gamma Distributed Random Variates

In order to produce random variates having the gamma distribution, the following result is used.

If x_i is a random variable with density

$$f(x_i) = e^{-x_i}, \quad 0 < x_i < \infty \quad \text{for } i = 1, 2, \dots, n \quad (4.2.1)$$

then the random variable

$$Y = \sum_{i=1}^n x_i \quad (4.2.2)$$

has the gamma density

$$g(y) = (1/\Gamma(n))y^{n-1} e^{-y}, \quad 0 < y < \infty \quad (4.2.3)$$

when the x_i 's are stochastically independent. Consequently, the generation of random variates from the gamma distribution involves only a simple

modification of the exponential method.

- 1) Using the method given in 4.1, generate n exponentially distributed random numbers x_1, x_2, \dots, x_n .
- 2) Set $Y = \sum_{i=1}^n x_i$ (if a scale parameter $B \neq 1$ is desired, set $Y = B + \sum_{i=1}^n x_i$).

4.3 Generation of Random Variates

having the Standard Normal Distribution

Using the composition principle of 3.1, the absolute normal density

$$f(t) = \sqrt{2/\pi} e^{-\frac{1}{2}t^2}, t \geq 0 \quad (4.3.1)$$

is represented as a mixture of three densities

$$f(t) = p_1 g_1(t) + p_2 g_2(t) + p_3 g_3(t). \quad (4.3.2)$$

Again, g_1 is a series of rectangles defined for $t \leq 3$, g_2 represents the toothlike regions above g_1 and below f defined for $t \leq 3$, and g_3 is the tail of the density (4.3.1) defined for $t \geq 3$.

Random variates from g_1 are generated by means of the technique given in 3.1. Define

$$\begin{aligned} M &= 3/\Delta \text{ where } \Delta = 10^{-1} \text{ (decimal machine)} \\ \text{or } \Delta &= 16^{-1} \text{ (binary machine), and} \end{aligned} \quad (4.3.3)$$

discrete variables

$$t_i = (i - 1) \cdot \Delta \text{ for } i = 1, 2, \dots, M. \quad (4.3.4)$$

Assign each t_i a probability

$$P_i = \Delta \cdot f(t_{i+1}) = \Delta \cdot \sqrt{2\pi} \cdot e^{-\frac{1}{2}t_{i+1}^2} = .t_{1i} \delta_{xi} \dots \quad (4.3.5)$$

Applying the storage technique described in Chapter 2, store the t_i 's based on the high order 4 octal digits (binary machine) or the high order 3 decimal digits. Note that as in the exponential case $t_{1i} = 0$ for all i , and that the sets of discrete values are overlapped (see the A-table in the listing of GEN3 in the appendix). Now define values

$$Q_0 = 0, \quad Q_r = \beta^{-r} \sum_{i=1}^M \delta_{xi} \quad (4.3.6)$$

and

$$\Pi_0 = 0, \quad \Pi_r = \sum_{j=1}^r \sum_{i=1}^M \delta_{ji}; \quad (4.3.7)$$

also define

$$S_0 = 0, \quad S_k = \beta^k \cdot \sum_{r=1}^k Q_r \quad (4.3.8)$$

and

$$N_i = \Pi_{k-1} - \beta \cdot S_{k-1}, \quad (4.3.9)$$

for a binary machine $r, k = 1, 2, 3, 4$ and $\beta = 8$ for a decimal machine, $r, k = 1, 2, 3$ and $\beta = 10$. The S-values and N-values are also recorded in the program listing of GEN3 (note that the N-values have been adjusted to allow for overlap between sets).

The procedure for the generation of random variates from this truncated portion of g_1 is: (binary case)

- 1) Generate a uniform (0,1) random number $u = .d_1 d_2 \dots$
- 2) If $d_1 d_2 < S_2$, set $Y = A(d_1 d_2 + N_2) + \Delta \cdot (.d_5 d_6 \dots)$;
otherwise go to (3).

- 3) If $d_1 d_2 d_3 < S_3$, set $Y = A\{d_1 d_2 d_3 + N_3\} + \Delta \cdot (.d_5 d_6 \dots)$;
otherwise go to (4).
- 4) If $d_1 d_2 d_3 d_4 < S_4$, set $Y = A\{d_1 d_2 d_3 d_4 + N_4\} + \Delta \cdot (.d_5 d_6 \dots)$;
otherwise generate a random variable from g_2 , g_3 , or the
residual of g_1 .

Random variates from the residual of g_1 and from g_2 are produced by first choosing one of the rectangles of g_1 with appropriate probability and then testing to determine whether the required variate should come from the rectangular residue of g_1 or from the toothlike region of g_2 . As presented in 3.1, this is accomplished as follows:

In M consecutive locations store each t_i once, and for each t_i define a pair of values as in (3.1.12), (3.1.13), and (3.1.14):

$$C(j) = S_4 + R[A(1)], \quad (4.3.10)$$

$$B(k) = C(k) + T[A(k)], \quad (4.3.11)$$

$$C(k+1) = C(k) + R[A(k+1)]. \quad (4.3.12)$$

Note that in the listing of GEN3 this portion of the A-table occupies locations 690 thru 6BF. However, for convenience it is assumed that location 690 correspond to $A(1)$. $R[A(k)]$ and $T[A(k)]$ denote respectively the area of the rectangular remnant and the area of the toothlike region corresponding to the t -value occupying location $A(k)$. For the normal case:

$$R[A(k)] = P[A(k)] - .\delta[1,A(k)]\delta[2,A(k)]\delta[3,A(k)]\delta[4,A(k)] \quad (4.3.13)$$

where $P[A(k)]$ is defined in (4.3.5) as

$$\begin{aligned} P[A(k)] &= \Delta \cdot f(A(k) + \Delta) = \Delta \cdot \sqrt{2/\pi} \cdot e^{-\frac{1}{2}(A(k)+\Delta)^2} \\ &= .\delta[1,A(k)]\delta[2,A(k)]\dots, \end{aligned} \quad (4.3.14)$$

and

$$T[A(k)] = F(A(k) + \Delta) - F(A(k)) = P[A(k)] \quad (4.3.15)$$

where

$$F(x) = \sqrt{2/\pi} \int_0^x e^{-\frac{1}{2}t^2} dt \quad (4.3.16)$$

Once having determined whether a variate with a rectangular density (g_1) or a variate with a toothlike density (g_2) is required, the schemes given in 3.1 and in 3.5 are used to produce the variate as required. This procedure is summarized as follows:

- 1) Let u be a uniform (0,1) random number such that $S_4 \leq u < p_1 + p_2 = F(3)$.
- 2) Test $u < B(k)$ for $k = 1, 2, \dots, M$ until the inequality is satisfied for some $k = k'$.
- 3) If $u < C(k')$, then generate a new uniform (0,1) random number v and put $Y = A(k') + \Delta \cdot v$ (this produces a variate from g_1); otherwise go to (4).
- 4) Generate independent uniform (0,1) random numbers u_1 and u_2 .
- 5) If $\max(u_1, u_2) \leq D(k')^*$, set $Y = A(k') + \Delta \cdot \min(u_1, u_2)$. This produces a variate from g_2 as described in 3.5; otherwise go to (6).
- 6) Set $W = -\frac{1}{2}(\Delta \cdot \min(u_1, u_2) - \Delta) \cdot [2 \cdot A(k') + \Delta \cdot \min(u_1, u_2) + \Delta]$

* The $D(k)$'s are tabulated values of the a_i/b_i ratios given in (3.5.12) and in (3.5.14) corresponding to the t_i stored in $A(k')$.

and test $|u_1 - u_2| < E(k') + (e^W - 1)^{**}$. If yes, set

$Y = A(k') + \Delta + \min(u_1, u_2)$. If no go to step (4) and try again.

A random variate from the density g_3 is produced by the technique presented in 3.6. An absolute normal variable $|Y|$, conditioned by $|Y| > 3$, is required. The procedure is summarized as follows:

- 1) Generate pairs of uniform (0,1) random numbers u_1 and u_2 until the condition $u_1^2 + u_2^2 \leq 1$ is satisfied.
- 2) Form pairs x_1 and x_2 as described in (3.6.1) and (3.6.2)

$$x_1 = u_1 \cdot [\{ 9 - 2\ln(u_1^2 + u_2^2) \} / \{ u_1^2 + u_2^2 \}]^{1/2}$$

$$x_2 = u_2 \cdot [\{ 9 - 2\ln(u_1^2 + u_2^2) \} / \{ u_1^2 + u_2^2 \}]^{1/2}$$

- 3) Test $x_1 > 3$. If yes, set $Y = x_1$. If no, test $x_2 > 3$. If yes, put $Y = x_2$. If no, go to step (1) and repeat the procedure.

In [12] Marsaglia presents a schematic and the required constants for the generation of random variates from the absolute normal distribution. This method results in random variates having the density (4.3.1). In order to obtain random variates from the standard normal distribution.

$$f(t) = (1/\sqrt{2\pi}) e^{-t^2/2}, -\infty < t < \infty, \quad (4.3.17)$$

a random + or - sign must be attached at some point in the procedure. In GEN3 this is done by generating a uniform (0,1) random number u and testing $u < b_2$. If yes, a - sign is affixed to the variate; otherwise a positive variate is returned.

The procedures presented in this chapter derive their speed from the

** This expression is simply an alternative way of testing $b_i |u-v| \leq f(A(k') + \Delta + \min(u_1, u_2))$.

fact that the vast majority of variates generated come from the truncated portion of g_1 . It requires only the generation of a uniform (0,1) random number, a series of compare instructions, and a table lookup. While the generation of variates from g_2 and g_3 may be somewhat time consuming, such variates are required only rarely. Consequently, the average time per generation is quite fast. As is evident these methods are exact. Precision is limited only by the word size of the computer.

CHAPTER 5
RESULTS OF TESTS PERFORMED ON
SEQUENCES OF PSEUDO-RANDOM NUMBERS

5.1 Random Numbers From The
Uniform (0,1) Distribution

The tests performed on sequences of independent uniform (0,1) random numbers were suggested by MacLaren and Marsaglia in [9]. The stringency of these tests is justified in that the procedures for the generation of random variates from other distributions depend heavily upon the method used to obtain uniform (0,1) random numbers.

The tests made were chi-square (χ^2) goodness of fit tests on the distribution of the random numbers, pairs of random numbers, triples of random numbers, and the maximum and minimum of n random numbers. In general a sequence of ℓ variables Y_1, Y_2, \dots, Y_ℓ was calculated from the sequence of random uniform variates. The range of the Y_i was divided into m cells of equal probability, p , and the number of occurrences, O_i , in each cell counted. The χ^2 statistic

$$\chi^2 = \sum_{i=1}^m (O_i - \ell \cdot p)^2 / \ell \cdot p \quad (5.1.1)$$

with $m - 1$ degrees of freedom was calculated and transformed to a standard normal deviate

$$T = (2 \cdot \chi^2)^{1/2} - (2 \cdot (m - 1) - 1)^{1/2} \quad (5.1.2)$$

for each test. This form is valid for $m \geq 31$. The significance level of this normal deviate was then computed as

$$S = \int_0^T (1/\sqrt{2\pi}) e^{-\frac{1}{2}t^2} dt . \quad (5.1.3)$$

In each test a total of 100,000 uniform numbers was generated. For convenience the numbers were generated in 10,000 sets of 10 uniform numbers each. Five separate runs were made and the tests performed as follows:

- 1) Uniformity. The unit interval was divided into 1000 segments. Each segment has an expected value of 100 occurrences. The χ^2 statistic was computed as

$$\chi^2 = \sum_{i=1}^{1000} (O_i - 100)^2 / 100 \quad (5.1.4)$$

where O_i is the number of observed occurrences in the i^{th} segment.

- 2) Pairs. Successive pairs of uniform numbers were taken as the coordinates of a point in the unit square. The unit square was partitioned into 100 cells, each with an expected value of 500 occurrences. The χ^2 statistic was computed as

$$\chi^2 = \sum_{i=1}^{10} \sum_{j=1}^{10} (O_{ij} - 500)^2 / 500 \quad (5.1.5)$$

where O_{ij} is the number of observed occurrences in the ij^{th} cell.

- 3) Triples. Successive triples of uniform numbers were taken as the coordinates of a point in the unit cube (every tenth number was skipped). The unit cube was partitioned into 1000 cells with an expected

value of 30 occurrences in each cell. The χ^2 statistic was computed as

$$\chi^2 = \sum_{i=1}^{10} \sum_{j=1}^{10} \sum_{k=1}^{10} (O_{ijk} - 30)^2 / 30 \quad (5.1.6)$$

where O_{ijk} is the number of observed occurrences in the ijk^{th} cell.

4) Maximum of n random values. If u_1, u_2, \dots, u_n are independent uniform (0,1) deviates, then

$$W = \max(u_1, u_2, \dots, u_n) \quad (5.1.7)$$

should have the distribution

$$P(W \leq a) = F(a) = a^n \quad \text{for } 0 < a < 1, \text{ and} \quad (5.1.8)$$

$$F(W) = [\max(u_1, u_2, \dots, u_n)]^n \quad (5.1.9)$$

should be uniformly distributed over the interval (0,1). A total of 10,000 W 's were generated for each n , and $F(W)$ was tested for uniformity using the χ^2 goodness of fit test for 100 equal subintervals. The χ^2 statistic was computed as

$$\chi^2 = \sum_{i=1}^{100} (O_i - 100)^2 / 100 \quad (5.1.10)$$

where O_i is the number of observed occurrences in the i^{th} subinterval.

In this test n takes the values 2 and 5.

5) Minimum of n random values. This test is the same as that for the maximum of n except

$$W = \min(u_1, u_2, \dots, u_n) \quad \text{and} \quad (5.1.11)$$

and

$$F(W) = 1 - (1 - W)^n \quad (5.1.12)$$

which should be uniformly distributed over the unit interval. The 10,000 W 's were segmented into 100 equal subintervals and the χ^2 test computed as in (5.1.10). In this test n takes the values 3 and 10.

The results of these tests are reported in Table 5.1.1. The results compare favorably with those presented in [9] and justify the use of this particular uniform (0,1) random number generator.

TABLE 5.1.1

RUN #1

TEST	CHI-SQUARE	D.F.	SIGNIFICANCE
UNIFORMITY	975.400	999	0.30186688
PAIRS	89.6720	99	0.25987495
TRIPLES	1034.7333	999	0.78918171
MAX OF 2	97.1800	99	0.46241028
MAX OF 5	80.7800	99	0.09257838
MIN OF 3	109.5200	99	0.77766504
MIN OF 10	75.0400	99	0.03713434

RUN #2

TEST	CHI-SQUARE	D.F.	SIGNIFICANCE
UNIFORMITY	1047.1600	999	0.85902187
PAIRS	92.8440	99	0.34129536
TRIPLES	1022.4667	999	0.70302930
MAX OF 2	104.4800	99	0.66267689
MAX OF 5	96.9800	99	0.45671365
MIN OF 3	102.1000	99	0.60032472
MIN OF 10	97.7400	99	0.47836695

TABLE 5.1.1 (continued)

RUN #3

TEST	CHI-SQUARE	D.F.	SIGNIFICANCE
UNIFORMITY	1006.3400	999	
PAIRS	93.8720	99	0.56949803
TRIPLES	1052.0000	99	0.36930366
MAX OF 2	93.5400	99	0.88131258
MAX OF 5	104.5000	99	0.36019046
MIN OF 3	102.7400	99	0.66318213
MIN OF 10	107.0600	99	0.61749364
			0.72480337

RUN #4

TEST	CHI-SQUARE	D.F.	SIGNIFICANCE
UNIFORMITY	970.2600	999	
PAIRS	103.9040	99	0.26223433
TRIPLES	1084.4667	99	0.64798213
MAX OF 2	83.2400	99	0.97021768
MAX OF 5	96.4200	99	0.12861638
MIN OF 3	76.7000	99	0.44078290
MIN OF 10	116.5000	99	0.04945168
			0.89040197

TABLE 5.1.1 (continued)

RUN #5

TEST	CHI-SQUARE	D.F.	SIGNIFICANCE
UNIFORMITY	1052.4000	999	0.88303445
PAIRS	82.5080	99	0.11705985
TRIPLES	1013.5333	999	0.63124447
MAX OF 2	92.6200	99	0.33528034
MAX OF 5	98.8400	99	0.50965471
MIN OF 3	140.6600	99	0.99687357
MIN OF 10	136.3000	99	0.99333696

5.2 Random Numbers Having A Discrete Distribution

Sequences of random numbers from the binomial density

$$f(x) = \binom{N}{x} \cdot p^x \cdot (1-p)^{N-x}, \quad 0 < p < 1, \quad x=0,1,\dots,N \quad (5.2.1)$$

were generated for a variety of parameter combinations. The range of the variable was divided into k intervals of probability q_i , such that the expected value of each interval, $m \cdot q_i$, was at least 10. The χ^2 statistic was computed as

$$\chi^2 = \sum_{i=1}^k (O_i - m \cdot q_i)^2 / m \cdot q_i \quad (5.2.2)$$

where O_i is the number of observed occurrences in the i^{th} interval, and m is the size of the sample generated. The significance level was computed as

$$S = (1/\Gamma(v/2)) \cdot \int_0^{\chi^2/2} e^{-t} \cdot t^{(v/2 - 1)} dt \quad (5.2.3)$$

where v = degrees of freedom, in this case $v = k-1$. The results of these tests are recorded in table 5.2.1. A sample size of 1000 random numbers were generated for each test.

Random numbers from the Poisson density

$$f(x) = \lambda^x e^{-\lambda} / x!, \quad 0 < \lambda < \infty, \quad x = 0,1,\dots,\infty \quad (5.2.4)$$

and from the negative binomial density

$$f(x) = \binom{r+x-1}{x} \cdot p^r \cdot (1-p)^x, \quad r \geq 0, \quad 0 < p < 1, \quad (5.2.5)$$

$$x = 0,1,2,\dots,\infty;$$

were tested in the same manner as random numbers from the binomial density (5.2.1). The results of these tests are recorded in table 5.2.2 and 5.2.3 respectively.

TABLE 5.2.1

RESULTS OF CHI-SQUARE GOODNESS OF FIT TESTS
ON RANDOM NUMBERS FROM THE BINOMIAL DENSITY

P	N	D.F.	CHI-SQUARE	SIGNIFICANCE
0.100	10	4	6.03494	0.8034458581
0.100	20	6	6.44102	0.6243600635
0.100	30	7	2.52631	0.0748969943
0.100	40	9	6.33123	0.2936355148
0.100	50	9	12.92191	0.8338284965
0.200	10	5	7.56431	0.8180614798
0.200	20	8	16.34847	0.9623445588
0.200	30	10	15.13733	0.8728654417
0.200	40	11	9.00770	0.3788180126
0.200	50	13	12.00444	0.4727203102
0.300	10	7	7.97911	0.6655606411
0.300	20	9	10.98540	0.7232888983
0.300	30	11	17.78713	0.9133467551
0.300	40	13	12.08372	0.4792085395
0.300	50	15	21.66071	0.8829858823
0.400	10	7	7.29406	0.6010818180
0.400	20	10	12.18194	0.7269356404
0.400	30	12	20.64248	0.9441325898
0.400	40	14	10.05124	0.2415723570
0.400	50	16	12.35961	0.2811236434
0.500	10	8	6.86156	0.4483592431
0.500	20	10	12.57912	0.7518355135
0.500	30	12	10.41700	0.4205727006
0.500	40	14	23.28293	0.9441810635
0.500	50	16	9.42171	0.1049637735
0.600	10	7	6.92477	0.5632454466
0.600	20	10	18.27415	0.9494885520
0.600	30	12	10.17474	0.3993657394
0.600	40	14	11.25580	0.3341658613
0.600	50	16	10.03645	0.1352818684

TABLE 5.2.1 (continued)

P	N	D.F.	CHI-SQUARE	SIGNIFICANCE
0.700	10	7	3.47126	0.1617396643
0.700	20	9	10.70221	0.7033259190
0.700	30	11	12.02199	0.6380075901
0.700	40	13	13.88283	0.6178529373
0.700	50	15	18.60988	0.7680404078
0.800	10	5	4.53985	0.5254310881
0.800	20	8	16.74590	0.9671335640
0.800	30	10	10.97150	0.6402573630
0.800	40	11	9.02074	0.3800224541
0.800	50	13	15.05620	0.6961423824
0.900	10	4	3.31137	0.4928666476
0.900	20	6	4.40204	0.3775593406
0.900	30	7	6.78871	0.5487935778
0.900	40	9	6.74965	0.3368342656
0.900	50	9	5.04256	0.1694175512

RESULTS OF CHI-SQUARE GOODNESS OF FIT TESTS ON
RANDOM NUMBERS FROM THE POISSON DENSITY.

LAM	D.F.	CHI-SQUARE	SIGNIFICANCE
5.00	10	13.57652	0.8068012578
10.00	15	8.51890	0.0986988385
15.00	17	16.33853	0.5000245564
20.00	20	18.35403	0.4359007928
25.00	21	28.70494	0.8787325544
30.00	23	24.94425	0.6468336304
35.00	25	26.91360	0.6397852204
40.00	27	43.19749	0.9750174127
45.00	28	22.55633	0.2450911949
50.00	30	23.26187	0.1956567510
55.00	30	42.01434	0.9286283430
60.00	31	34.81330	0.7087406320
65.00	33	29.49011	0.3573849816
70.00	34	29.73472	0.3231975827
75.00	35	26.18810	0.1410044103
80.00	35	26.80832	0.1619395204
85.00	36	21.61020	0.0278060381
90.00	37	28.29852	0.1527661009
95.00	38	30.51639	0.1992224404
100.00	39	37.10850	0.4435994584

TABLE 5.2.2

TABLE 5.2.3
 RESULTS OF CHI-SQUARE GOODNESS OF FIT TESTS ON RANDOM
 NUMBERS FROM THE NEGATIVE BINOMIAL DENSITY.

P	R	D.F.	CHI-SQUARE	SIGNIFICANCE
0.100	2	41	36.08627	0.3115372475
0.100	4	57	51.64328	0.3244292796
0.100	6	67	77.44254	0.8201664431
0.100	8	72	54.21390	0.0583581269
0.100	10	76	58.52535	0.0683503608
0.200	2	23	20.00081	0.3581367902
0.200	4	33	27.73007	0.2730820886
0.200	6	40	43.00986	0.6563763897
0.200	8	45	42.83166	0.4357673160
0.200	10	49	38.69528	0.1455417220
0.300	2	15	11.10495	0.2548804445
0.300	4	22	25.94653	0.7460151638
0.300	6	28	17.46899	0.0612580512
0.300	8	31	25.10221	0.2368731801
0.300	10	34	18.65952	0.0151991805
0.400	2	11	12.32864	0.6605459138
0.400	4	16	7.20666	0.0309309696
0.400	6	20	20.18198	0.5534007984
0.400	8	23	23.12594	0.5466209213
0.400	10	25	22.04679	0.3669595654
0.500	2	8	4.96920	0.2391366628
0.500	4	12	13.04976	0.6345655810
0.500	6	15	16.03409	0.6202105433
0.500	8	17	25.16863	0.9090016643
0.500	10	19	24.55579	0.8243107298
0.600	2	6	13.11971	0.9588261108
0.600	4	9	14.42509	0.8920130256
0.600	6	11	13.29194	0.7253265963
0.600	8	13	17.84469	0.8364937782

TABLE 5.2.3 (continued)

P	R	D.F.	CHI-SQUARE	SIGNIFICANCE
0.600	10	14	13.95647	0.5470410000
0.700	2	5	6.42472	0.7329438096
0.700	4	7	11.10162	0.8657526004
0.700	6	8	9.98931	0.7342226732
0.700	8	10	4.82568	0.0974859792
0.700	10	11	14.96953	0.8161040822
0.800	2	3	1.03482	0.2071720727
0.800	4	5	2.88184	0.2818034287
0.800	6	6	1.48994	0.0398398009
0.800	8	7	4.59512	0.2907654804
0.800	10	8	2.13612	0.0234479888
0.900	2	2	0.67942	0.2880246110
0.900	4	3	0.23420	0.0281120703
0.900	6	3	4.99992	0.8281970919
0.900	8	4	2.59324	0.3719796012
0.900	10	4	2.34658	0.3276973945

5.3 Random Numbers From Continuous
Distribution Functions

Sequences of random numbers from the exponential

$$f(t) = e^{-t}, \quad 0 < t < \infty \quad (5.3.10)$$

and normal

$$f(x) = (1/\sqrt{2\pi}) e^{-x^2/2}, \quad -\infty < x < \infty \quad (5.3.2)$$

densities were generated, and segmented into 20 intervals of approximately equal probability, $p_i = 0.05$. Samples of $m = 1000$ random numbers were generated and the χ^2 statistic computed as

$$\chi^2 = \sum_{i=1}^{20} (O_i - m \cdot p_i)^2 / m \cdot p_i \quad (5.3.3)$$

where O_i is the number of observed occurrences in the i^{th} interval. The significance level was calculated as given in (5.2.3) with $v = 19$. The results of these tests are given in tables 5.3.1 and 5.3.2.

Sequences of random numbers from the gamma density

$$f(t) = (1/\Gamma(n)) t^{n-1} e^{-t}, \quad 0 < t < \infty, \quad (5.3.4)$$

were tested in a somewhat different manner. If t is a variable with density (5.3.4) then

$$y = 2 \cdot (t)^{1/2} - (4 \cdot n - 1)^{1/2} \quad (5.3.5)$$

should have an approximate standard normal density when n is large.

Sequences of m gamma distributed random numbers were generated and using the transformation (5.3.5), transformed to m approximately standard normal variables. As before, the normal distribution was partitioned into 20

RESULTS OF CHI-SQUARE GOODNESS OF FIT TESTS ON RANDOM
NUMBERS FROM THE EXPONENTIAL DENSITY, $F(X) = \text{EXP}(-X)$.

D.F.	CHI-SQUARE	SIGNIFICANCE
19	20.76000	0.6497876
19	11.24000	0.0844989
19	17.64000	0.4534074
19	22.68000	0.7482694
19	13.68000	0.1979842
19	18.52000	0.5120074
19	11.48000	0.0933650
19	17.36000	0.4345072
19	24.24000	0.8128932
19	37.24000	0.9925960

TABLE 5.3.1

RESULTS OF CHI-SQUARE GOODNESS OF FIT TESTS ON RANDOM
NUMBERS FROM THE NORMAL DENSITY

D.F.	CHI-SQUARE	SIGNIFICANCE
19	19.59699	0.5808149783
19	25.42150	0.8528693863
19	18.29836	0.4974018037
19	6.71247	0.0044039802
19	14.75770	0.2621316817
19	19.91207	0.6001134403
19	20.76446	0.6500395713
19	23.50146	0.7840244000
19	11.96570	0.1129139927
19	18.02468	0.4792096579

TABLE 5.3.2

intervals of approximately equal probability $p_i = 0.05$, and the χ^2 statistic computed as before. The results of these tests are recorded in table 5.3.3.

RESULTS OF CHI-SQUARE GOODNESS OF FIT TESTS ON TRANSFORMED
GAMMA RANDOM NUMBERS

SAMPLE SIZE	N	D.F.	CHI-SQUARE	SIGNIFICANCE
100	50	19	12.30156	0.1276659399
100	100	19	14.27597	0.2326276646
100	200	19	13.79290	0.2043727108
200	50	19	13.99520	0.2160288470
200	100	19	20.25710	0.6207407158
200	200	19	23.61244	0.7885588803
500	50	19	23.13291	0.7684615439
500	100	19	21.91446	0.7114659110
500	200	19	32.02104	0.9689153752

TABLE 5.3.3

5.4 Discussion of Results.

It should be emphasized that no test or series of tests can assure the suitability of a given sequence of random numbers for a particular problem. When possible the sequence should be tested on a similar problem with known solution. However, the results presented here establish the fundamental reliability of the generation procedures presented in this paper. The significance levels obtained should follow the uniform distribution on (0,1) if the corresponding null hypothesis tested are valid. The above results are in agreement with this condition.

CHAPTER 6
DESCRIPTION AND USE
OF GENERATOR PROGRAMS

6.1 GENO

GENO supplies the user with a very fast procedure for the generation of independent uniform (0,1) random numbers with density

$$f(u) = 1, \quad 0 < u < 1. \quad (6.1.1)$$

The method used was suggested by Dr. Rolf Bargmann [1] and is a form of the multiplicative congruential procedure,

$$u_{n+1} = a \cdot u_n \bmod 2^{32} \quad (6.1.2)$$

where $a = 5^{13}$. This procedure requires that u_0 , the starting value, be an odd positive integer.

GENO is written as a subroutine with entry point URAN. It is called as follows:

$$U = \text{URAN}(IOD), \quad (6.1.3)$$

where "IOD" is the starting value and must be an odd positive integer. The desired random number U in (6.1.3), is returned to the calling program in single precision real mode.

GENO requires 20 words (80 bytes) of storage space and an average time per generation of 25 usec.

6.2 GEN1

GEN1 provides the user with a fast procedure for the generation of exponentially distributed random numbers

$$f(x) = e^{-x}, \quad 0 < x < \infty. \quad (6.2.1)$$

The method used is that described in 4.1, and presented by Marsaglia in [13].

GEN1 is written as a subroutine with entry point RANEXP. It is called from a FORTRAN main program as follows:

$$X = \text{RANEXP}(IOD). \quad (6.2.2)$$

The parameter "IOD" primes the scheme for the generation of uniform (0,1) random numbers, "IOD" must be an odd positive integer to insure the proper generation of u. The required random number x in (6.2.2) is returned to the calling program in single precision real mode.

The average time per generation is approximately 70 usec. Memory requirement for GEN1 is 307 words (1228 bytes).

6.3 GEN2

GEN2 supplies the user with a fast procedure for the generation of gamma ($\alpha = n$, $\beta = 1$)

$$g(y) = 1/\Gamma(n) \cdot y^{n-1} \cdot e^{-y}, \quad 0 < y < \infty \quad (6.3.1)$$

distributed random numbers. The technique used is described in 4.2.

GEN2 is written as a subroutine with entry point RANGAM, and is called from a FORTRAN main program as follows:

$$Y = \text{RANGAM}(IOD, N). \quad (6.3.2)$$

As before, "IOD" must be an odd positive integer and $n = "N"$ in (6.3.1).

The required variate is returned in single precision real mode.

GEN2 requires 313 words of memory space and approximately $65 \cdot N$ usec per generation for $N > 1$.

6.4 GEN3

GEN3 supplies the user with a fast procedure for the generation of random numbers from the standard normal density

$$f(x) = (1/\sqrt{2\pi}) e^{-\frac{1}{2}x^2}, -\infty < x < \infty. \quad (6.4.1)$$

The procedure is described in 4.3 and by Marsaglia in [12].

GEN3 is written as a subroutine and is called from a FORTRAN main program as follows:

$$X = \text{RANORM}(IOD). \quad (6.4.2)$$

The parameter "IOD" must be an odd positive integer, and the required variate X in (6.4.2) is returned to the calling program in single precision real mode.

Memory requirement for GEN3 is 442 words. The average time per generation is approximately 65 usec.

6.5 GEN4

GEN4 provides a fast procedure for the generation of random numbers from the binomial distribution

$$f(x) = \binom{n}{x} p^x (1-p)^{n-x}, 0 < p < 1, x=0,1,\dots,n. \quad (6.5.1)$$

GEN4 is written as a subroutine with two entry points, BSETUP and RANBI. The user calls BSETUP with the parameters "P" and "N" e.g.

CALL BSETUP(P,N). (6.5.2)

BSETUP calculates point probabilities using the recursion relation

$$f(0) = (1 - P)^N \text{ and} \quad (6.4.3)$$

$$f(x+1) = f(x) + P \cdot (N-x)/((x+1) \cdot (1-P)).$$

The four sets of discrete values are then stored using the procedure described in Chapter 2. BSETUP need be called only once for a given "P" and "N", but must be called at least once prior to calling RANBI.

RANBI addresses the very fast scheme that generates the required variates. It is called as follows:

X = RANBI(IOD) (6.4.4)

where "IOD" must be an odd positive integer, and X is returned in single precision real mode.

In its present form GEN4 requires the "P" be in the real mode, $0 < P < 1$, and that "N" be in the integer mode and not exceed 255. The program requires 948 words of memory and an average time per generation of approximately 35-40 usec. Figure 2.1 describes the flowchart of GEN4 except for the section that computes point probabilities.

If the user desires random numbers returned in the integer mode, he need only delete

STMT

199	ST	0, RES
200	MVI	RES, X'46'
201	AE	0, RES

and change the designation RANBI to IANBI in his calling statement and in statements

```

170          ENTRY RANBI
171          USING RANBI, 15
173          DC    CL7' ..RANBI'
174RANBI     STM   1,3, 24(13)

```

Storage requirements may also be altered simply by changing

```

STMT
146          C      7, = F'2000'
222          DS    2000XL1.

```

6.5 GENS

GENS supplies a fast routine for the generation of random numbers from the poisson density

$$f(x) = \lambda^x e^{-\lambda} / x! , 0 < \lambda < \infty , x = 0,1,\dots \quad (6.5.1)$$

Since the domain of x is infinite, it is necessary to truncate the distribution. This is done at both low and high point values whenever a point probability is less than 16^{-4} .

GENS is written as a subroutine with two entry points, PSETUP and RANPOI. The user calls PSETUP with the parameter λ , e.g.

```
CALL PSETUP(ALAM). \quad (6.5.2)
```

PSETUP calculates point probabilities using the recursion relation

$$f(0) = e^{-\lambda}$$

$$f(x+1) = f(x) + \lambda / (x+1), \quad (6.5.3)$$

and stores the four sets of discrete values by means of the technique described in Chapter 2. PSETUP need be called only once for a given λ , but must be called at least once before the initial call to RANPOI.

The entry point RANPOI addresses the very fast scheme that performs the actual generation of the required variate. It is called as follows:

$$X = \text{RANPOI}(\text{IOD}), \quad (6.5.4)$$

where "IOD" must be an odd positive integer and the desired random number X is returned in single precision real mode.

GEN5 requires 972 words of memory space and an average time per generation of approximately 35-40 usec. The flow of GEN5 follows that presented in figure 2.1 except for the section that computes the point probabilities.

If the user desires his result returned in the integer mode, he need only delete

STMT		
190	ST	0,RES
191	MVI	RES, X'46'
192	AE	0,RES

and change the designation RANPOI to IANPOI in his calling statement and in statements

STMT		
161	ENTRY	RANPOI
162	USING	RANPOI
164	DC	CL7' RANPOI'
165 RANPOI	STM	1,3,24(13)

Storage requirement may also be altered by changing

STMT

137	C	7, = F'2000'
211 A	DS	2000XL1

6.6 GEN6

GEN6 provides a fast routine for generating random numbers from the negative binomial density

$$f(x) = \binom{x+r-1}{x} p^r (1-p)^x, \quad 0 < p < 1, \quad x = 0, 1, \dots, \quad (6.6.1)$$

and $r \geq 0$.

The domain of x is infinite, consequently, the distribution is truncated at both low and high point values whenever a point probability is less than 16^{-4} .

GEN6 is written as a subroutine with two entry points, NBSETU and RANEBI. NBSETU receives the parameters, p and r , and calculates point probabilities using the recursion relation

$$f(0) = p^r \quad (6.6.2)$$

$$f(x+1) = f(x) \cdot (1-p) \cdot (x+r)/(x+1).$$

Having computed the point probabilities, NBSETU then stores the four sets of discrete values using the procedure described in Chapter 2. NBSETU is called as follows:

CALL NBSETU(P,IR), (6.6.3)

where "IR" must be in the integer mode and $0 < "P" < 1$. NBSETU need be

called only once for a given "P" and "IR", but must be called at least once prior to the initial call to RANEBI.

RANEBI addresses the very fast scheme for the generation of the required variates, and is called as follows:

$$X = \text{RANEBI}(IOD), \quad (6.6.4)$$

where "IOD", the primer, must be an odd positive integer, and X is returned in single precision real mode.

GEN6 requires 944 words of storage space and an average time per generation in the 35-40 usec range. The flow-chart of GEN6 follows that of figure 2.1 except for the section of GEN6 that calculates the point probabilities.

6.7 GEN7

GEN7 provides a fast procedure for generating discrete random variables with any specified probability distribution denoted by a vector $P = \{p_1, p_2, \dots\}$. It is written as a subroutine with two entry points, DSETUP and RANDIS. DSETUP is called with 3 parameters, "P", "IX", "N"; where "P" is a vector of probabilities (real mode), "IX" is a vector of corresponding discrete values (integer mode), and "N" is the number of elements in vectors, e.g.

$$\text{CALL DSETUP}(P, IX, N). \quad (6.7.1)$$

The discrete values of the "IX" vector are stored in four sets based on the high order 4 digits of the elements of the "P" vector. This procedure is given in Chapter 2 and illustrated in figure 2.1. DSETUP need be called

only once for a particular set of parameters, however, it must be called at least once prior to the initial call to RANDIS.

RANDIS addresses the very fast scheme that performs the actual generation of the required variates. It is called as follows:

$$X = \text{RANDIS}(IOD), \quad (6.7.2)$$

where "IOD" must be an odd positive integer and "X", the generated variate is returned in single precision real mode.

In its present form GEN7 requires that the elements of the "IX" vector be in the integer mode and that no value exceed 255. Also, the parameter "N", the number of elements in the vectors, may not exceed 256. If the user wishes his discrete variables, members of the "IX" vector, to be in the real mode or to be in the integer mode but with values greater than 255, he need only make the changes indicated by table 6.1. If he desires to have his result returned in the integer mode he need only delete the following

STMT		
149	ST	0,RES
150	MVI	RES, X'46'
151	AE	0,RES,

and change the designation RANDIS to IANDIS in his calling statement and in the following

STMT		
120	ENTRY	RANDIS
121	USING	RANDIS, 15
123	DC	CL7' RANDIS'
124 RANDIS	STM	1,3,24(13).

Storage requirement may also be altered simply by changing

STMT

97 C 7, = F'2000'

168 A DS 2000XL1

In its present form GEN7 requires 904 words of memory space and
35-40 usec per generation.

STMT	INTEGER > 255			REAL MODE		
87	L	6, = F'1'	LA	6,4	LA	6,4
	C	7, = F'2000'	C	7, = F'SPAC'	C	7, = F'SPAC'
97			L	2, 0(3,11)	LE	2, 0(3,11)
99	L	2, 0(3,11)	ST	2, A(12)	STE	2, A(12)
	STC	2, A(12)	C	0, = F'16'	C	0, = F'16'
101	CR	0,4	EX SLL	2, 2	EX SLL	2, 2
106	IC	0, A(2)	SPAC XL4	A DS	SPAC XL4	A DS
148 EX	DS	2000 XL1				
168 A						
				ALSO INSERT BETWEEN STMT 95 AND STMT 96	ALSO INSERT BETWEEN STMT 95 AND STMT 96	
				SLL 8, 2	SLL 8, 2	
					AND DELETE STMT 149	
					THROUGH STMT 151.	

TABLE 6.1

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APPENDIX A

Assembler Listing of GEN0

NOT REPRODUCIBLE

LRC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
000000				1 GENI START O	
				2 *****	
				3 * SOURCE--L.E.CANNON, UGA, DEPT. OF STAT., APR 1970	
				4 * PURPOSE--TO SUPPLY THE USFA WITH A FAST PROCEDURE FOR GENERATING JPN-	
				5 * INFORMATION, RANDOM NUMBERS.	
				6 * USAGE---X=URAN(LIND). "LIND" MUST BE AN AND POSITIVE INTEGER. THE RE-	
				7 * QUIRED VARIATE IS RETURNED IN SINGLE PRECISION REAL MODE.	
				8 * EXAMPLE:X=URAN(.56365) *X WILL BE UNIFORM ON [0,1].	
				9 * METHOD--THE METHOD IS THE MULTPLICATIVE CONGRUENTIAL PROCEDURE	
				10 * X=(171)*X+11*A	
				11 * WHERE A=54413. AS SUGGESTED BY DR. R. F. KAGGIANI.	
				12 *****	
				13 ENTRY URAN	
				14 USING URAN,15	
				15 DC XL1107	
				16 DC CL7, URAN*	
				17 UPAN STM# 1•3,24(13)	
				18 SER CLEAR FPRO FOR RESULT.	
				19 ARANCH BR TEMP.	
				20 R04C M 2,RPAT	
				21 U0050 ST 3,LA	
				22 S0008 ST 3,S	
				23 ST 3,RES	
				24 NVI RES,X'40'	
				25 AE 0,RES	
				26 LM 1•3,24(13)	
				27 SCR 15,14	
				28 TFMP L 3,LA	
				29 F028 00050 NYC BRANCH,TEMP REPLACE INSTR. AT BRANCH. AT TEMP.	
				30 00034 0203 F036 F028 00050 L 3,(1,1)	
				31 00034 5830 1000 00000 L 3,(0,3)	
				32 00035 5831 1000 00012 B BRANCH+4	
				33 RES DS 1F	
				34 RPAT DC XL4•48C27395,	
				35 LA OS IF	
				36 END	

APPENDIX B

Assembler Listing of GENI

PAGE 1
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LCC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
00000000					1 GEN1 START 0
00000000					2 *****
00000000					3 * SOURCE---L. E. CANNON, IIGA, DFPT, OF STAN. * APR 1970
00000000					4 * PURPOSE---TO SUPPLY THE USER WITH A FAST PROCEDURE FOR GENERATING EX-*
00000000					5 * PENSATIONAL DISTRIBUTED RANDOM NUMBERS.
00000000					6 * USAGE---MARSAGLIA. IT MUST BE AN ODD INTEGER. IT PRIMES THE GEN-*
00000000					7 * FRATING SCIFI.
00000000					8 * EXAMPLE---X=RANEXP(1000000000).
00000000					9 * METHOD---THE METHOD IS THAT OF MARSAGLIA AS DESCRIBED IN "COMMUNI-
00000000					10 * CATIONS OF ACM", VOL. 7, NO. 5, MAY 1964. THE PROGRAM DOCUMENT-
00000000					11 * TATION WILL FOLLOW THE NOTATION USED BY MARSAGLIA.
00000000					12 *****
00000000					13 ENTRY RANEXP PROLOGUE, STANDARD FOR 1607/65.
00000000					14 USING RANEXP, IS
00000000					15 OR X107,
00000000					16 CC CL7, RANEXP
00000000					17 QANEXP ST4 100,241131
00000000					18 BRANCH B TFMPL EXECUTED ONLY ON THE FIRST CALL.
00000000					19 SEP 0,0 RESULT SUMMED IN FPPG.
00000000					20 H 4,RPAT OBTAIN NEW U.
00000000					21 ST S,LA STORE FOR NEXT GENERATION.
00000000					22 SR 0,6
00000000					23 MI 4,4
00000000					24 SLDL 4,6
00000000					25 C 4,0,5,37, D102, C100400 WITH 171(BASE 171=4,518 BASE 8).
00000000					26 AC 11,NA 6,0,24137 BRANCH NA.
00000000					27 IC 6,0,240614) OBTAIN DISCRETE RANDW VARIABLE FROM D.
00000000					28 SLL 5,5 RS=USDOAD(0,0)
00000000					29 BC 15,EX BRANCH 67.
00000000					30 NA SLDL 4,3 RA=010201.
00000000					31 DD4C4 11 C 4,0,F,45A, 4,518(BASE 10)=717(BASE 9).
00000000					32 0004E 37 AC 11,NA D1020304>4,518 BRANCH NR.
00000000					33 0N2F4 33 BC 6,0,1441) OBTAIN DISCRETE RANDW VARIABLE FROM D.
00000000					34 SLL 5,3 RS=DSOVO(0,0)
00000000					35 AC 15,EX BRANCH 61.
00000000					36 SLDL 4,3 RA=010201.
00000000					37 C 4,0,F,38639, 01020304>3861 BRANCH NC.
00000000					38 BC 11,NC OBTAIN PROPER ADDRESS FOR D TABLE.
00000000					39 S 4,0,F,38639, 01020304>3861 BRANCH NC.
00000000					40 IC 6,0,141, OBTAIN DISCRETE VARIABLE FROM D.
00000000					41 LR 4,4 RA=010201.
00000000					42 SRLN 4,1,16 RS=DISCRETE VARIABLE + HI ORDER 16 BITS OF
00000000					43 ST 5,RES RS=DISCRETE VARIABLE + HI ORDER 16 BITS OF
00000000					44 HWI RES, X,41, RS=DISCRETE VARIABLE + HI ORDER 8
00000000					45 AF 0,RES Bits ZERO OUT FOR CONVERSION TO FLOAT-
00000000					46 LM 1,9,241131 ING PRINT. FPBD CONTAINS RESULT + HI.
00000000					47 ACR 15,14, AR=010201) RESTORE REGISTERS.
00000000					48 L 5,LA AR=010201) BRANCH TO MAIN LINE.
00000000					49 RA 2,4 RESTORE U TO RS.
00000000					50 SR 3,3 UP=PL0) BRANCH N.
00000000					51 CL 5,P(13) UP=PL0) BRANCH N.
00000000					52 DC 11,NO UP=PL0) BRANCH N.
00000000					53 SR 3,2 UP=PL0) BRANCH N.
00000000					54 AR 3,2 UP=PL0) BRANCH N.
00000000					55 CL 5,R(13) UP=PL0) BRANCH N.

LPC	OBJECT CODE	ADDR1 A0NP2	STMT	SOURCE STATEMENT	
000004	4740 F046	00000E	96	AC 11, NF	
000004	5551 F234	0029C	97	CL 5,C1)	
000004	47B1 F048	00000D	58	BC 11, NF	
000004	5C49 F130	0013R	59	H 4, NPAT	
000004	5051 F12A	00130	61	ST 5,LA	
000004	1B4A	61 C0	SR 4,4	C(1B1-CUCB1) BRANCH NF.	
000004	A073 DCN2	00002	67	SPI 3,2	
000004	4161 F04C	00374	63	IC 4,0,11	
000004	47F3 F05C	00164	64	AC 15,FL	
000004	5A51 F12A	00130	65	L S,LA	
000004	D731 F004	F0AE 0000C	ncn66	PVC BRANCH,TEMP	
000004	1A57 LG07	00000	66	REPLACE BRANCH BY TMP.	
000004	5A53 SU03	00000	67	L 5,J,1,1	
000004	5C51 S12A	00130	68	L 5,C1,4	
000004	5C54 F004	00110	69	ST 5,LA	
000004	47F1 F004	00110	70	H APANCH4,	
000004	5C49 F110	00130	71	H 4, NPAT	
1U0004	1A52	72	LR 8,2	OBTAI NF U.	
1U0004	1A52	73	AR 9,2	USFO FOR INDEXING.	
1U0004	4554 F134	0013C	74	CL 5,0,0	
1U0004	47A0 F0CE	00006	75	AC 11, -6	
1U0004	5C49 F110	00138	76	H 4, NPAT	
1U0004	1A52	77	SP 8,2	DATAIN NEW U.	
1U0004	1A52	78	LR 5,5	R9=4.	
1U0004	1A52	79	LR 9,2		
1U0004	5C49 F130	0013A	80	SR 8,2	
1U0004	1A54	81	NR 9,2		
1U0004	47F2 47F1 F0FA	0013A	82	CL 9,4	
1U0004	1A55 F1765	0000A	83	PC 11,0,6	
1U0004	8A89 F0FA	0000C	84	LR 6,4	
1U0004	5C53 F129	00130	85	8,0,0	
1U0004	1A56	86	ST 5,LA	LOAD SMALLEST U IN R4.	
1U0004	5C52 47F1 F040	0000A	87	LA 5,4	
1U0004	1A52	88	AC 15,0,0	CONTINUE(R9 HAS BOTH INCREMENT & DECREMENT).	
1U0004	5551 F156	0013C	89	AR 3,2	R3=10.
1U0004	4741 F0FE	00106	90	C1 5,P11)	
1U0004	5C51 F12C	00134	91	ST 11,ND	
1U0004	5261 F12C	00136	92	ST 3,0,4,6,	
1U0004	5C61 F12C	00134	93	MVI J,0,4,6,	
1U0004	5C61 F130	00138	94	AF 0,0,0	
1U0004	4551 F12A	00130	95	H 4, NPAT	
1U0004	5C51 F12A	0035C	96	ST 5,LA	
1U0004	47F1 F014	0001C	97	H 4, P61	
1U0004	5C51 F12C	98	AC 15,NI	STORE IP INTL J.	
1U0004	1A52	99 RES	DS 1,F	CONVERT IP TO FLOATING POINT.	
1U0004	10114	101 J	DS 1,F	NORMALIZE AND LOADED INTO FPR9 +W+.	
200114	4PC27395	102 NPAT	XL 6,4,4C27395,	BASE 10 0,0	
200114	10011C 10011C00	103 Q	X'0000000000000000	BASE 10 1,0	
200114	00000C00	104	X'0000000000000000	BASE 10 1,0	
200114	F41LCF47	105	X'F41LCF47'	BASE 10 0,97927562313	
00014A	FEEADBA	106	XC'FFFC294A'	BASE 10 0,99947710842	
01014C	FFFFFACT7	107	XC'FFFFFC72C'	BASE 10 0,9999999597312	
00015A	FFFFFFF4C	108	XC'FFFFFC74C'	BASE 10 0,999999945812	
00015A	FFFFFFFFFF	109	DC 10 0,999999999963	BASE 10 1,000100000000	
00015A	FFFFFFFFFF	110	DC X'FFFFFFFFFF'	BASE 10 1,000100000000	

LCC	OBJECT CTYPE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
00015C	F1F2A221	1111 0		DC	X*F1F2A221*
00016C	F21FF280	112		DC	X*F21FF280*
000164	FACB910C	113		DC	X*F21FCB910C*
00016A	F2C11720	114		DC	X*F2C11720*
0001AC	F2D5F418	115		DC	X*F2D5F418*
000170	F3074280	116		DC	X*F3074280*
000176	F31FF405	117		DC	X*F31FF405*
000178	F3261F80	118		DC	X*F3261F80*
000180	F35C2594	119		DC	X*F35C2594*
00017C	F377F605	120		DC	X*F377F605*
000185	F4021107	121		DC	X*F4021107*
00019A	F40A1107	122		DC	X*F40A1107*
000184	F42715AE	123		DC	X*F42715AE*
000192	F434114C	124		DC	X*F434114C*
00019C	F464720F	125		DC	X*F464720F*
000191	F495C2594	126		DC	X*F495C2594*
000194	F4921107	127		DC	X*F4921107*
00019A	F49A1107	128		DC	X*F49A1107*
00019C	F495C1CAF	129		DC	X*F495C1CAF*
0001A2	F4A671A9	130		DC	X*F4A671A9*
0001A6	F4C4JAF1	131		DC	X*F4C4JAF1*
0001A9	F4722F716	132		DC	X*F4722F716*
0001AC	F4A1112A	133		DC	X*F4A1112A*
0001AH	F4941C80	134		DC	X*F4941C80*
0001A6	F4C59A8F	135		DC	X*F4C59A8F*
0001A2	F4D4JAF1	136		DC	X*F4D4JAF1*
0001C1	F4E497468	137		DC	X*F4E497468*
0001C4	F4911E34	138		DC	X*F4911E34*
0001CA	F4931E94	139		DC	X*F4931E94*
0001CC	F4A90C57	140		DC	X*F4A90C57*
0001CD	F4A4C69F	141		DC	X*F4A4C69F*
0001D4	F4B3110D0	142		DC	X*F4B3110D0*
0001DA	F4B3199F	143		DC	X*F4B3199F*
0001E1	F4B4249F	144		DC	X*F4B4249F*
0001E5	F4B5A41A	145		DC	X*F4B5A41A*
0001E4	F4B74461	146		DC	X*F4B74461*
0001EA	F4B71A1A	147		DC	X*F4B71A1A*
0001EC	F4B8A357	148		DC	X*F4B8A357*
0001F0	F4C36990	149		DC	X*F4C36990*
0001FA	F4C3773	150		DC	X*F4C3773*
0001F8	F4C44745	151		DC	X*F4C44745*
0001FC	F4B13700	152		DC	X*F4B13700*
000200	F4B13700	153		DC	X*F4B13700*
000208	F4B65FB	154		DC	X*F4B65FB*
00020C	F4B60771	155		DC	X*F4B60771*
000211	F4B13700	156		DC	X*F4B13700*
000214	F4B705E58	157		DC	X*F4B705E58*
000218	F4B13700	158		DC	X*F4B13700*
00021C	F4B650C	159		DC	X*F4B650C*
000220	F4A658E	160		DC	X*F4A658E*
000224	F4A7012	161		DC	X*F4A7012*
000224	F4C4B35	162		DC	X*F4C4B35*
00022C	F4C221A	163		DC	X*F4C221A*
000230	F4E949A	164		DC	X*F4E949A*
000234	EF96A01	165		DC	X*F4F96A01*

LIN	OBJECT CODE	ACM1 ADDR2	STMT	SOURCE STATEMENT
000218	F0060D36		166	DC X'F8C0D043A6'
00021C	F01452AH		167	DC X'F81452AH'
000247	F01F294		168	DC X'F81F294'
000244	F0201DCA		169	DC X'F8201DCA'
000248	F032FE11		170	DC X'F832FE11'
00024C	F0B1CA79		171	DC X'F8B1CA79'
000249	F044544		172	DC X'F844544'
000254	F04C417C		173	DC X'F84C417C'
000254	F04FA491		174	DC X'F84FA491'
00025F	F177nFD	C	175	DC X'F177D5CF'
000260	F15F1061		176	DC X'F1FF15F1061'
000264	F24FC7F4		177	DC X'F24FC7F4'
000268	F2545637		178	DC X'F2545637'
00026C	F2CC0B9		179	DC X'F2CC0B9'
000271	F2T64655		180	DC X'F2T64655'
000274	F313600A		181	DC X'F313600A'
000274	F4104010		182	DC X'F4104010'
00027C	F319CFC7		183	DC X'F3210CFC7'
000280	F31CA378		184	DC X'F331CA378'
000284	F3181928		185	DC X'F34181928'
000284	F426C010		186	DC X'F426C010'
00028C	F49AEC59		187	DC X'F49AEC59'
00028D	F419KA3A		188	DC X'F419KA3A'
00028E	F550n271		189	DC X'F550n271'
000294	F54FFF12		190	DC X'F54FFF12'
00029C	F4nC576		191	DC X'F5nC576'
0002A3	F654P07		192	DC X'F5454P07'
0002A4	F4AEC41		193	DC X'F5AEC41'
0002A4	F6F18261		194	DC X'F56F3918261'
0002A4	F72Fn177		195	DC X'F72Fn177'
0002A6	F7640487		196	DC X'F7640487'
0002A6	F7328457		197	DC X'F7328457'
0002A8	F7CT7nR9		198	DC X'F7CT7nR9'
0002A9	F7FFCFY3		199	DC X'F7FFCFY3'
0002A9	F921510C		200	DC X'F821510C'
0002A4	F946AD3H		201	DC X'F846AD3H'
0002A8	F7C74AC		202	DC X'F7C74AC'
0002CC	F9A48F7		203	DC X'F9A48F7'
0002D1	F9215179		204	DC X'F9215179'
0002D4	F9E15179		205	DC X'F9E15179'
0002D4	F9Fn159		206	DC X'F8Fn159'
0002D4	F93761A		207	DC X'F93761A'
0002D4	F91759CE		208	DC X'F91759CE'
0002E4	F9843171		209	DC X'F9843171'
0002F4	F977F5905		210	DC X'F977F5905'
0002F4	F04540111		211	DC X'F94540111'
0002F4	F93761A		212	DC X'F93761A'
0002F4	F9CC90C		213	DC X'F9CC90C'
0002FA	F8E789C9		214	DC X'F9F8E789C9'
0002FA	F9039CC4		215	DC X'F9039CC4'
0002FC	F977F5905		216	DC X'F977F5905'
0002FC	FA19E797		217	DC X'F977F5905'
000304	FA2E264F		218	DC X'F977F5905'
000304	F4337605		219	DC X'F977F5905'
00030C	F452537F		220	DC X'F977F5905'
000310	F464n368			BASE 10 0.97816010835

LOC	OBJECT CODE	ACON1 ADDRESS	STMT	SOURCE STATEMENT
7002114	FAT3667	221	DC	X'FAT3667'
200311H	FAB144E5	222	DC	BASE IN 0.9783194676A
000311C	F491E916	223	DC	BASE IN 0.9785465211A
2003120	FAA87084	224	DC	BASE IN 0.9790413110A
000174	FAB112F	225	DC	BASE IN 2.97911222095
2003128	FAC1C7C	226	DC	BASE IN 0.9794177549
000312C	FAD74114	227	DC	BASE IN 0.9795402011A
0003130	FAD74116	228	DC	BASE IN 1.9794542115
0003134	F51101C	229	DC	BASE IN 0.98104646117
3001114	FAD1CFF8	230	DC	BASE IN 0.9824860111
200314C	F1115246	231	DC	BASE IN 0.983136211A
000314D	F81430E4	232	DC	BASE IN 3.98317410511A
200314E	FAD26114	233	DC	BASE IN 0.98477407495
000314F	FAD74110	234	P	BASE IN 0.9810234782
000314G	FH36E44P	235	DC	BASE IN 0.98119504112
000314H	F40D04C	236	DC	BASE IN 0.9813057099
000314I	F81430E4	237	DC	BASE IN 0.98145012546
0003150	F84C97A	238	DC	BASE IN 0.9815014669
200315C	FH4F3A91	239	DC	BASE IN 1.9816375593
000315D	FFF1111E	240	D	BASE IN 0.98164942111
3001144	FFF1F91A	241	DC	BASE IN 0.981653737
000315A	FFF1F91C	242	DC	BASE IN 0.981651779
00031AC	FFF1F91C	243	DC	BASE IN 0.981651779
2003173	FFFFFFF7	244	DC	BASE IN 0.981651779
0001174	00	245	D	BASE IN 0.981651779
0001174	CA	246	DC	BASE IN 0.981651779
2003176	0R	247	DC	BASE IN 0.981651779
2003177	10	248	DC	BASE IN 0.981651779
0003178	19	249	DC	BASE IN 0.981651779
0003179	20	250	DC	BASE IN 0.981651779
000317A	15	251	DC	BASE IN 0.981651779
000317B	25	252	DC	BASE IN 0.981651779
000317C	34	253	DC	BASE IN 0.981651779
000317D	01	254	DC	BASE IN 0.981651779
000317E	07	255	DC	BASE IN 0.981651779
000317F	01	256	DC	BASE IN 0.981651779
0003180	C5	257	DC	BASE IN 0.981651779
0003181	04	258	DC	BASE IN 0.981651779
0003182	C4	259	DC	BASE IN 0.981651779
0003183	04	260	DC	BASE IN 0.981651779
0003184	07	261	DC	BASE IN 0.981651779
2003185	09	262	DC	BASE IN 0.981651779
0003185	0C	263	DC	BASE IN 0.981651779
0003187	05	264	DC	BASE IN 0.981651779
0003188	0F	265	DC	BASE IN 0.981651779
0003189	1F	266	DC	BASE IN 0.981651779
000318A	14	267	DC	BASE IN 0.981651779
000318B	11	268	DC	BASE IN 0.981651779
000318C	14	269	DC	BASE IN 0.981651779
200318D	13	270	DC	BASE IN 0.981651779
000318E	12	271	DC	BASE IN 0.981651779
000318F	1E	272	DC	BASE IN 0.981651779
0003190	1C	273	DC	BASE IN 0.981651779
0003191	1E	274	DC	BASE IN 0.981651779
0003192	19	275	DC	BASE IN 0.981651779

LIN	OBJECT CNAME	ADDR1 ADDR2	STMT	SOURCE STATEMENT	PAGE
000393	17		274	DC X'1F'	4
000394	16		277	DC X'1A'	
000395	14		278	DC X'1A'	
000396	23		279	DC X'23'	
000397	15		280	DC X'1F'	
000398	22		281	DC X'22'	
000399	25		282	DC X'21'	
00039A	21		283	DC X'21'	
00039A	10		284	DC X'0'	
00039C	28		285	DC X'28'	
00039D	30		286	DC X'10'	
00039E	31		287	DC X'31'	
00039F	24		288	DC X'24'	
0003A0	2F		289	DC X'2F'	
0003A1	31		290	DC X'1F'	
0003A1	35		291	DC X'29'	
0003A2	29		292	DC X'26'	
0003A3	26		293	DC X'30'	
0003A4	39		294	DC X'26'	
0003A5	2E		295	DC X'37'	
0003A6	17		296	DC X'2A'	
0003A7	24		297	DC X'12'	
0003A8	32		298	DC X'20'	
0003A9	2C		299	DC X'2C'	
0003AA	2C		300	DC X'1F'	
0003AB	1F		301	DC X'3C'	
0003AC	1C		302	DC X'27'	
0003AD	21		303	DC X'4'	
0003AE	34		304	DC X'3A'	
0003AF	34		305	DC X'31'	
0003B0	11		306	DC X'38'	
0003B1	18		307	DC X'1D'	
0003B2	10		308	DC X'3H'	
0003B3	39		309	DC X'62'	
0003B4	02		310	DC X'01'	
0003B5	C3		311	DC X'03'	
0003B6	03		312	DC X'03'	
0003B7	C3		313	DC X'03'	
0003B8	01		314	DC X'03'	
0003B9	01		315	DC X'03'	
0003B9	01		316	DC X'14'	
0003B9	14		317	DC X'05'	
0003B9	14		318	DC X'0F'	
0003B9	5		319	DC X'CF'	
0003B9	CF		320	DC X'14'	
0003B9	OF		321	DC X'12'	
0003B9	OF		322	DC X'12'	
0003B9	C3		323	DC X'12'	
0003B9	12		324	DC X'12'	
0003B9	12		325	DC X'12'	
0003B9	12		326	DC X'1A'	
0003B9	14		327	DC X'25'	
0003B9	25		328	DC X'25'	
0003B9	25		329	DC X'25'	
0003B9	25		330	DC X'25'	

SOURCE STATEMENT
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LOC	OBJECT CODE	ADDR1 ADDR2	STMT	SOURCE STATEMENT
0001CA 21			331	DC X'21'
0001CA 21			332	DC X'21'
0001CC 10			333	DC X'1D'
0001CC 10			334	DC X'1D'
C001CF 10			335	DC X'10'
0001CF 2F			336	DC X'2F'
000300 2F			337	DC X'2F'
000301 3E			338	DC X'1F'
000302 1E			339	DC X'3E'
000303 3F			340	DC X'3F'
000304 2A			341	DC X'76'
000305 2B			342	DC X'26'
000306 2B			343	DC X'26'
000307 19			344	DC X'39'
000308 15			345	DC X'39'
000309 35			346	DC X'49'
00030A 36			347	DC X'39'
00030A 39			348	DC X'49'
00030A 15			349	DC X'2E'
00030C 2E			350	DC X'2F'
00030C 2F			351	DC X'2F'
00030E 2E			352	DC X'37'
00030F 17			353	DC X'37'
00030F 37			354	DC X'37'
000351 37			355	DC X'37'
000352 37			356	DC X'37'
000353 37			357	DC X'37'
000354 17			358	DC X'20'
000355 20			359	DC X'20'
000356 20			360	DC X'20'
000357 2C			361	DC X'20'
000358 2D			362	DC X'2C'
000359 2C			363	DC X'2C'
00035A 2C			364	DC X'2C'
00035C 2C			365	DC X'2C'
00035D 2C			366	DC X'2C'
00035E 2D			367	DC X'2C'
00035F 2C			368	DC X'2C'
000360 2C			369	DC X'3F'
000361 2C			370	DC X'3C'
000362 3C			371	DC X'3C'
000363 3C			372	DC X'3C'
000364 3C			373	DC X'3C'
000365 3F			374	DC X'3F'
000366 4F			375	DC X'27'
000367 1C			376	DC X'1A'
000368 3C			377	DC X'1A'
000369 3C			378	DC X'3A'
00036A 3A			379	DC X'3A'
00036B 27			380	DC X'3A'
00036C 3A			381	DC X'31'
00036D 3A			382	DC X'38'
00036E 3E			383	DC X'48'
00036F 3E			384	DC X'38'
000400 38				BASE 10 3.5000

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LIN	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
C00431	39			186	DC K*18*
300432	28			387	DC K*1A*
000433	30			388	DC K*3D*
000434	10			390	DC K*1D*
C00435	30			391	DC K*1D*
000436	10			392	DC K*3A*
C00437	14			393	DC K*3B*
C00438	34			394	DC K*3B*
C00439	34			395	DC K*3A*
C0047A	1A			396	DC K*3A*
C0047B	35			397	DC K*1A*
0004AC	02			398	DC K*1A?
0004CC	02			399	DC K*1B?
00040F	C3			400	DC K*1C5*
00040F	C4			401	DC K*1C5*
000410	C5			402	DC K*1A?
000411	C5			403	DC K*1A?
000412	C5			404	DC K*1A?
000413	C4			405	DC K*1A?
000414	C4			406	DC K*1A?
000415	C6			407	DC K*1A?
000416	C6			408	DC K*1A?
000417	04			409	DC K*1A?
000418	04			410	DC K*1A?
000419	07			411	DC K*1A?
00041A	07			412	DC K*1A?
00041B	CD			413	DC K*1A?
00041C	OF			414	DC K*1A?
00041D	CE			415	DC K*1A?
00041E	CF			416	DC K*1A?
00041F	OF			417	DC K*1A?
000420	OF			418	DC K*1A?
000421	OF			419	DC K*1A?
000422	11			420	DC K*1A?
000422	11			421	DC K*1A?
000424	16			422	DC K*1A?
000425	16			423	DC K*1A?
000426	16			424	DC K*1A?
000427	12			425	DC K*1A?
000428	1P			426	DC K*1A?
000429	1P			427	DC K*1A?
00042A	18			428	DC K*1A?
00042A	18			429	DC K*1A?
00042B	1A			430	DC K*1A?
00042C	1F			431	DC K*1A?
000427	1E			432	DC K*1A?
000428	1E			433	DC K*1A?
000429	1E			434	DC K*1A?
000431	1A			435	DC K*1A?
000432	1A			436	DC K*1A?
000433	1A			437	DC K*1A?
000434	1A			438	DC K*1A?
000435	1A			439	DC K*1A?
000436	1A			440	DC K*1A?
000437	1A				

LIN	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
000418 23					
000439 1F					
00043A 22					
00043B 22					
00043C 22					
00043D 25					
00043E 25					
00043F 21					
000440 21					
000441 21					
000442 17					
000443 17					
000444 10					
000445 1C					
000446 1C					
000447 26					
000448 30					
000449 26					
00044A 29					
00044B 26					
00044C 26					
00044D 26					
00044E 32					
00044F 20					
000450 2C					
000451 27					
000452 27					
000453 31					
000454 31					
000455 0C					
000456 0C					
000457 01					
000458 C4					
000459 C4					
00045A C4					
00045B C4					
00045C 08					
00045D 0A					
00045E C4					
00045F 0A					
000460 C4					
000461 DC					
000462 DC					
000463 FC					
000464 DC					
000465 CC					
000466 0D					
000467 0D					
000468 0D					
000469 10					
00046A 10					
00046B 16					
00046C 16					
00046D 16					
00046E 16					

LIC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
00046F 18		496		DC	X'14'
000479 1A		497		DC	X'15'
000471 1C		498		DC	X'1C'
000472 1C		499		DC	X'1C'
000473 1C		500		DC	X'1C'
000476 1C		501		DC	X'1C'
000475 1C		502		DC	X'1C'
000476 1E		503		DC	X'1F'
000477 1E		504		DC	X'19'
000478 1E		505		DC	X'19'
000479 1E		506		DC	X'19'
00047A 1E		507		DC	X'19'
00047B 1E		508		DC	X'19'
00047C 17		509		DC	X'17'
00047D 17		510		DC	X'17'
00047E 17		511		DC	X'17'
00047F 17		512		DC	X'17'
00047F 17		513		DC	X'17'
00047A 17		514		DC	X'17'
00047B 17		515		DC	X'17'
00047C 18		516		DC	X'18'
00047D 18		517		DC	X'17'
00047E 18		518		DC	X'17'
00047F 18		519		DC	X'17'
00047A 16		520		DC	X'17'
00047B 16		521		DC	X'1F'
00047C 16		522		DC	X'1F'
00047D 16		523		DC	X'19'
00047E 16		524		DC	X'20'
00047F 16		525		DC	X'20'
00047A 15		526		DC	X'20'
00047B 15		527		DC	X'19'
00047C 15		528		DC	X'19'
00047D 15		529		DC	X'19'
00047E 15		530		DC	X'19'
00047F 15		531		DC	X'20'
00047A 14		532		DC	X'20'
00047B 14		533		DC	X'20'
00047C 14		534		DC	X'20'
00047D 14		535		DC	X'20'
00047E 14		536		DC	X'20'
00047F 14		537		DC	X'20'
00047A 13		538		DC	X'20'
00047B 13		539		DC	X'20'
00047C 13		540		DC	X'20'
00047D 13		541		DC	X'20'
00047E 13		542		DC	X'01'
00047F 13		543		DC	X'11'
00047A 12		544		DC	X'11'
00047B 12		545		DC	X'04'
00047C 12		546		DC	X'04'
00047D 12		547		DC	X'08'
00047E 12		548		DC	X'08'
00047F 12		549		DC	X'07'
00047A 11		550		DC	X'08'

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LFC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
0004A6	09	551		DC	X'009'
0004A7	CC	552		DC	X'0C'
0004A8	0D	553		DC	X'0D'
0004A9	CE	554		DC	X'1F'
0004AA	10	555		DC	X'10'
0004AB	13	556		DC	X'13'
0004AC	02	557		DC	X'02'
0004AD	C2	558		DC	X'02'
0004AE	02	559		DC	X'02'
0004AF	C3	560		DC	X'03'
0004B0	01	561		DC	X'03'
0004B1	03	562		DC	X'03'
0004B2	05	563		DC	X'05'
0004B3	05	564		DC	X'05'
0004B4	C6	565		DC	X'06'
0004B5	C6	566		DC	X'06'
0004B6	C7	567		DC	X'07'
0004B7	0A	568		DC	X'0A'
0004B8	CA	569		DC	X'0A'
0004B9	C9	570		DC	X'09'
0004BA	CF	571		DC	X'0F'
0004BB	14	572		DC	X'14'
0004BC	15	573		DC	X'15'
0004BD	11	574		DC	X'11'
0004BE	12	575		DC	X'12'
		576	END		
		577			=F'37'
		578			=F'456'
		579			=F'3863'
		580			=F'3639'

0004C0 CC000075
 0004C2 C00000C8
 2004CA 000L0F17
 0004CC 00000E37

NOT REPRODUCIBLE

APPENDIX C

Assembler listing of GEN2

PAGE 1

LOC	ACT CODE	ADONI ADDR2	STAT	SOURCE STATEMENT	FORTRAN	PAGE
00000				1 GEN2 START 1		6/10/71
2				2		
3				3 SOURCE--L'F.CANNON, UNA. DEPT. OF STAT., APR 1970		
4				4 PURPOSE--TO SUPPLY THE USER WITH A FAST PROGRAM FOR GENERATING FA-		
5				5 PONENTIAL AND GAMMA DISTRIBUTED AND NUMBERS.		
6				6 USAGE--X-RANGAM(L,N). L MUST BE AN NON INTEGFR. IT PRIMES THE		
7				7 GENERATING SCHEM. N IS THE NUMBER OF POTENTIAL RANDOM		
8				8 NUMBERS SUMMED IN GENERATING A RANDOM NUMBER THAT IS GAMMA		
9				9 LA(N,L).		
10				10 EXAMPLE--RANGAM(55555,50) = WILL OF GAMMA(L=55555,R=1).		
11				11 NOTE--IF NOT AN INDEFINITLY DISTRIBUTED RANDOM NUMBER WILL BE		
12				12 GENERATED.		
13				13 METHOD--THE METHOD IS THAT OF MASSAGLIA AS DESCRIBED IN "CUMMING- CATION'S OF ACTM", VOL 6, NO. 5, MAY 1966. THE PROGRAM OBTAINED 14 FROM WILL FOLLOW THE NOTATION USED BY MASSAGLIA.		
15				15		
16				16		
17				17 ENTRY RANGAM PROLOGUE. STANDARD FOR 360/650.		
18				18 USING RANGAM, IS		
19				19 DC MILITARY.		
20				20 DC GL74 RANGAM.		
21				21 RANGAM STW 14-12-120131		
22				22 L 11-4-11 LOAD ADDRESS OF N INTN Q11.		
23				23 L 11-6-11 LOAD N INTN R11.		
24				24 IA 12-1 USED AS AN INDEX.		
25				25 IA 13-1 USED AS AN INCREMENT.		
26				26 SFR 14-1 RESULT STORED IN SFR.		
27				27 MARCH 0 TEMP EXECUTED ONLY ON THE FIRST CALL.		
28				28 M 4-RAT. OBTAIN NEW "U".		
29				29 ST 5-LA STORE U FOR NEXT GENERATION.		
30				30 SFR 6-1 SET FPR2=1. MW=0.		
31				31 SR 6-1		
32				32 NL 6-4		
33				33 SLDL 6-4		
34				34 C 6-4 TWO HIGH ORDER OCTAL DIGITS OF U IN R4(01012)		
35				35 C 6-4 P370 D1023>37 BRANCH NA.		
36				36 BC 11-NA D1023>37 BRANCH NA.		
37				37 TC 6-0+236(4) OBTAIN DISCRETE RANDOM VARIABLE FROM Q.		
38				38 SLL 5-6 R4=05047... R5=D50407... 39 BC 15-EX BRANCH FX.		
40				40 SLDL 4-3 R4=010203...		
41				41 C 6-4-4556* 4556(R4SF 12)=71048(E 1).		
42				42 AC 11-NA D1023>456 BRANCH NA.		
43				43 IC 6-0-144(4) OBTAIN DISCRETE RANDOM VARIABLE FROM Q.		
44				44 SLL 5-7 R5=D50407... 45 AC 15-EX BRANCH FX.		
46				46 SLDL 4-3 R4=010203...		
47				47 C 6-4 F 3P63* 3P63(R4SF 12)=74277(BASE PI).		
48				48 RC 11-NC C1023>3963 BRANCH NC.		
49				49 S 6-0+F 3639* OBTAIN PRIMER ADDRESSES FOR D TABLE.		
50				50 IC 6-0(4) OBTAIN DISCRETE VARIABLE FROM D.		
51				51 FX LR 6-6 R4= DISCRETE VARIABLE.		
52				52 SPOL 4-16 R5= DISCRETE VARIABLE + HI ORDER 16 BITS OF		
53				53 SJ 5-RES RS(MW LO DANE 16 BITS) WITH HI ORDER 8		
54				54 HWI RES,X42* BITS ZEROED OUT FOR CONVERSION TO FLOAT-		
55				55 AF 2-RES ING POINT. FPR2 CONTAINS RESULT + "W". 55 AER 0-2 FINAL RESULT SUMMED IN FPR0.		

LNC	OBJECT CODE	ADDR1 ADDR2	STAT	SOURCE STATEMENT
000000 97CA F016	0001E0 94EC 00C0	0001E J001C	56	PXLE 12-10. BRANCH L00P N TIMES.
000000 07FE	0001E0 5E51 F140	5A J001C	47	IN 14,12,12113) RESTORE REGISTERS.
000192 00C0 0203 F016	0001E0 5E51 F140	59 TEMP	60	BRANCH TO MAIN LINE.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	61	S1A REPLACES INSTR. AT R0ANCH.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	62	HVC BRANCH, TEMP REPLACE BRANCH HV T0MP.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	63	5,0,1,1 ADDRESS OF PRIMA IN R5.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	64	5,0,1,51 PR IMPR IN R5.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	65	ST SALA STOP DRAWER FOR USE IN URAN.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	66	BRANCH#4 BRANCH TO INSTRUCTION FOLLOWING "RANCH".
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	67	ST SALA RESTORE U TO R5.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	68	SP CL S,P(1)
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	69	AC 11,N0
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	70	SP 3,2
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	71	NE AR 3,2
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	72	CL 5,0,1,1
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	73	NC 11,NR
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	74	CL 5,C(1)
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	75	AC 11,N0
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	76	N 4,HPAT
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	77	ST SALA STORE FOR NEXT GENERATION.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	78	GN SR
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	79	SP1 3,2
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	80	IC 4,0,1,1
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	81	AC 15,0,1L
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	82	N 4,RPAT
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	83	LR N,? USED FOR P-DEFINING.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	84	AR 9,2 AA=10. OBTAIN DISCRETE VARIABLE.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	85	CL 5,2(R) U>Q11 Q11 CONTINUE LOOP.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	86	AC 11,0,-4 H 4,HPAT
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	87	ST 5,0,1,4 CATAIN NFW II.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	88	SA 8,2
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	89	LA 8,5
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	90	LN 9,2
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	91	SP 6,2
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	92	NH AC 11,N0
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	93	CLP 5,6
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	94	AC 11,0,-6
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	95	LR 6,5
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	96	R,9,0,NH
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	97	ST 5,0,1,4 LOAD SMALLEST II IN R6.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	98	C011F1R0 HAS OCTH INCREMENT & DELIMITED.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	99	ST 5,0,1,4 STORE FOR NEXT GENERATION.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	100	R,5,MIN(U1,U2,...,U10)).
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	101	AC 15,GN
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	102	SP 3,2 R>IP. U>P(IP) CONTINUE LOOP.
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	103	AC 11,N0
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	104	ST 3,J
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	105	MVI J,K+4,G
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	106	AF 2,J
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	107	AC 4,RPAT
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	108	ST 5,0,1,4 U>P(1).
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	109	AC 15,N1
000192 5E51 F140	0001E0 5E51 F140	0001E J001C	110	DS IF

LCC PROJECT CONF	ADDR1 ADDR2	STAT	SOURCE STATEMENT
000148		DS	IF
00014C F4C27194	111 LA	DS	IF
000154 CCCCCCCC	112 J	DS	XL4+4PC27195*
000154 00000000	113 NEAT	DC	
00015C F4A1CF47	114 Q	DC	4PC27195*
000160 FF1A94A	115	DC	K+OHC1H1H1H00*
000164 FF1FRC72	116	DC	K+OHC1H1H1H00*
00016A FF1FF4C	117	DC	K+OHC1H1H1H00*
00016C FF1FFF	118	DC	K+OHC1H1H1H00*
000170 FF1FF55	119	DC	X+FF1A94A
000174 F124721	120	DC	X+FF1A94A
000178 F231F19D	121	DC	X+FF1A94A
00017C F28C41M	122	DC	X+FF1A94A
000180 F2C1A730	123	DC	X+FF1A94A
000184 F2D5F414	124	DC	X+FF1A94A
000188 F31742M4	125	DC	X+FF1A94A
00019C F31742M4	126	DC	X+FF1A94A
00019E F31742M4	127	DC	X+FF1A94A
0001A0 F31742S	128	DC	X+FF1A94A
0001A4 F31742S	129	DC	X+FF1A94A
0001A8 F31742S	130	DC	X+FF1A94A
0001A9 F31742S	131	DC	X+FF1A94A
0001B0 F31742S	132	DC	X+FF1A94A
0001B4 F31742S	133	DC	X+FF1A94A
0001B8 F31742S	134	DC	X+FF1A94A
0001C4 F31742S	135	DC	X+FF1A94A
0001CA F31742S	136	DC	X+FF1A94A
0001C8 F31742S	137	DC	X+FF1A94A
0001D4 F31742S	138	DC	X+FF1A94A
0001E0 F31742S	139	DC	X+FF1A94A
0001E4 F31742S	140	DC	X+FF1A94A
0001E8 F31742S	141	DC	X+FF1A94A
0001F4 F31742S	142	DC	X+FF1A94A
0001FA F31742S	143	DC	X+FF1A94A
0001FC F31742S	144	DC	X+FF1A94A
0001FD F31742S	145	DC	X+FF1A94A
0001FA F414C69F	146	DC	X+FF1A94A
0001GA F4A20A6A	147	DC	X+FF1A94A
0001GC F471F734	148	DC	X+FF1A94A
0001GK F4C1C96	149	DC	X+FF1A94A
0001GK F4C1C96	150	DC	X+FF1A94A
0001GK F4C1C96	151	DC	X+FF1A94A
0001GK F4C1C96	152	DC	X+FF1A94A
0001GK F4C1C96	153	DC	X+FF1A94A
0001GK F4C1C96	154	DC	X+FF1A94A
0001GK F4C1C96	155	DC	X+FF1A94A
0001GK F5740661	156	DC	X+FF1A94A
000203 F5936A1A	157	DC	X+FF1A94A
000276 F5A36357	158	DC	X+FF1A94A
000123A F5C36RA0	159	DC	X+FF1A94A
0002JJC F9DC3771	160	DC	X+FF1A94A
00021W F5F42745	161	DC	X+FF1A94A
000214 FACA550A	162	DC	X+FF1A94A
00021A FA27C49A	163	DC	X+FF1A94A
10121C FA33C760	164	DC	X+FF1A94A
201177W FA4567F8	165	DC	X+FF1A94A

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LOC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT	
000224	F4190771	166		DC	*F1FA5A0D771	RASF 10 0.977701411485
000228	F4631F01	167		DC	*F1FA5A0D771	RASF 10 0.97769696411
00022C	F4766F54	168		DC	*F1FA5A0D771	RASF 10 0.97467642549
000271	F49FC1CA	169		DC	*F1FA5A0D771	RASF 10 0.97840779506
000274	F49F410C	170		DC	*F1FA5A0D771	RASF 10 0.97596467042
000278	F4AFA4FF	171		DC	*F1FA5A0D771	RASF 10 0.97939297367
00027C	F4AF00B2	172		DC	*F1FA5A0D771	RASF 10 0.97047750279
0002A0	F4CF4F415	173		DC	*FACT-0380	RASF 10 0.970161531
000246	F400003A	174		DC	*F1FA5A0D771	RASF 10 0.97706613301
00024A	F417063A	175		DC	*F1FA5A0D771	RASF 10 0.97461571204
00024C	F4F70A11	176		DC	*F1FA5A0D771	RASF 10 0.9816925649
000250	F4C4043A	177		DC	*F1FA5A0D771	RASF 10 0.974615295141
000254	F4161549	178		DC	*F1FA5A0D771	RASF 10 0.9701771647
00025A	F41F7996	179		DC	*F1FA5A0D771	RASF 10 0.96163746157
00025C	F42110C4	180		DC	*F1FA5A0D771	RASF 10 0.98170611594
000260	F412FF11	181		DC	*F1FA5A0D771	RASF 10 0.98156707942
000264	F432A179	182		DC	*F1FA5A0D771	RASF 10 0.97156112113
00026A	F444AA86	183		DC	*F1FA5A0D771	RASF 10 0.98101461370
0002CC	F4C417C	184		DC	*F1FA5A0D771	RASF 10 0.98163746157
000270	F44A371	185		DC	*F1FA5A0D771	RASF 10 0.98170611594
000274	F47701F0	186	C	DC	*F1FA5A0D771	RASF 10 0.98171060981
000278	F4161549	187		DC	*F1FA5A0D771	RASF 10 0.98156707942
00027C	F4FC2F58	188		DC	*F1FA5A0D771	RASF 10 0.97156112113
0002A0	F4240437	189		DC	*F1FA5A0D771	RASF 10 0.98101461370
000246	F4C40437	190		DC	*F1FA5A0D771	RASF 10 0.98163746157
00024A	F41F4106	191		DC	*F1FA5A0D771	RASF 10 0.98170611594
00024C	F417063A	192		DC	*F1FA5A0D771	RASF 10 0.98156707942
0002A0	F4161549	193		DC	*F1FA5A0D771	RASF 10 0.97156112113
000274	F429C277	194		DC	*F1FA5A0D771	RASF 10 0.98101461370
000278	F431C410	195		DC	*F1FA5A0D771	RASF 10 0.98163746157
00027C	F431C410	196		DC	*F1FA5A0D771	RASF 10 0.98170611594
0002A0	F4267070	197		DC	*F1FA5A0D771	RASF 10 0.98156707942
0002A4	F4751559	198		DC	*F1FA5A0D771	RASF 10 0.97156112113
0002A8	F4609346	199		DC	*F1FA5A0D771	RASF 10 0.98101461370
0002AC	F4540421	200		DC	*F1FA5A0D771	RASF 10 0.98163746157
0002B0	F40FF212	201		DC	*F1FA5A0D771	RASF 10 0.98170611594
0002B4	F465E554	202		DC	*F1FA5A0D771	RASF 10 0.98156707942
0002BA	F451F007	203		DC	*F1FA5A0D771	RASF 10 0.97156112113
0002BC	F44F1071	204		DC	*F1FA5A0D771	RASF 10 0.98101461370
0002C1	F4191A3	205		DC	*F1FA5A0D771	RASF 10 0.98163746157
0002C4	F422F0177	206		DC	*F1FA5A0D771	RASF 10 0.98170611594
0002D0	F40FF212	207		DC	*F1FA5A0D771	RASF 10 0.98156707942
0002D4	F4751559	208		DC	*F1FA5A0D771	RASF 10 0.97156112113
0002D8	F44F1071	209		DC	*F1FA5A0D771	RASF 10 0.98101461370
0002D9	F4191A3	210		DC	*F1FA5A0D771	RASF 10 0.98163746157
0002E0	F44A00A	211		DC	*F1FA5A0D771	RASF 10 0.98170611594
0002E4	F44A00A	212		DC	*F1FA5A0D771	RASF 10 0.98156707942
0002E8	F47C1FAC	213		DC	*F1FA5A0D771	RASF 10 0.9704926761
0002F4	F44455E7	214		DC	*F1FA5A0D771	RASF 10 0.97156112113
0002F8	F4C770H9	215		DC	*F1FA5A0D771	RASF 10 0.98101461370
0002F9	F47FFCFA1	216		DC	*F1FA5A0D771	RASF 10 0.98163746157
0002F9	F4F15379	217		DC	*F1FA5A0D771	RASF 10 0.98170611594
0002F0	F45F0158	218		DC	*F1FA5A0D771	RASF 10 0.98156707942
0002F4	F421F2CE	219		DC	*F1FA5A0D771	RASF 10 0.9735757562
0002F3	F43F1996	220		DC	*F1FA5A0D771	RASF 10 0.97421630871

LCC	OBJCT CNT	ADDBL ADRT?	STAT	SOURCE STATEMENT
000100	F57552905		221	NC X'F9755A619708
000104	F57552911		222	NC X'F9755A619711
000109	F9824014		223	DC X'F5D1741A'
00010C	F9CC980C		224	DC X'F9CFC920C'
000109	F8ED05C9		225	DC X'F5F5C919'
000114	AUJ674		226	DC X'F8A144C44'
000114	FAL2E797		227	DC X'F8A144C44'
000114	PAZL44F		228	DC X'F8A2E70F4
000120	FA3Y1B05		229	DC X'F4357405
CUR124	FA52467F		230	DC X'F8A55174
000132	F46A1A17		231	DC X'F8A612A0
0C932C	F3716C67		232	DC X'F8A7172F1
1D9130	F4915F59		233	DC X'F8A135F59
020134	FANNE54		234	DC X'F8A11911544
1CC132	FA2-1720A		235	DC X'F8A1732A9
000133C	FAP9127F		236	DC X'F8A1172F9
000143	FC515620		237	DC X'FAC515620
000143	FAD75114		238	DC X'FAD75114
0001364	FAL43406		239	DC X'F8A543406'
1D914C	FAF107C		240	DC X'F8A1447C'
1D9152	FA166118		241	DC X'F8A166118'
0001354	FH115464		242	DC X'FH115464
0001354	FH114564		243	DC X'FH114564
000135C	FH2454564		244	DC X'FH2454564
0001361	FR2P919		245	DC X'FH2P919
1D9164	FN34656A		246	DC X'FH34656A
0001359	FA4105469		247	DC X'FH4105469
000135C	FH459014		248	DC X'FH459014
0001370	FH4C97A1		249	DC X'FH4C97A1
000137	FRE6E911	P	250	DC X'FH6E911
0001379	FFEA101C		251	DC X'FFCA101C
000137C	FFFF09FA		252	DC X'FFF09FA
0001382	FFFEE1C		253	DC X'FFE1C
000134	FFFFFFF7		254	DC X'FFF7
760134	FFFFFFFFFF		255	DC X'FFFFFFF0
000139C	UN	D	256	DC X'0000
000139A	UN		257	DC X'00A0
7D01AF	CR		258	DC X'1100
000139F	10		259	DC X'1100
000139A	15		260	DC X'1100
0001391	2C		261	DC X'2000
0001397	03		262	DC X'3000
0001398	05		263	DC X'5000
0001393	24		264	DC X'7000
0001394	36		265	DC X'1000
0001355	C1		266	DC X'0200
0001394	07		267	DC X'1140
000139C	07		268	DC X'0050
000139D	09		269	DC X'0070
000139F	0C		270	DC X'0090
000139F	0D		271	DC X'00B0
000139F	07		272	DC X'00D0
000139D	09		273	DC X'00F0
000139F	0C		274	DC X'0110
000139F	0D		275	DC X'0124

LIC	OBJECT NAME	OWNER ACCOUNT	STAT	SCOPCF STATEMENT	PAGE
J001140 0E			276	DC X*1F*	4
J001141 0F			277	DC X*1F*	
00019A2 14			278	DC X*14*	
J001343 11			279	DC X*11*	
J001346 16			280	DC X*16*	
J001347 13			281	DC X*13*	
J001348 12			282	DC X*12*	
J001347 18			283	DC X*18*	
J001349 15			284	DC X*15*	
J001345 16			285	DC X*16*	
J001347 16			286	DC X*16*	
J001346 17			287	DC X*17*	
J001345 15			288	DC X*15*	
J001340 14			289	DC X*14*	
J001343 10			290	DC X*10*	
J001344 26			291	DC X*1F*	
J001345 1F			292	DC X*22*	
J001343 22			293	DC X*25*	
J001341 25			294	DC X*11*	
J001342 21			295	DC X*11*	
J001343 10			296	DC X*10*	
J001344 26			297	DC X*30*	
J001345 10			298	DC X*11*	
J001346 13			299	DC X*24*	
J001347 24			300	DC X*25*	
J001348 2F			301	DC X*2F*	
J001349 3F			302	DC X*3F*	
J001345 29			303	DC X*29*	
J001344 26			304	DC X*35*	
J001345 31			305	DC X*19*	
J001346 2F			306	DC X*8E*	
J001347 37			307	DC X*37*	
J001348 24			308	DC X*24*	
J001349 32			309	DC X*12*	
J001345 26			310	DC X*20*	
J001347 2C			311	DC X*2C*	
J001343 3F			312	DC X*1F*	
J001344 3C			313	DC X*3C*	
J001345 27			314	DC X*27*	
J001346 34			315	DC X*34*	
J001347 14			316	DC X*11*	
J001348 21			317	DC X*39*	
J001349 18			318	DC X*13*	
J001344 3D			319	DC X*39*	
J001345 28			320	DC X*20*	
J001346 28			321	DC X*02*	
J001347 02			322	DC X*03*	
J001348 01			323	DC X*13*	
J001349 03			324	DC X*63*	
J001345 01			325	DC X*01*	
J001346 03			326	DC X*13*	
J001347 05			327	DC X*75*	
J001348 05			328	DC X*13*	
J001349 06			329	DC X*14*	
J001345 14			330	DC	

LIN	PROJECT CODE	ANRRL ADDP?	STMT	SOURCE STATEMENT	DC	X'270	BASE	IN 2.6375
00040E	27		386	DC	X'1A1	BASE	IN 3.6410	
00040F	1A		387	DC	X'1A1	BASE	IN 3.6410	
200410	3A		388	DC	X'1A0	BASE	IN 3.6450	
000411	3A		389	DC	X'1A0	BASE	IN 3.6450	
200412	1A		390	DC	X'1A0	BASE	IN 3.6450	
000412	1A		391	DC	X'1A0	BASE	IN 3.6450	
000413	1A		392	DC	X'1A0	BASE	IN 3.6450	
000414	3A		393	DC	X'1A1	BASE	IN 3.6450	
200414	1A		394	DC	X'1A1	BASE	IN 3.6450	
000415	3A		395	DC	X'1A1	BASE	IN 3.6450	
200415	1A		396	DC	X'1A0	BASE	IN 3.6450	
000416	1A		397	DC	X'1A0	BASE	IN 3.6450	
000417	1A		398	DC	X'1A0	BASE	IN 3.6450	
000418	3A		399	DC	X'1A0	BASE	IN 3.6450	
000419	3A		400	DC	X'1A0	BASE	IN 3.6450	
000419	1A		401	DC	X'1A0	BASE	IN 3.6450	
000420	3A		402	DC	X'1A0	BASE	IN 3.6450	
000420	1A		403	DC	X'1A0	BASE	IN 3.6450	
000421	3A		404	DC	X'1A0	BASE	IN 3.6450	
200421	1A		405	DC	X'1A0	BASE	IN 3.6450	
000422	1A		406	DC	X'1A0	BASE	IN 3.6450	
000423	3A		407	DC	X'1A0	BASE	IN 3.6450	
100424	1A		408	DC	X'02*	BASE	IN 3.6450	
000425	02		409	DC	X'02*	BASE	IN 3.6450	
000426	02		410	DC	X'02*	BASE	IN 3.6450	
200426	02		411	DC	X'02*	BASE	IN 3.6450	
000427	05		412	DC	X'02*	BASE	IN 3.6450	
200427	05		413	DC	X'02*	BASE	IN 3.6450	
000428	05		414	DC	X'02*	BASE	IN 3.6450	
000428	04		415	DC	X'02*	BASE	IN 3.6450	
000429	05		416	DC	X'02*	BASE	IN 3.6450	
200429	05		417	DC	X'02*	BASE	IN 3.6450	
00042F	06		418	DC	X'1A0	BASE	IN 3.6450	
00042F	06		419	DC	X'02*	BASE	IN 3.6450	
20042F	06		420	DC	X'02*	BASE	IN 3.6450	
000431	0C		421	DC	X'02*	BASE	IN 3.6450	
000431	07		422	DC	X'02*	BASE	IN 3.6450	
000431	07		423	DC	X'02*	BASE	IN 3.6450	
000432	06		424	DC	X'02*	BASE	IN 3.6450	
000432	06		425	DC	X'02*	BASE	IN 3.6450	
200432	06		426	DC	X'02*	BASE	IN 3.6450	
000437	0F		427	DC	X'02*	BASE	IN 3.6450	
000437	0F		428	DC	X'02*	BASE	IN 3.6450	
200437	0F		429	DC	X'02*	BASE	IN 3.6450	
000438	16		430	DC	X'12*	BASE	IN 3.6450	
000438	16		431	DC	X'12*	BASE	IN 3.6450	
200438	16		432	DC	X'12*	BASE	IN 3.6450	
000439	16		433	DC	X'16*	BASE	IN 3.6450	
200439	16		434	DC	X'16*	BASE	IN 3.6450	
00043F	12		435	DC	X'12*	BASE	IN 3.6450	
000440	1A		436	DC	X'1A*	BASE	IN 3.6450	
000441	1B		437	DC	X'1B*	BASE	IN 3.6450	
200442	1B		438	DC	X'1B*	BASE	IN 3.6450	
000443	1B		439	DC	X'1B*	BASE	IN 3.6450	
000444	1E		440	DC	X'1E*	BASE	IN 3.6450	

LRC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT	PC
0004445	1E	441	DC	X'1E'	RASF 10 1.8750	
0004446	1F	442	DC	X'1F'	BASE 10 1.8750	
3C447	15	443	DC	X'19'	BASE 10 1.8750	
0004448	1A	444	DC	X'1A'	BASE 10 1.8750	
0004445	1R	445	DC	X'1A'	BASE 10 1.8750	
000444A	1B	446	DC	X'1A'	BASE 10 1.8750	
000444A	1A	447	DC	X'1A'	BASE 10 1.8750	
000444C	14	448	DC	X'1A'	BASE 10 1.8750	
303444D	1A	449	DC	X'1A'	BASE 10 1.8750	
203444F	1A	450	DC	X'1A'	BASE 10 1.8750	
62C444F	1A	451	DC	X'1A'	BASE 10 1.8750	
3034450	23	452	DC	X'23'	BASE 10 2.0750	
0004451	1F	453	DC	X'1F'	BASE 10 2.0750	
0004452	22	454	DC	X'22'	BASE 10 2.0750	
3034453	22	455	DC	X'22'	BASE 10 2.0750	
3034454	22	456	DC	X'22'	BASE 10 2.0750	
0004455	25	457	DC	X'25'	BASE 10 2.0750	
0004456	25	458	DC	X'25'	BASE 10 2.0750	
0004457	21	459	DC	X'21'	BASE 10 2.0750	
3034458	21	460	DC	X'21'	BASE 10 2.0750	
3034459	21	461	DC	X'21'	BASE 10 2.0750	
0004455	1C	462	DC	X'1D'	BASE 10 2.0750	
0004454	1C	463	DC	X'1D'	BASE 10 2.0750	
000445H	10	464	DC	X'10'	BASE 10 2.0750	
000445C	10	465	DC	X'10'	BASE 10 2.0750	
000445D	10	466	DC	X'10'	BASE 10 2.0750	
000445E	24	467	DC	X'28'	BASE 10 2.0750	
000445F	28	468	DC	X'28'	BASE 10 2.0750	
000446C	1D	469	DC	X'1A'	BASE 10 2.0750	
3034461	2F	470	DC	X'2F'	BASE 10 2.0750	
2034462	29	471	DC	X'29'	BASE 10 2.0750	
0004463	26	472	DC	X'26'	BASE 10 2.0750	
0004464	26	473	DC	X'26'	BASE 10 2.0750	
0004465	2E	474	DC	X'12'	BASE 10 2.0750	
0004466	12	475	DC	X'20'	BASE 10 2.0750	
3034467	2C	476	DC	X'2C'	BASE 10 2.0750	
3034468	2C	477	DC	X'27'	BASE 10 2.0750	
0004469	27	478	DC	X'27'	BASE 10 2.0750	
000446A	27	479	DC	X'11'	BASE 10 2.0750	
303446B	31	480	DC	X'10'	BASE 10 2.0750	
303446C	30	481	DC	X'10'	BASE 10 2.0750	
000446D	C6	482	DC	X'10'	BASE 10 2.0750	
000446F	00	483	DC	X'00'	BASE 10 0.0000	
000445F	C1	484	DC	X'01'	BASE 10 0.0023	
0004470	C4	485	DC	X'04'	BASE 10 0.0250	
3034471	C4	486	DC	X'04'	BASE 10 0.0250	
3034472	04	487	DC	X'04'	BASE 10 0.0250	
3034473	09	488	DC	X'09'	BASE 10 0.0475	
0004474	09	489	DC	X'09'	BASE 10 0.0475	
0004475	C3	490	DC	X'0B'	BASE 10 0.0625	
0004476	08	491	DC	X'08'	BASE 10 0.0625	
3034477	C9	492	DC	X'18'	BASE 10 0.0875	
0004478	0A	493	DC	X'1A'	BASE 10 0.0875	
0004479	0C	494	DC	X'0C'	BASE 10 0.1750	
000447A	CC	495	DC	X'0C'	BASE 10 0.1750	
000447B	CC	496	DC	X'0C'	BASE 10 0.1750	

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LLOC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
100447C CC		496		DC	X'00C'
000470 0C		497		DC	X'00C'
00047E CC		498		DC	X'00C'
20047F 0D		499		DC	X'00D'
000480 05		500		DC	X'005'
200481 10		501		DC	X'00A'
000482 10		502		DC	X'00A'
200483 16		503		DC	X'00B'
200484 16		504		DC	X'00B'
000485 16		505		DC	X'00A'
200486 16		506		DC	X'00A'
000487 1A		507		DAKF	X'19A'
200488 1A		508		DC	X'19A'
000489 1C		509		DC	X'19C'
20048A 1C		510		DC	X'19C'
00048B 1C		511		DC	X'19C'
20048C 1C		512		DC	X'19C'
20048D 1C		513		DC	X'19C'
20048E 1C		514		DC	X'19C'
00048F 19		515		DC	X'19F'
20048G 19		516		DC	X'19F'
00048H 16		517		DC	X'19H'
20048I 16		518		DC	X'19I'
00048J 16		519		DC	X'19J'
20048K 19		520		DC	X'19K'
00048L 17		521		DC	X'19L'
20048M 17		522		DC	X'19M'
00048N 17		523		DC	X'19N'
20048P 17		524		DC	X'19P'
00048Q 17		525		DC	X'19Q'
20048R 17		526		DC	X'19R'
00048S 1A		527		DC	X'19S'
20048T 1A		528		DC	X'19T'
00048U 23		529		DC	X'19U'
20048V 21		530		DC	X'19V'
20048W 1F		531		DC	X'19W'
00048X 1F		532		DC	X'19X'
20048Y 1A		533		DC	X'19Y'
00048Z 20		534		DC	X'19Z'
20048A 21		535		DC	X'19A'
20048B 20		536		DC	X'19B'
00048C 20		537		DC	X'19C'
20048D 19		538		DC	X'19D'
00048E 24		539		DC	X'19E'
20048F 24		540		DC	X'19F'
20048G 24		541		DC	X'19G'
00048H 23		542		DC	X'19H'
20048I 23		543		DC	X'19I'
00048J 24		544		DC	X'19J'
20048K 24		545		DC	X'19K'
00048L 24		546		DC	X'19L'
20048M 24		547		DC	X'19M'
00048N 24		548		DC	X'19N'
20048O 16		549		DC	X'19O'
00048P 03		550		DC	X'00P'

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LIN	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
551				DC	X'0000
552	000484 00			DC	X'0001
553	000485 01			DC	X'0001
554	000486 01			DC	X'0010
555	000497 01			DC	X'0100
556	000498 04			DC	X'0400
557	000499 04			DC	X'0400
558	0004AA CP			DC	X'0400
559	000499 08			DC	X'0400
560	0004AC 07			DC	X'0800
561	0004B2 C4			DC	X'0800
562	0004BF 09			DC	X'0001
563	0004BF 0C			DC	X'0001
564	0004C0 C0			DC	X'0001
565	0004C1 0E			DC	X'0001
566	0004C2 10			DC	X'1000
567	0004C3 11			DC	X'1300
568	0004C4 02			DC	X'0200
569	0004C5 02			DC	X'0200
570	0004C6 C2			DC	X'07F0
571	0004C7 C2			DC	X'1200
572	0004CH 04			DC	X'0300
573	0004C9 C3			DC	X'0300
574	0004CA 05			DC	X'0500
575	0004C9 05			DC	X'0005
576	0004CC C6			DC	X'0006
577	0004CF 05			DC	X'0005
578	0004CF 07			DC	X'0007
579	0004CF 04			DC	X'0A00
580	0004D0 04			DC	X'0A00
581	0004D1 09			DC	X'0009
582	0004D2 VF			DC	X'000F
583	0004D3 14			DC	X'1400
584	0004D4 15			DC	X'1500
585	0004D5 11			DC	X'1100
586	0004D6 12			DC	X'1200
587				END	
588					=F'37*
589					=F'456*
590					=F'3863*
591					=F'3639*

APPENDIX D

Assembler Listing of GEN3

NOT REPRODUCIBLE

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LOC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
				MVI	RFC,X'0C10
000FC92	92C1 0302	C0114		56	15,0+6
000FC96	AFF3 00AC		0009F	57	MVI
000FC9A	9241 43n2	00314		58	RES,X'41'
000FC9E	7401 1102		00314	59	AF
C00142	5ED0 32A6		002C8	A0	0'RES
1000AA	98F9 000C		000RC	61	L 13,4FFA+6
0000AA	92FF 300C	0000C		62	LM 14,0,12(13)
C00140	61F5 92FE			63	MVI 12(12),X'FF+
0001406	0201 501E	0009E	(00020	64	PCR 15,14
000140A	5853 1000	00010	00020	65	REPLACES INSTR. AT BRANCH.
000140E	5857 50C0	00020	00020	66	MVC
0001412	5050 A7FE	00020	00020	67	BKINCH, TEMP
0001416	47F3 9012	00024	00024	68	L 5,0+11
000141A	5C53 92FE	00310	00310	69	L 5,0+11
000141E	4121 0004	00024	00024	70	L 5,0+11
0001422	5E51 A69E	00020	00020	71	L 2,6
0001426	41A1 41AE	00020	00020	72	LA 2,6
000142A	1162	00020	00020	73	CL 5,0+11
000142E	1A67			74	BC 11,14
0001432	554A 83C6			75	IF U>FFF4FLNFO
0001436	47A7 90CA	00308		76	IF U>FFF4FLNFO(BASS 16) GO TO L4.
000143A	5555 46B8	00020		77	AR 6,2
000143E	47A3 90FC	00020		78	CL 5,0+6
0001442	5C41 A2FA	00020		79	CL 11,0+6
0001446	5E53 92FE	0030C		80	AC 11,0,LA 4,RPAT
0001450	4951 0002	00310		81	ST 5,0+6
0001454	47F3 A1B4	00020		82	SR 6,2
0001458	5C43 92FA	00020		83	L5 5,0+6
0001462	1F75			84	4,RPAT
0001466	5C43 92FA	00020		85	7,5
0001470	5051 A7FE	00030		86	4,RPAT
0001474	1575	00020		87	SY 5,0+6
0001478	47A3 9106	00018		88	CLQ 7,5
0001482	1E57	00020		89	RC 11,0,LA 4,RPAT
0001486	5E51 92FE	00030		90	7,5
0001490	4861 0001	00020		91	LA 7,5
0001494	47A3 9126	00020		92	SLI 6,1
0001498	4864 8546	00058		93	CLR 3,0+6
0001502	1D24	00020		94	AC 11,0,LA 4,RPAT
0001506	1P37	00020		95	SLI 6,1
0001510	8A12 0010	000010		96	LW 4,RPAT
JU-12A	1537			97	SLI 7,5
0001514	47A3 9126	00138		98	ST 5,0+6
0001518	4864 8546	00001		99	CLR 3,0+6
0001522	1D24	00020		100	AC 11,0,LA 4,RPAT
0001526	1P37	00020		101	SLI 6,1
0001530	8A12 0010	000010		102	LW 4,RPAT
JU-12A	1537			103	SLI 7,5
0001534	47A3 9126	00138		104	IF MAX(U1,U2)<0111 GO TO LS.
0001538	4864 8546	00001		105	R7=MAX-MIN.
0001542	1D24	00020		106	SLI 7,5
0001546	1P37	00020		107	ST 7,5,11,IF CONV TO FLOATING POINT.
0001550	8A12 0010	000010		108	AF 2,0+6
0001554	3822	00280		109	SEP 2,0+6
0001558	4864 8546	000010		110	SRL 5,1,2

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LRC	OBJECT CODE	ACDR1 ADDR2	STMT	SOURCE STATEMENT
C00159	9C60 92A2	00284	111	ST S,M MVI W,X+400
00015C	9240 92A2	C02A4	112	LF 4,4F+0,06275.
00016C	7E40 A6C2	00604	113	LER
000169	2864		114	SF 6,4
00016A	7E40 92A2	00284	115	AE 6,4
00016A	7A60 92A2	00284	116	SQ 3,3
00016F	1A31		117	IC 3,A+716161
000170	4A16 A67E	00600	118	R3=TABLE RANDOM VARIABLE.
000170	9913 0001	00001	119	SIL 3,1
000174	5033 92A2	00284	120	R3=R302.
00017C	5245 92A2	00284	121	MVI W,X+450
00017A	7A63 92A2	00284	122	AE 6,4
00016E	3F46	00008	123	MFP 4,6
00016B	7E43 96C0	00284	124	ME 4,4F+0,050
00016A	7043 92A2	00284	125	SIE 4,4
00016A	7043 92A2	00284	126	CALL EXP(M)
00016A	7B01 B6C4	0060C	128	SE 0,0
00016A	85A3 FCC2	00002	129	SLL 6,2
00016F	7FC6 A106	00118	140	MF 0,F(6)
000162	5A53 A2FF	00310	141	S,LA
00016A	3902		142	CER 0,7
00016A	4703 90FC	00LFE	143	AC 11,L8
2CC16C	5A53 0293	002AC	144	L,MIN
00016D	7E03		145	SER 0,7
000167	8P63 0002	00002	146	SRL 6,2
00016A	1A44		147	LS 4,4
00016A	4346 A67E	00600	148	IC 6,A+716161
16016C	470 40606	00000	149	FI 0,0
000163	5E53 92FE	00000	150	L,LA
000164	1B77		151	S,R 7,7
000166	5C60 92F4	0030C	152	H,4,RPAT
00016A	1A66		153	S,R 6,4
000169	5019 82FF	00310	154	S,T 5,LA
00016F	FF53 00008	00000	155	SRI 5,9
000164	5053 8260	00000	156	ST 5,11
00016A	5A53 97FE	00310	157	L,LA
000165	1F74		158	MR 6,5
00016E	1F74	00310	159	ALR 7,6
00016A	5E53 97FE	0030C	160	L,LA
000164	5C60 92FA	0030C	161	H,4,RPAT
000161	5053 92FE	00310	162	ST 5,LA
00016C	1B46		163	S,R 4,4
00016F	9053 0008	00008	164	S,T 5,8
000202	5053 9244	002RC	165	ST 5,02
00020A	5A53 97FE	00310	166	L,LA
000224	1C45		167	MR 6,5
00020C	1E74		168	ALR 7,4
00021F	4710 916E	00100	169	RO L4
F.00212	A870 0C09	00008	170	SRL 7,8
000216	5C10 97AE	002C0	171	ST 7,5
00021A	9241 92AE	002C0	172	MVI 5,40+
00021F	7460 92AE	002C0	173	AE 0,5
100222	7C91 92AE	002C0	174	SIE 0,5
00023E	7C00 8ACE	00600	175	CALL 4,0G,(S)
			187	HE 0,0,E-2,0.
				CALCULATE INT101002+02*02*0.

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LRC	OBJECT CODE	ADDR	ADDR2	STMT	SOURCE STATEMENT
000242	7A00 9402	006F4		168	AF 0.=E*9.0*
000242	7D00 92AE	007C0		189	DF 0.5
000242	7D00 92AE	002C0		197	STE 0.5
000246	9246 9246	002B8	66288	151	CALL SD,-(S)
00026A	9222			203	PVI U1,X=40*
00026C	7A20 9246	002A8		205	SEP 2.2
000270	3C20			206	AF 2.01
000272	7920 8606	0029A		207	MER 2.1
000274	6720 827E	00290		208	CE 2.=F*3.0*
00027A	9240 42A4	002BC		209	AC 2.110
00027E	3R22			210	PVI U2,X=40*
00027F	7A21 9244	002AC		211	SER 2.2
000280	1C70			212	AF 2.02
000284	1P00			213	MER 2.1
000294	7-21 8606	006FA		214	SFR 0.0
00029C	47C1 91AE	001E0		215	CE 2.=F*1.0*
00029C	5P50 42F5	00110		216	AC 1.1L4
0002A6	9C43 A2FA	0C10C		217	L 5.LA
0002CA	5053 82FE	07310		218	5.RPAT
0002F1	11245			219	LTP 5.LA
0002F5	6720 9297	002A4		220	DP 0.5
0002A2	3123			221	LNFR 2.7
0002A4	18C2			222	LER 0.2
0002A6	47FA 8090	007A2		223	A OUT
0002AC				224	MIN DS
0002A6				225	DIF DS
0002A6				226	W DS
0002HA				227	U1 DS
00027C				228	U2 DS
00027C				229	S DS
0002C6				230	AREA DS
00030C	48C27795			231	BPAT DS
000310				232	LA DS
000314				233	RFS DS
000318	42139300			234	E DC E'16.-0*
00031C	41051773			235	DC E'13.343857463*
000320	42139138			236	DC E'19.-6.92e768*
000324	42111111			237	DC E'12.-0.5664466
000328	41F72556			238	DC E'14.-1.5-6759
00032C	41C42720			239	DC E'12.-6.3-55216*
000330	42124925			240	DC E'1A.-2.5-14278*
000334	41F0340F			241	DC E'15.-0.2-2223*
00033A	41A66396			242	DC E'10.-4-124568*
00033C	41RA6FA4			243	DC E'11.-4-197461*
000340	42155555			244	DC E'21.-3-13333*
000344	4198C29F			245	DC E'9.-54-2115044*
000344	4160340F			246	DC E'28.-4.4-4444*
00034C	419E6FAF			247	DC E'9.-95-59629
000350	42174521			248	DC E'23.-272727*
000356	42266249			249	DC E'36.-57-12856*
00035F	422JC100			250	DC E'32.-0-
00035C	4192463C			251	DC E'9.-165-202116*
000360	418705E22			252	DC E'8.-491-31467*
000364	41A0C1682			253	DC E'8.-816C1104*

LCC	OBJCT/CONE	ACCR1	ACCR2	STMT	SOURCE	STATEMENT
000168	422A/AAA			254	DC	E142-05466666666
000174	417A32F			255	DC	E17-637254400
000175	41800046			256	DC	E18-167666639
000176	42333333			257	DC	E151-199000000
000177	427F741			258	DC	E17-901368655
000178	418271A5			259	DC	E17-152602877
000179	424610C0			260	DC	E16-6
000180	4150568E			261	DC	E14-3445742
000181	416969C1			262	DC	E1A-525576115
000182	4187A77E			263	DC	E1A-931567219
000183	42800100			264	DC	E12R-0
000184	41629702			265	DC	E17-161078474
000185	415444CF			266	DC	E15-267374
000186	41111110			267	DC	E175H-10
000187	416546C2			268	DC	E16-3-10693163
000188	415944126			269	DC	E15-3420055
000189	415A2A11			270	DC	E15-635362894
000190	41590006			271	DC	E15-517357695
000191	41501510			272	DC	E15-620506743
000192	41515130			273	DC	E15-1-1117028
000193	41524241			274	DC	E14-934270510
000194	41672455			275	DC	E11-4-94677734
000195	4187A761			276	DC	E145-3333333
000196	42555544			277	DC	E175-6-10
000197	42199904			278	DC	E10-1-5101050
000198	41A677A6			279	DC	E17-1347696660
000199	417A3313			280	DC	E14-722616549
000200	41402268			281	DC	E15-931865632
000201	4159F5F0			282	DC	X1F70C2088-
000202	40C91F8			283	AC	X1F8212221-
000203	418659D3			284	DC	X1CAAF7303-
000204	41A77-00			285	DC	X1F8A7740-
000205	41402268			286	DC	X1FAT-CF40-
000206	4153F337			287	DC	X1FG3213974
000207	41A07274			288	DC	X1F9P002260-
000208	4SC0CF0			289	DC	X1F9C00450-
000209	41415B			290	DC	X1F9F4500-
000210	4A35A1DE			291	DC	X1F4A1110-
000211	4A717C49			292	DC	X1FA4-7-FR9-
000212	41A740n7			293	DC	X1F0B21160-
000213	4AD039			294	DC	X1FA0-3-702-
000214	41035455			295	DC	X1F811CA5-
000215	416210F4C1AE			296	DC	X1FB4C190-
000216	417C76A9			297	DC	X1FB7C1968-
000217	41402268			298	DC	X1FAC1502-
000218	418A1AC			299	DC	X1F0B21160-
000219	41C72008			300	DC	X1FC002240-
000220	4C314070			301	DC	X1FC31-070-
000221	416969E2			302	DC	X1FC51-062-
000222	41A740F9			303	DC	X1FCB1100-
000223	41A401F7			304	DC	X1FC462067-
000224	41A740FF			305	DC	X1FCAC1ER-
000225	41C657E			306	DC	X1FCFF57E-
000226	41C690E			307	DC	X1FD0E850DF-
000227	4031E552			308	DC	X1FD41F952-

LCC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE	STATEMENT
000444	FF66019F	309		OC	X*FH746039F*	
000444	FF591010	310		OC	X*FD9103U*	
00044C	FC7015AF	311		OC	X*FH70158F*	
000450	F0AE5116	312		OC	X*FD8E5016*	
000454	FCA316A	313		OC	X*FDA329E6*	
000458	F0P7A1AF	314		OC	X*FD01A0AF*	
000474	FF215FF6	315		OC	X*FD01A2FA*	
000478	FF2227B3	316		OC	X*FD01A50E*	
00047C	FF651111	317		OC	X*FDCA77*	
00047D	FF81110F	318		OC	X*FD0101D*	
00048C	FE00E49	319		OC	X*FD11FD49*	
00047A	FF110F0C	320		OC	X*FF110F0C*	
00047A	FF215FF6	321		OC	X*FD214FF6*	
00047A	FF2227B3	322		OC	X*FD2127B3*	
00047C	FF651111	323		OC	X*FD01A50E*	
00047D	FF81110F	324		OC	X*FD0101D*	
000486	FF513125	325		OC	X*FF013425*	
00048A	FF611425	326		OC	X*FF01A14D*	
0004AC	FF10C2A3	327		OC	X*FF10C2A3*	
0004A9	FF13C0FA	328		OC	X*FF13C0E3*	
000496	FF4119FD	329		OC	X*FF411050*	
0004A8	FF7C1C75A	330	C	OC	X*FD7C1C75A*	
0004A9	FF117206	331		OC	X*FF117206*	
0004A9	FF611407	332		OC	X*FD611407*	
0004A9	FF4119FD	333		OC	X*FF411050*	
0004A9	FF4119FD	334		OC	X*FP41101A*	
0004AC	FF10C452	335		OC	X*FD0101D*	
0004A0	FF31C777	336		OC	X*FD917277*	
0004A6	FF9165782	337		OC	X*F9165782*	
0004AA	FF4119FD	338		OC	X*FF411050*	
0004C0	FF215906	339		OC	X*FD215906*	
0004C7	FF40F0C1	340		OC	X*FA0193C1*	
0004C4	FA9A1C6	341		OC	X*FA0110C6*	
0004CC	FA0111AF	342		OC	X*FA01112F*	
0004C4	FF4119FD	343		OC	X*FD4119FD*	
0004C4	FF4119FD	344		OC	X*FD4119FD*	
0004C4	FF4119FD	345		OC	X*FD4119FD*	
0004C4	FF4119FD	346		OC	X*FD4119FD*	
0004C4	FF4119FD	347		OC	X*FD4119FD*	
0004C4	FF4119FD	348		OC	X*FD4119FD*	
0004C4	FF4119FD	349		OC	X*FD4119FD*	
0004C4	FF4119FD	350		OC	X*FD4119FD*	
0004C4	FF4119FD	351		OC	X*FD4119FD*	
0004C4	FF4119FD	352		OC	X*FD4119FD*	
0004C4	FF4119FD	353		OC	X*FD4119FD*	
0004C4	FF4119FD	354		OC	X*FD4119FD*	
0004C4	FF4119FD	355		OC	X*FD4119FD*	
0004C0	FF215906	356		OC	X*FD215906*	
0004C4	FF10C456	357		OC	X*FD0101D*	
0004C0	FF4119FD	358		OC	X*FD4119FD*	
00051C	F16F5A6	359		OC	X*FD0101D*	
0004A0	FF31A9A	360		OC	X*FD9176A*	
000514	FF9140B8	361		OC	X*FD01ACB8*	
000518	FF0914FA	362		OC	X*FD0104FA*	
00051C	FF0HE9C9	363		OC	X*FD05EDC9*	

LNC	PROJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
000520	F00A4966			364	X'FD00RRR064.'
000524	FFP72A8A			365	DC X'FDF22F8A'
00052A	FFE47D92			366	DC X'FD0F49D92'
00052C	FFE51290			367	DC X'FF055A99'
000529	FF1299E8			368	DC X'FE1229E8'
000514	FF22D621			369	DC X'FE22D621'
000518	FF22C7971			370	DC X'FF2C0971'
00051C	FF6957888			371	DC X'FF6957888'
0005140	FFE2310			372	DC X'FE7F2310'
0005144	FFACF946			373	DC X'FECFC946'
0005148	FF223046			374	DC X'FFE23046'
0005152	FF1732F6			375	DC X'FF1732F6'
0005156	FF2FAEC3			376	DC X'FF2546F3'
000554	FF15796H			377	DC X'FF3F7940'
000554	FF1A			378	DC X'FF75'
000554	FF1A			379	DC X'FF39'
00055C	FC29			380	DC X'FC29'
00055E	FF74			381	DC X'FF75'
000560	FF7D			382	DC X'FF77'
000562	FFF1			383	DC X'FFF9'
000564	FD71			384	DC X'FD71'
000565	FFE1			385	DC X'FFHF'
000568	FF15			386	DC X'4E35'
00056A	FFB8			387	DC X'FEB6'
00056C	FAF1			388	DC X'F1F1'
00056E	FDFA			389	DC X'FDFA'
000570	FFC2			390	DC X'F5C2'
000572	FF35			391	DC X'FFB5'
000574	FF54			392	DC X'F952'
000576	FOA1			393	DC X'FD2A'
000578	C374			394	DC X'F374'
00057A	FER2			395	DC X'F0A2'
00057C	FD7D			396	DC X'FD7D'
00057F	FOH2			397	DC X'FD02'
000580	FO57			398	DC X'FD07'
000582	FCED			399	DC X'CE07'
000584	FO2F			400	DC X'FD92F'
000586	F93A			401	DC X'FA31'
000588	FD2F			402	DC X'FD2F'
00058A	FCAC			403	DC X'FCAC'
00058C	E189			404	DC X'FLK9Y'
00058E	FPAS			405	DC X'FA45'
000590	FO2D			406	DC X'FC27'
000592	FC6A			407	DC X'FA31'
000594	FC33			408	DC X'FC6A'
000594	FRE7			409	DC X'FC0C3'
000594	FA22			410	DC X'FRF7'
000594	A02D			411	DC X'FD007'
00059C	FC29			412	DC X'FC29'
00059E	FA22			413	DC X'FR22'
0005A0	FA66			414	DC X'FB64'
0005A0	FA9F			415	DC X'FA64'
0005A0	FAE1			416	DC X'FA7F'
0005A0	FA9F			417	DC X'FAE1'
0005A0	FA9F			418	DC X'FA9F'

LFC	OBJECT CODE	AC001 AD002	STMT	SOURCE STATEMENT
			419	DC X17EF9*
			420	DC X0645*
			421	DC X1F6RD*
			422	DC X1F76*
			423	DC X1FCAC*
			424	DC X1FCBA*
			425	DC XFB45*
			426	A
			427	DC X001*
			428	DC X001*
			429	DC X001*
			430	DC X001*
			431	DC X001*
			432	DC X002*
			433	DC X002*
			434	DC X002*
			435	DC X002*
			436	DC X002*
			437	DC X002*
			438	DC X002*
			439	DC X002*
			440	DC X002*
			441	DC X002*
			442	DC X002*
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			447	DC X002*
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			462	DC X002*
			463	DC X002*
			464	DC X002*
			465	DC X002*
			466	DC X002*
			467	DC X002*
			468	DC X002*
			469	DC X002*
			470	DC X002*
			471	DC X002*
			472	DC X002*
			473	DC X002*

LINC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
000150	C6			474	DC X'060
000150	C7			475	DC X'070
00015A	C7			476	DC X'070
00015A	C7			477	DC X'070
00015C	C5			478	DC X'070
00015D	C6			479	DC X'060
00015E	C9			480	DC X'090
00015F	CA			481	DC X'0A0
000160	C4			482	DC X'0A0
00015F	0F			483	DC X'0F0
00015F	C9			484	DC X'0F0
00015F	C5			485	DC NC
10015F	C9			486	DC NC
10015F	C9			487	DC NC
00015A	C5			488	DC NC
00015A	C5			489	DC NC
000157	C5			490	DC NC
00015B	C8			491	DC NC
00015B	C0			492	DC NC
00015A	C1			493	DC NC
00015C	C1			494	DC NC
00015D	C1			495	DC NC
00015E	C1			496	DC NC
00015F	C1			497	DC NC
000160	C1			498	DC NC
000151	C1			499	DC NC
000152	C1			500	DC NC
000153	C1			501	DC NC
000154	C1			502	DC NC
000155	C1			503	DC NC
000156	C1			504	DC NC
000157	C1			505	DC NC
000158	C1			506	DC NC
000159	C1			507	DC NC
00015A	C1			508	DC NC
00015B	C1			509	DC NC
00015C	C1			510	DC NC
00015D	C1			511	DC NC
00015E	C1			512	DC NC
00015F	C1			513	DC NC
000160	C10			514	DC X'100
000141	C1F			515	DC X'1F0
000122	C21			516	DC X'200
000113	C20			517	DC X'200
000101	C10			518	DC X'200
000101	C22			519	DC X'220
000101	C23			520	DC X'230
000101	C21			521	DC X'210
000101	C27			522	DC X'270
000100	C00			523	DC X'000
000101	C01			524	DC X'010
000101	C05			525	DC X'050
000101	C05			526	DC X'050
000101	C05			527	DC X'050
000101	C05			528	DC X'050

LOC	OBJFCR CODE	ADDR1 ADDR2	STAT	SOURCE STATEMENT
000015	05			DC X'05*
000020	07			DC X'07*
000021	07			DC X'07*
000022	09			DC X'09*
000023	08			DC X'08*
000024	08			DC X'08*
000025	08			DC X'08*
000026	08			DC X'08*
000027	08			DC X'08*
000028	00			DC X'00*
000029	0F			DC X'0F*
00002A	10			DC X'10*
00002B	11			DC X'11*
00002C	12			DC X'12*
00002D	12			DC X'12*
00002E	12			DC X'12*
00002F	12			DC X'12*
00002F	13			DC X'13*
000030	13			DC X'13*
000031	13			DC X'13*
000032	14			DC X'14*
000033	14			DC X'14*
000034	15			DC X'15*
000035	14			DC X'14*
000036	15			DC X'15*
000037	19			DC X'19*
000038	19			DC X'19*
000039	19			DC X'19*
00003A	15			DC X'15*
00003A	19			DC X'19*
00003B	14			DC X'14*
00003C	1C			DC X'1C*
00003D	1C			DC X'1C*
00003E	1C			DC X'1C*
00003F	1C			DC X'1C*
000040	1C			DC X'1C*
000041	10			DC X'10*
000042	10			DC X'10*
000043	10			DC X'10*
000044	10			DC X'10*
000045	10			DC X'10*
000046	10			DC X'10*
000047	10			DC X'10*
000048	10			DC X'10*
000049	21			DC X'21*
00004A	21			DC X'21*
00004B	22			DC X'22*
00004C	22			DC X'22*
00004D	26			DC X'26*
00004E	25			DC X'25*
00004F	24			DC X'24*
000050	00			DC X'00*
00004C	01			DC X'01*
000052	04			DC X'04*
000051	07			DC X'07*
000053	07			DC X'07*
000054	07			DC X'07*
000055	07			DC X'07*

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LIC	OBJECT CODE	AC001 ADDR2	STMT	SOURCE STATEMENT
0000016	0:		584	DC X'0E'
0000017	JE		585	DC X'7E'
0000018	OF		586	DC X'0F'
0000019	10		587	DC X'10'
0000020	10		588	DC X'10'
0000021	11		589	DC X'11'
0000022	11		590	DC X'11'
0000023	12		591	DC X'12'
0000024	14		592	DC X'14'
0000025	14		593	DC X'14'
0000026	14		594	DC X'14'
0000027	15		595	DC X'15'
0000028	15		596	DC X'15'
0000029	15		597	DC X'15'
0000030	15		598	DC X'15'
0000031	15		599	DC X'15'
0000032	17		600	DC X'17'
0000033	18		601	DC X'18'
0000034	18		602	DC X'18'
0000035	18		603	DC X'18'
0000036	18		604	DC X'18'
0000037	1C		605	DC X'1C'
0000038	1C		606	DC X'1C'
0000039	1F		607	DC X'1F'
0000040	1F		608	DC X'1F'
0000041	21		609	DC X'1E'
0000042	21		610	DC X'1E'
0000043	24		611	DC X'21'
0000044	24		612	DC X'24'
0000045	24		613	DC X'24'
0000046	24		614	DC X'24'
0000047	24		615	DC X'24'
0000048	25		616	DC X'25'
0000049	25		617	DC X'25'
0000050	25		618	DC X'25'
0000051	26		619	DC X'25'
0000052	26		620	DC X'25'
0000053	26		621	DC X'25'
0000054	26		622	DC X'25'
0000055	26		623	DC X'25'
0000056	26		624	DC X'25'
0000057	29		625	DC X'29'
0000058	29		626	DC X'29'
0000059	29		627	DC X'29'
0000060	29		628	DC X'29'
0000061	29		629	DC X'29'
0000062	29		630	DC X'29'
0000063	24		631	DC X'2A'
0000064	24		632	DC X'2A'
0000065	24		633	DC X'2A'
0000066	24		634	DC X'2A'
0000067	25		635	DC X'2B'
0000068	25		636	DC X'2B'
0000069	25		637	DC X'2B'
0000070	25		638	DC X'2B'

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LRC	OBJCT	CNT	ADR1	ADR2	STMT	SOURCE STATEMENT
260647	25				639	DC X*2D*
CON1AF	2E				640	DC X*2E*
CON1BF	2F				641	DC X*2F*
CON1CF	OF				642	DC X*0F*
CON1D91	12				643	DC X*12*
CON1E92	CC				644	DC X*CC*
CON1E91	0F				645	DC X*9E*
CON1E94	11				646	DC X*11*
CON1E95	13				647	DC X*13*
CON1E95	UN				648	DC X*UD*
CON1E97	10				649	DC X*10*
CON1E98	17				650	DC X*17*
CON1E99	15				651	DC X*15*
CON1E94	C4				652	DC X*7B*
CON1E94	19				653	DC X*19*
CON1E9C	04				654	DC X*09*
CON1E97	19				655	DC X*19*
CON1E9F	04				656	DC X*0A*
CON1E9F	04				657	DC X*06*
CON1E9D	C7				658	DC X*07*
CON1E91	14				659	DC X*1A*
CON1E92	1E				660	DC X*1E*
CON1E91	19				661	DC X*1A*
CON1E94	05				662	DC X*95*
CON1E95	1F				663	DC X*1F*
CON1E96	10				664	DC X*10*
CON1E97	74				665	DC X*64*
CON1E9P	1F				666	DC X*1E*
CON1E99	21				667	DC X*21*
CON1E9A	C3				668	DC X*53*
CON1E9H	28				669	DC X*24*
CON1AC	26				670	DC X*24*
CON1AD	22				671	DC X*22*
CON1AF	01				672	DC X*01*
CON1EF	26				673	DC X*2C*
CON1E92	2C				674	DC X*2C*
CON1AH	01				675	DC X*01*
CON1E91	25				676	DC X*25*
CON1E93	2A				677	DC X*28*
CON1E94	29				678	DC X*29*
CON1E94	02				679	DC X*2A*
CON1E95	2A				680	DC X*2E*
CON1E97	2C				681	DC X*20*
CON1E94	2F				682	DC X*2F*
CON1E93	14				683	DC X*14*
CON1E94	02				684	DC X*02*
CON1E95	09				685	DC X*09*
CON1E9C	16				686	DC X*16*
CON1E91	20				687	DC X*20*
CON1AE	23				688	DC X*23*
CON1AF	27				689	DC X*27*
					690	END "F*46*
					691	"F*475"
					692	"F*3981"
					693	

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LIN	OBJECT CODE	ADDR1	ADDR2	SIMT	SOURCE STATEMENT
0096CC	0000E77	694			=F'3703.
000600	FF4F13F0	695			=X'FF4F10F0.
3Cn604	4017r1C0	696			=E'0.0625.
3n06DP	4C801700	697			=F'0.5.
009A0C	41130000	698			=E'1.0.
0096F0	C1200000	699			=E'1-2.0.
0096E4	41150r00C	700			=E'9.0.
0006F8	41330000	701			=E'3.0.

NOT REPRODUCIBLE

APPENDIX E

Assembler Listing of GEN4

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PROGRAM 6/3C/71

LIC	OBJECT CODE	ADDR	STRT	SOURCE STATEMENT
00000				1 GENA
				2 *****
				3 * SOURCE---L.E. TANINIGA, CFP. OF STATE JUN 1970
				4 * PURPOSE--TO SUPPLY THE USER WITH A FAST RANDOM NUMBER.
				5 * AIMALLY USEFUL FOR RANDOM NUMBER.
				6 * USAGE---1. CALL ASSETUP(I,IR). I IS THE PROBABILITY OF A SUCCESS.
				7 * IR IS THE RANGE OF THE VARIABLE. ASSETUP SETS UP THE TABLE
				8 * FROM WHICH RANDI GENERATES A RANDOM NUMBER. BSETUP ALFFN AF
				9 * CALLED ONLY ONCE FOR A GIVEN SET OF PARAMETERS.
				10 * 2. XEPARITION. IR1 MUST BE AN ODD POSITIVE INTEGER. LT
				11 * PRIMES THE GENERATING SCHEME. X IS THE RANDOM NUMBER GENFE.
				12 * ATTN. IT IS RETURNED IN THE REAL NAME.
				13 * EXAMPLE--CALL ASSETUP(1,5.2)
				14 * 0011 11100
				15 * X=RANDI(5.25)
				16 *
				17 * METHOD---THE METHOD IS THAT OF MARAGLIA AS DESCRIBED IN "COM. ACM"
				18 * 6(MARAGLIA, 1961), 37.
				19 *****
				20 ASSETUP
				21 ENTRY ASSETUP
				22 R 126,115
				23 RC XL1,77
				24 DC [LT, ASSETUP]
				25 STM 14,12,12(13)
				26 ALR 1n,n
				27 USING w,1n
				28 LR 16,13
				29 L4 13,AF4
				30 ST 13,R(10,16)
				31 ST 14,4(R,13)
				32 L 7,n,11
				33 LF 2,6(1,21) FPR2=P.
				34 LF 4,4(1,0) FPR4=I-P=Q.
				35 LER 6,6
				36 SFR 4,2
				37 STE 6,4(1,1)
				38 L 6,6,6) R6=N,
				39 LR 5,4 CONVERT N TO WORDS.
				40 SLL 5,2
				41 SER 6,4
				42 ST 6,NUM CONVERT TO FLOATING POINT.
				43 HVI NIM,X,16*
				44 AF A,NUM
				45 L 7,AF-1*
				46 AR 6,7 MER 0,4
				47 RMM 6,7,8-2 FPR3=QeeN,
				48 LA 6,4
				49 SP 5,1
				50 SER 4,4
				51 LA 12,1
				52 CE 0,LLIM
				53 AC 6,LO BRANCH TO LN
				54 AC 6,LO OTHERWISE STORE R3=0 INTO XST.
				55 ST 3,XST

NOT REPRODUCIBLE

NOT REPRODUCIBLE

LPC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
000674	7C03 A2F6	002C8	56	STF	0,P(1)
00071N	1A34		57	AR	3,4
120C7A	7A49 AF4E	002C9	58 L1	AF	4,-E+1,0*
0007F	1C06		59	DER	0,6
0007G	1C06		60	MER	0,6
120C7F	7C00 A2F2	002A4	61	DE	0,0
120344	7C02	002C4	62	MER	0,2
120C4A	7C03 A2F3	002A4	63	CE	0,LLIM
120C4C	4247 40B8	002A4	64	BC	4,COVER
120C4D	7C03 A2F6	002A4	65	STF	0,P(1)
0007G	7C04 A2F6	002C9	66	SF	6,-E+1,0*
120350	4734 A2F8	002A4	67	AXE	3,6,LL
120C4C	47F3 A0FA	002C9	68	CONVQ	
120C4H	7A49 AF4E	002C9	69 L0	AE	4,-E+1,0*
120C4I	3C04		70	DFR	0,4
120C4K	1C04		71	MFR	0,6
120C4L	7D00 42F2	002A4	72	DF	0,0
120C4M	1C07	002C9	73	MFR	0,2
120C4F	1A56	002C9	74	SR	5,4
120C4O	7B49 AF4E	002C9	75	SE	6,-E+1,0*
120C4P	7C03 A2F6	002C9	76	4	12,-F+1,
120C4Q	7C03 A2F2	002C4	77	CE	0,LLIM
120C4R	4743 A0B8	002A0	78	BC	4,L3
120C4S	50C9 A2A6	002C0	79	ST	12,-X57
120C4T	1A54	002C0	80	SR	5,4
120C4U	47F7 A068	0007A	81	AC	15,L1
120C4V	1A53		82 CONVER	LA	5,3
120C4W	1A54		83	SR	5,4
120C4X	1M33		84	SR	3,3
120C4Y	1A23 A2R6	002C8	85 L3	LE	2,P(1)
120C4Z	7F20 A67A	006FC	86	AU	2,7,RA
120C5A	7021 A2P6	002C8	87	STE	2,P(1)
120C5C	5A91 A2P6	002C8	88	L	9,P(1)
120C5D	5A91 A2P6	00008	89	SLL	9,8
120C5E	1BBC	002C8	90	ST	9,P(1)
120C5F	1P49	00000	91	BXL	3,4,L3
120C5G	5C54 42A6	002C8	92	SP	12,12
120C5H	5A91 A2P6	00000	93	SR	11,11
120C5I	1A77		94	SR	7,7
120C5J	1A66		95	SR	6,4
120C5K	1A33		96	SR	3,1
120C5L	1M89		97 L4	SR	8,8
120C5M	5A93 A2A6	002C8	98	L	9,P(3)
120110	1A69	00004	99	SLDL	8,4
120112	8D81 9064	00004	100	AR	6,8
000106	1A78	00004	101	SLDL	8,4
000108	AERO 0004	00004	102	SLDL	8,4
00011C	1ARR	00004	103	SLDL	8,4
00011E	AM1 0004	00004	104	AR	11,8
000112	1AC3	00004	105	SLDL	8,4
000114	8714 A0E4	000F6	106	AR	12,8
00011A	5063 A6A6	006C8	107	AXE	3,4,L4
00011C	SC77 A6A4	006CC	108	ST	6,5
00011U	5080 AAE	006D0	109	ST	7,5+4
			110	ST	11,5+8

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AND STORE QNEW INTO P(0).

FP4+X0.

FP7=PI(X0/(X0+1)).

FP0=PI(X0*(N-X)/(X0+1)*Q).

FP0=PI(X0*P1(N-X)/(X0+1)*Q).

IF FP0>PI*X-LIM

BRANCH TO CONVER

OTHERWISE STORE FP0 INTO P(1).

CONTINUE

THIS SECTION OF CODE

COMPUTES SUCCESSIVE

PROBABILITIES UNTIL A

PROBABILITY=LLIM IS

ATTAINED, TMF PROPR

STARTING POINT IS

THEN STORED

AND A BRANCH IS

MADE IN L1 FOR FURTHER CALCULATIONS.

R3=1.

CLEAR RA.

R0=P(1).

R4=0.

R6=SUM(L1).

R8=0102.

R7=SUM(L1D2).

R8=010203.

R11=SUM(L1D204).

R12=SUM(L1D2034).

CONTINUE.

STORE SUM(L1) INTO S(1).

STORE SUM(L1D2) INTO S(1).

STORE SUM(L1D203) INTO S(1).

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LINE	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT		FOBAPR70	PAGE	
000124	59C9 A6C2	00004	111	ST	12,5+12			6/30/71	
00012A	A9B3 0004	00004	112	SLL	11,4	R11=SUM(10120111916.			
00012C	A570 0004	00004	113	SLL	7,4	R7=SUM(10216.			
000130	8460 7004	00004	114	SLL	6,4	R6=SUM(1011916.			
000134	1A11		115	SR	3,3	R3=1.			
000136	1AB9		116 L5	SP	8,8	CLEAR R4.			
000138	5853 A2B6	00008	117	L	9,P(3)	P9=P(1).			
00013C	90P0 C004	00004	118	STD	8,4	R8D1.			
000141	1A64		119	SR	6,8	R6=SUM(1011916-SUM(D1)).			
000142	1A78		120	SR	7,8	R7=SUM(10216-SUM(D2)).			
000144	1553		121	SR	11,8	R11=SUM(10120111916-SUM(D3)).			
000146	1A63		122	SR	8,8	CLEAR R9.			
00014A	40P7 0004	00004	123	SR	8,4	R8D2.			
00014C	1878		124	SR	7,8	R7=R7-SUM(D2).			
00014F	1B64		125	SR	11,8	R11=R11-SUM(D2).			
000150	1AB4		126	SR	8,8	CLEAR R4.			
000152	8701 0004	00004	127	SLDL	8,4	RA=03.			
000156	1P44		128	SP	11,8	911=R11-SUM(D3).			
00015A	A124		129	BXLF	3,6,LS	CONTINUE.			
00015C	1ACC		130	SP	12,12	STORE R6 INTO N(3).			
00015E	4C63	ADFA	00008	131	ST	6,4	STORE R7 INTO N(1).		
00015F	SC79	ADCA	0000C	132	ST	7,4	STORE R11 INTO N(2).		
000160	60A1 A6CE		00000	133	ST	11,8+N(8)			
000164	1A77		134	SR	7,7				
00016C	1A0C		135	SP	0,0				
00016E	41A1 1001	00001	136	LA	6,1	R3=1.			
000172	1A31		137 L6	SR	3,3	LOAD R2 WITH STARTING VALUE.			
000173	1A64		138 C0	119 L7	SR	8,8	CLEAR R4.		
000174	5951 A2B6	00020	140	L	9,P(1).	R9=P(1).			
00017F	ACB3 0034	00004	141	STD	8,4	R8D1,R8=D2,P8=03,P9=04.			
000182	SC93 A2B5		142	ST	9,P(1).	RESET ONE SHIFTED P(1).			
000184	5950 AFIA		143	C	8,F(1).	ANCH TO LB.			
00018A	47A0 4197		001A4	144	AC	AR 7,4			
00018F	1A78		145	AR	7,4	OTHERWISE STORE R7(DISPFET Q.V.)			
000190	5970 AFRE	00009	146	C	7,8+F7C30.	CHECK THAT TABLE LENGTH IS NOT EXCEEDED.			
000194	4729 ALAO	001R2	147	RC	2,L9	THF NUMBER OF TIMES CORRESPONDING			
000195	1A74		148	SR	7,4	TO THE DIGIT IN R8.			
00019A	422C A6CF	000FC	149	STC	2,A(12)				
00019E	81C4 A1P0	00194	150	BXLF	12,6,6-4				
0001A1	1AT5		151	AR	7,4	INCREMENT DISCRETE R.V.			
0001A4	LA74		152 L8	AP	2,6	CONTINUE.			
0001A5	A734 A166	00178	153	AXLE	3,6,L7	CONTINUE.			
0001A4	LA04		154	AP	0,6	CONTINUE 16 LDNP			
0001EC	1S06		155	CR	0,4	UNTIL FIRST 4 DIGITS OF			
0001EF	6740 A169	00172	156 L9	BC	4,16	EACH P(1) HAVE BEEN SHIFTED.			
000192	5A30 A6C2	00064	157 L9	L	3,5,12	THIS SECTION OF CODE			
00019A	5A11 A6CE	006E0	158	S	3,6,8	RESETS N(3)			
00019A	1S1C		159	CR	3,12	IF TABLE LENGTH IS EXCEEDED.			
00019C	4790 A1RA	001CA	160	BC	8,NUT				
0001C0	1B3C		161	SR	3,12				
0001C2	5A30 A6CE	006F0	162	A	3,4,8				
0001C6	5029 A6CF	005Fn	163	ST	3,N+R				
0001CA	5A01 A25E	0027n	164 OUT	DS	0H				
			165	L	13,AREA#4				

NOT REPRODUCIBLE

LOC	OBJECT CODE	ACRRE	ADDR2	SIMT	SOURCE STATEMENT
00115E	98FC D0NC	000NC	000NC	166	LW 14,12,12(13)
00115D	92FF D0NC	(000C		167	BVI 12(13),X1FFF
10115E	C7F			168	BCR 15,14
				169	DNPB 10
				170	ENTRY RANAL
				171	USING RANAL
				172	DR. X1107,
				173	DC CL7, RANAL.
				174	STM 1,3,24(13)
				175	SFR 0,0
				176	BRANCH 6 TEMP
				177	H SPAT
				178	SR 3,LA
				179	SR 2,2
				180	SR 0,0
				181	SLDL 2,4
				182	IF 01<5 THEN
				183	BRANCH TN FX
				184	R2-D107.
				185	IF 0102>=S111 THEN
				186	BRANCH TN 47
				187	S 2,N
				188	OTHERWISE OBTAIN PROPER TABLE ADDRESS
				189	BRANCH TN EX.
				190	R2-D1020.
				191	IF 01020>=S121 THEN
				192	BRANCH TN 43
				193	S 2,N6
				194	OBTAIN PROPER TABLE ADDRESS
				195	BRANCH TN FX.
				196	R2-D1020104.
				197	IF 01020304>=S131 THEN
				198	GO TO BRANCH/DATA IN NEW U.
				199	OBTAIN PROPER TABLE ADDRESS.
				200	DS 0,A12)
				201	ST 0,RES
				202	MVI PE5,0,A64.
				203	NORMALIZED FLT. PT. RESULT IN FOR.
				204	AF 0,RES
				205	LW 1,3,24(13)
				206	RETURN.
				207	REPLACES THE INSTRUCTION AT BRANCH.
				208	L 1,LA REPLACE INSTRUCTION AT BRANCH.
				209	MVC 1,0,11 LOAD PRIORITELY DONE ON FIRST CALL.
				210	L 3,G,1,2)
				211	3,LA BRANCH#4.
				212	DS 1,AF
				213	DS 1,IF
				214	DS 1,F
				215	DC E10,0,F,0105*
				216	DS 256F
				217	DS 4F
				218	DS 3F
				219	DC X14AC27395*
				220	DS 1,F

NOT REPRODUCIBLE

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LCC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
0006FC	40010000	221	ZERO	DC	X'400000000000'
0006F0		222	A	DS	2000XL1
0005CC	41100000	223		END	
0005C4	FFFFFFFFFF	224			=E1.0*
300FC4	0CC00001	225			=F-1*
300FC0	00C30000	226			=F1*
000F00	00C007D0	227			=F0*
		228			=F12000*

NOT REPRODUCIBLE

APPENDIX F

Assembler Listing of GEN5

PAGE 1

LCC	OBJECT CODE	ACM1 ADDRESS	START	SOURCE STATEMENT
000000				1 GFNS START A PRINT DM,MDEN,NODATA
				2 * 3 * 4 * SOURCE---Lef. CARTRIDGE ACCEPT. OF STATE. JUN 1972 5 * PURPOSE--TO SUPPLY THE USER WITH A FAST PROCEDURE FOR GENERATING 6 * PRIMES: DISTINQUISHED RANDOM NUMBERS. 7 * USAGF---CALL PSETUP(LAM). LAM IS THE MEAN OF THE POISSON DISTRIBUTION TO BE GENERATED. PSETUP SETS UP THE TABLE FROM 8 * RANDOM GENERATOR. IT PRIMES 9 * ONLY ONCE FOR A GIVEN LAM. 10 * 2. X=RANPO(LIND). "LIND" MUST BE AN ADD INTEGER. IT PRIMES 11 * THE CIPHERING SCHEME. X REPRESENTS THE POISSON VARIABLE 12 * DD 1=1..100 13 * X=RANPO(156951)
				14 *
				15 *
				16 * METHOD---THE METHOD IS THAT OF MARSAGLIA AS DESCRIBED IN "COMA. ACM 17 * 14 * 15 * 16 * 17 * 18 * 19 * 20 * 21 * 22 * 23 * 24 * 25 * 26 * 27 * 28 * 29 * 30 * 31 * 32 * 33 * 34 * 35 * 36 * 37 * 38 * 39 * 40 * 41 * 42 * 43 * 44 * 45 * 46 * 47 * 48 * 49 * 50 * 51 * 52 * 53 * 54 * 55 * 56 * 57 * 58 * 59 * 60 * 61 * 62 * 63 * 64 * 65 * 66 *
				0000C 07F3 F00C 00004 07 00005 0D7F2C5E3F407 0000C 90FC 000C 00010 000C 00012 000C 00017 1AF0 00018 4101 4232 00019 5U03 F00A 0001C 5UF3 D004 00022 5820 1000 00024 7827 2000 00029 3121 0002A 7021 A27A 00046 3022 0004A 1H55 0004E 4151 0006 0005F 4170 23FC 0005P 5FC1 AF85 00058 3B44 0005A 7901 4286 0005C 4740 4076 0005D 7C05 4284 0005E 5051 A2A2 0005F 1A56 0006A 7A43 AEBA 0006F 1C01 0007C 3D74 0007D 7901 42A6 0007E 4743 A090 0007A 7005 A7RA 0007F 8756 4058 00082 47F1 4C90 00083 7A43 AEBA 000CB8A 3C02
				0000C PSETUP 0000C 21 0000C 22 0000C 23 0000C 24 0000C 25 0000C 26 0000C 27 0000C 28 0000C 29 0000C 30 0000C 31 0000C 32 0000C 33 0000C 34 0000C 35 0000C 36 0000C 37 0000C 38 0000C 39 0000C 40 0000C 41 0000C 42 0000C 43 0000C 44 0000C 45 0000C 46 0000C 47 0000C 48 0000C 49 0000C 50 0000C 51 0000C 52 0000C 53 0000C 54 0000C 55 0000C 56 0000C 57 0000C 58 0000C 59 0000C 60 0000C 61 0000C 62 0000C 63 0000C 64 0000C 65 0000C 66
				PRINT PSETUP DC X1107 DC CL7,2SETUP STM 1412+12113) BALR 10,0 USING 010 LR 1413 LA 13,AKFA ST 14,81H,141 ST 14,41H,131 L 21111 LF 21112 LWER 212 STE 212,LAM CALL FPR,(LAM) LPER 212 SP 5,5 LA 6,4 LA 7,102, LWER 212 STE 212,LAM COMPUTE EXP(-LAM). IF FPR1-LAM>1M, BRANCH TO LO. OTHERWISE STORE FPR1 INTO P11. STORE STARTING POINT. FPR4=X. FPR4=EXP(-LAM), FPR4=FPR1-LAM/X. IF FPR4<LAM, BRANCH TO CONVER. OTHERWISE STORE FPR4 INTO P11. CONTINUE THIS SECTION OF CODE COMPUTES SUCCESSIVE PROBABILITIES

PAGE ?

LCC	OBJECT CODE	AC01 ANDR2	STMT	SOURCE STATEMENT	
00000C 1006		00E9A	67	DER 0,4 A 12,0,F01.	
00000E SAC0 AEB6		0079A	69	CE 0,4LLIM	
000102 7900 4286		0079A	70	THE PROPER STARTING POINT IS STORED.	
000106 4743 A074		0079A	71	A BRANCH IS	
00010A SCC1 A282		0029A	72	MADE TO L2 FOR SUB-	
00000E 4750 A058		0006A	72	SEQUENT CALCULATIONS.	
000CA2 1453		73 CONVER	SR 5,6		
000104 4140 0006		0006A	74	L4 4,4	
00000A 1031			75	SR 3,1	
000104 2821 428A		0029C	75 L3	LE 2,P(3)	
00000E 7620 A6AE		006C0	76	AU 7,2,ED0	
000102 7023 A28A		0029C	78	STE 2,P(3)	
00010E 1A43 426A		0029C	79	L 9,P(1)	
00010A 959 0008		0029C	80	SLL 9,8	
00010F 503 42PA		0029C	81	ST 7,P(3)	
000002 9736 A05A		000AA	82	BKLF 3,6,L3	
00000A 1PCC			P1		
00000A 1A8A			84	SR 12,12	
00010A 1A77			85	SR 11,11	
00010F 1A6A			86	SR 7,7	
00000E 1A71			87	SR 6,6	
00000D 1A43			87	SR 3,3	
000102 5FC1 A2RA		0029C	89	SR 8,4	
00010A 1A6A		nn004	90	SLDL R,4	
00010F 1A7A		00004	91	SLDL R,4	
000102 8D43 3004		00004	92	SLDL R,4	
000006 1ABA		00004	93	SLDL R,4	
00000A 1PC90 3004		00004	94	SLDL R,4	
00010C 1AC9		00004	95	AR 11,11	
00010F 5714 A0RE		00004	96	SLDL R,4	
000102 5063 A69A		0069C	99	AXLE 3,6,L4	
00010F 5073 468E		006AU	100	ST 6,5	
00000A 5081 A692		006A4	101	ST 7,5,4	
00010F 50C1 A696		006AB	102	ST 11,5,R	
000102 850 1004		00004	103	ST 12,5,12	
00010C 4573 0004		00004	104	SLL 11,4	
00010A 8963 0004		00004	105	SLL 7,4	
00010E 1331		00004	106	SLL 6,4	
00010A 1A84		00004	107 L5	SR 3,1	
000112 5893 A2AA		0029C	108	SR 8,4	
000116 8000 3004		00034	109	SLDL 9,(13)	
00011A 1A6A			110	SR 8,4	
00011C 1A7A			111	SR 6,5	
00011F 1A6A			112	SR 7,5	
000120 1A8A			113	SR 11,6	
000122 Pn80 0004		00004	114	SLDL 8,5	
000126 1B78			115	SP 7,5	
00012A 1FPA			116	SR 11,8	
000124 1A8A			117	SR 8,4	
00012C 8000 0004		00004	118	SLDL 8,4	
000130 1RA1			119	SR 11,8	
000132 6734 ACFE		00010	120	BKLF 3,6,L5	
000116 1ACC			121	SR 12,12	

FOBAPRTO 6/10/71

USING THE RECURSIVE RELATION UNTIL
A PROFILE-LIMIT IS ATTAINED.
THE PROPER STARTING POINT IS STORED.
A BRANCH IS
MADE TO L2 FOR SUB-
SEQUENT CALCULATIONS.

R3=1.
CONVERTING TO FIXED POINT.

R9=P(1).
STORE CONVERTED P(1).
CONTINUE

R3=L.
CLEAR RA.

R9=P(1).
R6=0,1.

RR=0,1,D2.

RT=SUM(D1,D2).

RA=0,1,D2,3.

R11=SUM(D1,D2,D3).

PH=0,1,D2,D3,D4.

R12=SUM(D1,D2,D3,D4).

CONTINUE.

STORE SUM(D1) INTO S(1).

STORE SUM(D1,D2) INTO S(1).

STORE SUM(D1,D2,D3) INTO S(1).

STORE SUM(D1,D2,D3,D4) INTO S(1).

R1=SUM(D1,D2,D3,D4).

R7=SUM(D1,D2,D3,D4).

R8=SUM(D1,D2,D3,D4).

R3=1.

CLEAR RA.

R9=P(1).

RA=0,1.

R6=SU4(D1)16-SUM(D1).

RT=SU4(D1,D2)=16-SUM(D1).

R11=SU4(D1,D2,D3)=16-SU4(D1).

CLEAR RA.

R8=0,3.

R11=R1-SUM(D3).

CONTINUE.

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LOC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT		POLARITY 6/30/71
000139	5C60 A69A	006AC		122	ST 6+N	STORE R6 INTO N(0).	
00013C	5070 A69F	006B0		123	ST 7+N+4	STORE R7 INTO N(1).	
000140	5C60 4642	006B4		124	ST 11+N+8	STORE R11 INTO N(2).	
000144	1877			125	SR 7+7		
000146	1800			126	SR 0+0		
000149	4160 0001	00001		127	LA 6+1		
00014C	1831	00294	128 L6	128	SR 3+3		
00014E	5A20 A2B2			129	L2 XST		
000152	1RFA			130 L7	SR 8+N		
000154	5953 420A	0029C		131	L0,0131		
000156	87K3 0004	CNCGS		132	SLNL 6+N		
00015C	5C63 A2PA	2029C		133	ST Q,P(1)	R01=R802,R6=03,R8=04.	
000160	55E2 AEHE	00E43		134	C R+N+0	RESTOR SHIFTEN P(1).	
000164	4760 A16C	0017F		135	BC 8+L9	IF RA=0.	
00016A	1A74			136	AR 7+R	BRANCH TO L8.	
000164	5577 AE92	00F44		137	C T,+F+2000*	OTHERWISE STORE R20 DISCRETE R.V.)	
00016F	4721 A17A	001AC		138	BC 2+L9	CHECK THAT TABLE LENGTH IS NOT EXCEEDED.	
000172	1876			139	SR 7+6		
000174	427C A6B2	006C4		140	STC 3+12	THE NUMBER OF TIMES CORRESPONDING	
00017A	87CA A162	U0174		141	AXLF 12+6+0+4	TO THE DIGIT IN RA.	
00017C	1A76			142	AR 7+R		
00017E	1A75			143 L8	AR 6+N		
000181	8734 4140	00152		144	AXLE 3+N,L7	INCREMENT DISCAFE R.V.	
000184	1AC6			145	AR 0+N	CONTINUE.	
000186	1706			146	CR 0+N	CONTINUE L6 LOOP	
00018A	4740 A13A	0014C		147	AC 4+N,L8	UNTIL FIRST 4 DIGITS OF	
00018C	5A10 A656	006A8		148 L9	L 3+N+12	FACH P(1) HAVE AFFN SHIFTED.	
000190	5B10 AA42	006A4		149	S 3+N+R	THIS SECTION OF CODE	
000194	1C1C			150	CR 7+12	RESISTS N(3)	
000176	47C3 A192	001A6		151	AC 6+OUT	IF TABLE LENGTH IS EXCEEDED.	
000194	1D3C			152	SR 7+12		
00019C	5A37 A6A2	006B4		153	A 3+N+P		
000198	5D10 A642	006A4		154	ST 3+N+R		
0001A4	5AD0 A216	00248	155 OUT	155	DS 0+N		
0001A8	94FC 000C	000NC		156	L 13,ARF+4	ENTRY RANP01	
0001AC	92FF 000C	round		157	LM 14+12+12(13)	USING RANP01,15	
0001D0	01FF			158	MV1 12(13),X,FF,	DC X1,07.	
0001E4	07			159	BC 15+14	BRANCH TO CALLING ROUTINE.	
0001E2	4009C10507D6C9			160	DROP 10		
0001E6	9013 D01A	00018		161	DC CL7+ RANP01		
0001E8	19C3			164	STM 1+3+24(13)		
0001E9	47F3 FG70	0022A	167 BRANCH	165 RANP01	SER 0,0	BRANCH MADE ON FIRST CALL ONLY.	
0001C4	5C29 F4FF	006B8		166	TEMP	GENERATE NEW U.	
0001CA	5035 FSQ2	006BC		168	ST 2+SPAT	STORE FOR NXFT CALL.	
0001CC	1A77			169	SR 2+2		
0001CF	19C1			170	SR 0,0		
0001D1	0021 0006	00004		171	SLNL 2+4	P2=01.	
0001D4	4929 F4E2	0069C		172	C 2+N	IF NICS THEN	
0001D8	4743 FC5A	00214		173	BC 6+EX	BRANCH TO EX	
0001DC	AC20 1004	20304	174	175 N1	SLDL 2+4	P2=0102.	
0001E0	5923 F4E6	006A0		176	C 2+5+4	IF N102>\$111 THEN	

LOC	OBJECT CONCE	ADDR1	ADDR2	STMT	SOURCE STATEMENT	
0001F4	47A0 F036	001FC0		177	NC 11.N2	BRANCH TO N2
0001FB	4823 F4F2	006AC	178	S 2.N	OTHERWISE OBTAIN PROPER TABLE ADDRESS	
0001FC	47F3 F05A	00214	179	E EX.	BRANCH TO EX.	
0001FU	AD23 0004	00004	180 N2	SLDL 2,4	R2=0102D1.	
0001F4	5923 F4EA	026A4	1A1	C 2,S+R	IF DID2D3>S(2) THFN	
0001FR	47B0 F14A	00204	1B2	BC 11.N3	BRANCH TO N3	
0001FC	5B23 F4F6	0C680	1B3	S 2,N+4	OBTAIN PROPER TABLE ADDRESS	
000200	47F3 F05A	00214	1B4	B EX.	BRANCH TO EX.	
0002A4	8C20 0004	00004	1B5 N3	SLDL 2,4	R2=C1D2D3D6.	
0002C8	5520 F4EE	006AB	1B6	C 2,S+12	IF DID2D3D6>S(3) THFN	
0002AC	47B3 F006	001C0	1B7	BC 11.BRANCH	GO TO BRANCH(OBTAIN NFW U).	
000210	5H21 F4FA	006B4	1B8	S 2,N+8	OBTAIN PROPER TABLE ADDRESS.	
000214	42D2 F504	006C4	1B9 FX	IC 0,A(2)	OBTAIN PROPER DISCRETE VALUE.	
000218	5000 F00A	0029C	1B9	ST 0,RES	CONVERT TO FLOATING POINT.	
00021C	9246 FCC6	00290	1B1	MVI QES,X,46*	QES,X,46*	
000222	7A01 F006	00290	1B2	AE 0,RES	NORMALIZED FLT. PT. RESULT IN FPQU.	
000224	9R13 0018	00018	1B3	LW 1,3,24(13)		
00022P	C7FE		1B4	BCR 15,14	RETURN.	
00022A	5810 F502	006BC	1B5 TEMP	L 3,LA	REPLACES THF INSTRUCTION AT RANCH.	
00022E	0203 F006 FC70 CC1C0 0C22A	006BC	1B6	HVC L 3,n(1)	REPLACE INSTRUCTION AT BANCH.	
000234	5E30 1000	0C000	1B7	L 3,n(1,3)	LOAD PRIMER(ONLY D0NF ON FIRST CALL).	
000238	5830 1600	0C000	1B8	L 3,LA		
00023C	5030 F502	006BC	1B9	ST RANCH+4	GO TO INSTRUCTION FOLLOWING RANCH.	
00024C	47F3 F00A	001C4	2001 AREA	DS 1RF		
00024E			2002 LAM	DS 1F		
00024F			2003 RES	DS 1F		
000254			204 XST	DS 1F		
00029A	3CFARR82		205 LLIM	DC F=0,0,0,015,		
00029C			206 P	DS 256F		
00029C			207 S	DS 4F		
0002AC			2CA N	DS 3F		
0006A8	48C27105		209 BPAT	DC XL4=48C27395*		
0006AC			210 LA	DS 1F	X'4000C000'	
0007C0	4000CU00		211 ZERO	DC 2000XL1		
0006C4			212 A	DS END		
000798	00000001		213		=F'1'	
00079C	4110C000		214		=E'1'	
00079C	0001nnnn		215		=F'0'	
0007A4	00000700		216		=F'20C0'	
0007A4			217			

APPENDIX G

Assembler Listing of GEN5

PAGE 1

FOURPART 6/30/71

LOC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
000000					1 GFM6 START A
					2 SOURCE---L.F. CANNON, UGA, DEPT. OF STAT., JUN 1970
					3 * PURPOSE---TO SUPPLY THE USER WITH A FAST PROCEDURE FOR GENERATING NEGATIVE BINOMIAL DISTRIBUTED RANDOM NUMBERS.
					4 * USAGE---1. CALL NHSETUP(IRI). P-PROBABILITY OF A SUCCESS, IR-NUMBER OF SUCCESSES. NSETU SETS UP THF TABLE FROM WHICH RANGE OF
					5 * AL COUNTERS A RANDOM NUMBER. NSETU NEED BE CALLED ONLY ONCE FOR A GIVEN P AND IR.
					6 * 2. X=RANFD(IRD). IRD MUST BE AN MCP INTEGER. IT PRIMES THE COUNTERING SCHEM. X REPRESENTS THE NUMBER OF FAILURES.
					7 * BEFORE THE IR, TM SUCCESS.
					8 * EXAMPLE--CALL NHSETUP(.5,10)
					9 DO 1 I=1,10
					10 * RANFD(I,.695)
					11 * METHOD---THE METHOD IS THAT OF MARSAGLIA AS DESCRIBED IN "COM. ACM"
					12 * 6 JULY 1963 137.
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					53 *
					54 *
					55 *

17 * METHOD---THE METHOD IS THAT OF MARSAGLIA AS DESCRIBED IN "COM. ACM"

18 * 6 JULY 1963 137.

19 *

20 ENTRY NSFTU

21 NSFTU 0 12101E1

22 NC XL1107

23 OC CL7* "5SFTRU"

24 SIM 14-17-1261131

25 MUL IN-O

26 USING *10

27 LR 14-13

28 LA 11-A7A

29 ST 13-410141

30 ST 14-410131

31 L 4-6111

32 L 4-6111

33 L 2-2111

34 SER 2-4

35 L 6-4111

36 L 6-6111

37 ST 6-6111

38 MUL 6-6111

39 DIV 6-6111

40 SCR 6-6

41 AE 6-SUC

42 BC R SKIP

43 BC L 7-0-F0-1*

44 AR 6-7

45 NFR 0-6

46 MMW 6-7-0-2

47 SR 5-5

48 LA 6-4

49 LA 7-1020

50 LA 12-0-F1*

51 SER 6-4

52 CE 0-LIM

53 BC 0-LO

54 SITE 0-PLS

55 ST 5-XST

IF P=0 CALL MM

BRANCH TO LO.

OTHERWISE STORE PEE TO P(0).

STORE RS=0 INTO XST.

LFC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
000C79 1A46	00007A 7A43 AE9E	QUEAN	96	AR 5.6	F004=X01.
000C7F 1C06	000080 1D94	QUEAN	57 L2	AE 4.0E+10	FPR4=P1(X+1).
000C82 1C02	000084 A2A2	QUEAN	54	DER 0.6	FPR4=P1(X+1)/X+1.
000C84 7900 A2A2	0002A4 200RF	QUEAN	59	MFR 0.4	FPR4=P1(X+1)-P1(X+1)/X+1.
000C86 6740 A2AC	0002B8 62	CONVER	60	MFR 0.2	IF P(X+1) < LIM.
000C8C 7035 A2A6	0002B8 63	CONVER	61	CE 0.11M	BRANCH TO CONVER.
000C90 7A53 AE9E	0002B8 64	STE 0.051	62	BC 4.CONVER	OTHERWISE STORE FPR4(P1(X)) INTO P(55).
000C92 6756 A068	0002B8 65	STE 0.051	63	AF 4.0E+10	INCREMENT X+1.
000C93 47F3 A0AC	0002B8 66	RXLE 5.0.L2	64	BC 4.CONVER	CONTINUE.
000C95 7A52 AE9E	0002B9 67 L7	STE 0.051	65	AF 4.0E+10	THIS SECTION OF CODE
000CA3 1C06	0002B9 66	DER 0.6	66	DER 0.6	CALCULATES SUCCESSIVE PROBABILITIES, USING THE RECURSIVE
000CA7 1C04	0002B9 69	MFR 0.2	67	A 12.0E+10	RELATION. UNTIL A PROBABILITY > LIM IS OBTAINED.
000CB7 2AC0 2FA6	0002B9 71	AE 6.0E+10	68	CF 0.11M	THE CORRESPONDING X-VALUE IS STORED INTO XST
000CB9 7A61 AF9F	00UFAN 72	CE 0.11M	69	BC 4.1D	AND A BRANCH
000CDAE 7501 A2A2	00UFAN 73	ST 12.45T	70	ST 12.45T	IS MADE TO L2
000CDE2 47E3 A0AA	00UFAN 74	AC 15.12	71	AC 15.12	FOR FURTHER CALCULATIONS.
000CD9 50C0 A29F	00UFAN 75	SR 5.6	72	SR 5.6	3.3%.
000CD9 47F0 A768	00UFAN 76	LA 4.4	73	LA 4.4	CONVERTING TO FIXED POINT.
000CD9 1956 1956	00UFAN 77 CONVER	SR 3.1	74	SR 3.1	3.3%.
000CD9 4160 2004	00UFAN 78	LE 2.0.P13I	75	LE 2.0.P13I	CONVERTING TO FIXED POINT.
000CD9 1H33	00UFAN 79 L3	AU 2.0.EQU	76	AU 2.0.EQU	3.3%.
000CD9 7A21 A2A6	00UFAN 80 R1	ST 2.0.P13I	77	ST 2.0.P13I	3.3%.
000CD9 7571 46EA	00UFAN 81	SR 11.11	78	SR 11.11	3.3%.
000CD9 7721 42A6	00UFAN 82	SR 7.7	79	SR 7.7	3.3%.
000CD9 5491 4244	00UFAN 83	SL0L 6.4	80	SL0L 6.4	3.3%.
000CD9 9C00 1H0R	00UFAN 84	SL0M 8.4	81	SL0M 8.4	3.3%.
000CD9 5091 A2A6	00UFAN 85	ST 7.6	82	ST 7.6	3.3%.
000CD9 8724 4094	00UFAN 86	BXF 3.6E13	83	BXF 3.6E13	CONTINU
000CD9 1HCC 1HCC	00UFAN 87	SR 12.12	84	SR 12.12	CLEAR AB.
000CD9 1P9A 1P9A	00UFAN 88	SR 11.11	85	SR 11.11	P9=0(11).
000CD9 1P77 1P77	00UFAN 89	SR 7.7	86	SR 7.7	RE01.
000CD9 1A66 1A66	00UFAN 90	SR 6.4	87	SR 6.4	R6=0(D010).
000CD9 1A31 1A31	00UFAN 91	SR 6.4	88	SR 6.4	R8=0(D02).
000CD9 1H04 1H04	00UFAN 92 L4	SR 3.3	89	SR 3.3	R7=SUM(D02D1).
000CD9 5P91 42A6	00UFAN 93	AR 11.0.8	90	AR 11.0.8	R11=SUM(D01D2D1).
000CD9 2D00 3C04	00UFAN 94	SL0L 6.4	91	SL0L 6.4	R6=0(D010).
000CD9 1A69 1A69	00UFAN 95	AR 6.4	92	AR 6.4	R12=SUM(D01D2D4).
000CD9 2D03 3004	00UFAN 96	SL0M 8.4	93	SL0M 8.4	R8=0(D02).
000CD9 1A70 1A70	00UFAN 97	AR 7.6	94	AR 7.6	R7=SUM(D01D2).
000CD9 8D83 0004	00UFAN 98	SL0L 6.4	95	SL0L 6.4	STORE SUM(D1) INTO S01.
000CD9 1A84 1A84	00UFAN 99	AR 7.6	96	AR 7.6	STORE SUM(D1P2P1) INTO S11.
000CD9 5C71 46AA	00UFAN 100	ST 7.6	97	ST 7.6	STORE SUM(D1D2D3) INTO S12.
000CD9 5CD7 464C	00UFAN 101	ST 11.0.8	98	ST 11.0.8	STORE SUM(D1D2D3D4) INTO S13.
000CD9 5C03 4612	00UFAN 102	ST 12.5.12	99	ST 12.5.12	R11=SUM(D02D3D4).
000CD9 89B3 1004	00UFAN 103	SLL 11.4	100	SLL 11.4	R7=SUM(D02D3D6).
000CD9 8912 8910 0004	00UFAN 104	SLL 7.6	101	SLL 7.6	R6=0(D010).
000CD9 89B0 0006	00UFAN 105	SLL 6.4	102	SLL 6.4	R3=1.
000CD9 1A33	00UFAN 106	SR 3.3	103	SR 3.3	

LNC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT	FORTRAN	PAGE
00012C	1A69	002B8	111 LS	SA	0,9 CLFAK RA. R9-P(1).	6/30/71	3
00012E	4F91 42A6	00004	112	SLNL	R4-N.		
0C0112	8C97 0004		113	SLNL	R4-N.		
0C0113	1A69		114	SR	R6-SUM(D1)16-SUM(D1).		
0C0113	1A70		115	SR	R7-SUM(D1)16-SUM(D1).		
0C0114	1A69		116	SR	R8-SUM(D1)16-SUM(D1).		
0C0115	1A69		117	SR	CLEAR R8. R8=0.		
0C0115	8C97 0004	00004	118	SLNL	R7=0.		
0C0112	1A70		119	SR	R7-R1-SU=1021.		
0C0114	1A69		120	SR	R1=BL1-SUM(D2).		
0C0116	1A69		121	SR	CLEAR R8. PA=0.		
0C0116	4A 3C11 0004		122	SLNL	R1=BL1-SUM(D3).		
0C011C	1A69		123	SR	CONTINUE.		
0C011C	8756 A11A		124	PALE	374-115		
0C0112	1ECC		125	SR	12-12		
0C0114	5D61 4686	000C8	126	ST	6*N		
0C0119	5C70 4684	000CC	127	ST	STORE M1 INTO N11.		
0C011C	5C71 4A8E	000D0	128	ST	STORE M1 INTO N12.		
0C011C	1A77		129	SR	31-N6		
0C0112	1A69		130	SR	7-7		
0C0114	4160 0001		131	LA	0,1		
0C0114	4160 1A69		132	SR	3-1		
0C0114	4160 129F	00001	133	SR	2-15T		
0C0114	1A69	002B0	134	SR	6,1		
0C0114	4A61 42A6		135	SLNL	R1=BL1.		
0C0114	6C00 0004	000C4	136	SLNL	R1=0, PA=02,R9=03,RE=74.		
0C0117	5C91 121A	000CA	137	ST	RFST(M SHIFTEN BL1).		
0C0112	5C43 4E84	000AC	138	C	IF BRIC.		
0C0110	4781 4188	0017A	139	EC	APAW=M70 UP.		
0C011A	1A76	140	AR	7-A	OTHERWISE STORE M2 DISCRETE & V.		
0C0116	4971 AFAE	00FCU	141	C	CHECK THAT TABLE LENGTH IS NOT EXCEEDED.		
0C011A	4121 4196	001A8	142	RC	2-L9		
0C011F	1A76	143	SR	7-6	TIME NUMBER OF TIMES CORRESPONDING		
0C0120	422C A6CE	000C6	144	STC	TO THE NIGHT IN RA.		
0C0119	81C6 A17E	00190	145	MUL	12-6,-e-4		
0C0114	1A76	146	AR	7-4			
0C0114	8754 419C	001E6	147	AC	INCREMENT DISCRETE R,V.		
0C0111	1A69	148	AR	3-4,L7	CONTINUE.		
0C0112	1A69	149	AR	0-4	CONTINUE 14 LOMP		
0C0114	4781 417A	001E9	150	CR	UNTIL FIRST 4 DIGITS OF		
0C0119	5A90 4A62	006C4	152	LC	EACH P111 HAVE BEEN SHIFTED.		
0C011C	5A91 A68E	006D0	153	S	THIS SECTION OF CODE		
0C0117	193C	001C9	154	CR	RESETS M:31		
0C0112	4781 A1AF	001C0	155	BC	IF TABLE LENGTH IS EXCEEDED.		
0C0116	1E9C	006ED	157	SR	3-12		
0C0114	5A90 A68E	006D0	158	A	3-4,e-4		
0C011C	5013 A68E	006D0	159	ST	3-N+8		
0C0110		001E9	160	DS	ON		
0C011A	5E99 A252	001E9	161	LM	13-4RE,A44		
0C0114	985C D90C	0030C	162	NVI	14-12,12(13)		
0C011C	017C D90C	163	NCR	15-14	DROP		
0C011C	017C	164	ENTRY RANGE!	10			
		165					

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LNC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
00106				166	USING RANEBS,15
00116	07			167	OC RLL,O7,
0011C	4020:1055C2C9			168	OC CLT, RANEBS,
0011S	9F13 D018	00018		169	RANEBS
0011A	3H00			170	ST 1,7,24(13)
0011C	47F3 FC70	00746	171	BRANCH	BRANCH MADE ON FIRST CALL ONLY.
0011G	SC20 F4F6	00604	172	B	TEMP
0011F	5C93 <507	00608	173	ST	2,RPAT
0011R	1R22		174	SR	3,LA
0011A	1803		175	SR	2,2
0011C	6,1E2 6,1D3 3D04	00004	176	SL0L	A,0
0011G	5922 F4F2	0058A	177	C	R2=01.
0011A	6743 F0A4	00210	178	BC	IF R1CS THEN
0011F	AC27 6004	00014	179	BRANCH TO EX	
0011C	5922 6004	00604	180	NI	R2=D1D2.
0011C	5921 F4F6	006AC	180	NI	IF D1D2>5(11) THEN
0011G	47B3 F,016	0070C	181	RC	BRANCH TO N7
0011A	5,22 F4F2	005CA	182	S	OTHERWISE CRAIN PROPER TABLE ADDRESS
0011C	47F3 F0A4	00230	183	SL0L	BRANCH TO EX.
0011G	AC27 3D04	0061A	184	NI	R2=D1D2D3.
0011A	5922 F4F4	00604	184	NI	IF D1D2N1>5(12) THEN
0011F	5921 F4F4	006C0	185	C	BRANCH TO N3
0011C	47B3 F,01A	00220	186	RC	MINTAIN PROPER TABLE ADDRESS
0011G	5n22 F4F6	006CC	187	S	BRANCH TO EX.
0011A	5,22 F4F2	00230	188	A	R2=D1D2D3D4.
0011F	5922 Jn3 3D04	00604	189	NI	IF D1D23D4>5(13) THEN
0011C	47F3 F4F6	006AC	190	C	GO TO BRANCH MINTAIN NEW U3.
0011G	5921 F4F6	0010C	191	RC	MINTAIN PROPER TABLE ADDRESS.
0011A	5,22 F4F4	00600	192	S	CONTAIN PROPER DISCRETE VALUE.
0011F	5921 410 F50A	006E0	193	FX	CONVERT TO FLOATING POINT.
0011C	5n3 F0D2	0070A	194	ST	0,RES
0011G	9264 F0D2	002AB	195	HVI	RESX=46.
0011A	7A03 Fm2	002AB	196	AF	0,RES
0011F	CALL m1A	00718	197	LM	NORMALIZED FLT. PT. RESULT IN FPAC.
0011C	37FF		198	REC	RETURN.
0011G	5811 F502	00608	199	TEMP	REPLACES THE INSTRUCTION AT BRANCH.
0011A	D201 F0D6 F070 C910C 00246	2C9	200	HVC	REPLACE INSTRUCTION AT BRANCH.
0011F	SP13 1m00	00600	201	L	LOAD PRIMONLY DONE ON FIRST CALL).
0011C	SP13 1m00	00600	202	L	LOAD PRIMONLY DONE ON FIRST CALL).
0011G	5D13 5402	00608	203	ST	3,0(1,3)
0011A	47B3 F0A4	0010E	204	R	GO TO INSTRUCTION FOLLOWING BRANCH.
0011F			205	AREA	3,LA
0011C			206	RES	4F
0011G			207	SUC	IF
0011A			208	ST	IF
0011F	3CFRA992	209	LLIN	DC	E,0,000015.
0011C			210	P	256F
0011G			211	S	4F
0011A			212	N	IF
0011F	4RC27393	213	BPAT	DC	X,4,4RC27393.
0011C			214	LA	IF
0011G	4C00GJ00	215	EQD	DC	X,400C0000.
0011A			216	A	END
0011F			217		=E,1.
0011C			218		=SF,-1.
0011G			219		=SF,1.
0011A			220		

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SOURCE STATEMENT

LOC OBJECT CODE ADDR1 ADDR2 STMT

221
222"F"0"
"F"20U0"CONFC OCO70000
OONEC O000007D0

APPENDIX H

Assembler Listing Of GEN7

NOT REPRODUCIBLE

PAGE 2
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LTC	OBJECT CODE	ACTR1	ADDR2	STMT	SOURCE STATEMENT
000074	8080 0004	00004		96	SLDL 0,4
000074	1AC9			97	AP 12,0
00007C	813A ADAC	0005E	SR	98	BKLE 2,4,14
000080	5C43 AKA	0016C	ST	99	CONTINUF.
000084	1070 AF0F	00120	ST	100	STORE SUM(011) INTO S111.
00008A	5C43 A012	00124	ST	101	STORE SUM(012) INTO S112.
00008C	40C9 4616	00678	ST	11,5+8	STORE SUM(010) INTO S110.
00008E	65A0 0004	00004	SLL	12,5+12	STORE SUM(01100000) INTO S110.
000094	8979 0004	00004	SLL	11,4	R11-SUM(01010116).
000098	8963 0004	00004	SLL	11,4	R7-SUM(01010116).
00009C	1A93	00004	SLL	11,4	RA-SUM(01111116).
00009E	1A93	64	SR	3,1	R11-SUM(01111116).
0000A1	5893 427A	0021C	SR	3,1	CLEAR RA.
0000AA	829A 004	00004	SLDL	0,4	R9-P11.
0000AA	1A63	68	SR	4,1	R4-SUM(0101016-SUM(011).
0000AA	1A79	69	SR	4,1	R4-SUM(0101016-SUM(011).
0000AC	1A84	70	SR	4,1	R7-SUM(0101016-SUM(011).
0000AE	1A84	71	SR	7,8	R11-SUM(0101016-SUM(011).
0000B0	AD02 0004	00004	SLDL	0,4	CLEAR RA.
0000B4	1A79	72	SR	11,8	R4-P02.
0000B6	1B04 1B04	00004	SLDL	0,4	R7-R11-SUM(02).
0000B8	1B04	73	SR	6,4	R11-R11-SUM(02).
0000B9	1B04 1B04	00004	SLDL	0,4	CLEAR RA.
0000C1	1A91	74	SR	11,8	R6-P03.
0000C1	8734 A09C	00004	SLDL	0,4	R11-R11-SUM(03).
0000C4	1AC9	75	SR	11,8	CONTINUF.
0000C6	5C63 A61A	0062C	P2	7,4	STORE RA INTO M110.
0000CA	5D0A 5D1A	00610	A3	7,4	STORE R7 INTO M111.
0000CF	5A73 A61E	00634	A4	7,4	STORE ALL 140N M121.
000102	1A77	00004	SLDL	0,4	R11-SUM(03).
0000D4	1A60	76	SR	11,8	LOAD R11 WITH ADDRESS OF X VECTOR.
0000D4	5E0A 5E06 1C04	00010	AN	12,12	R11-P11.
0000D8	5E0A 5E06 1C04	00004	SLDL	0,4	CLEAR RA.
0000F2	1P83 A20A	0021C	SR	9,1	R9-P11.
0000FA	5D0A 5D04	00004	SLDL	0,4	R4=01, R2=02, R8=03, RA=D4.
0000FA	5D03 A20A	0021C	ST	9,1	RESTORE SHIFTED P111.
0000F5	5C58 4E0A	0061C	RC	8,0,0	IF RA=0.
0000F2	4780 AF0E	00110	RC	8,0,0	BRANCH TO 10.
000104	1A79	00004	SLDL	0,4	OTHERWISE STORE R20 DISCRETE #, Y, I.
0000F4	5C70 AF0E	00F20	97	7,-,0 20000	CHECK THAT TABLE LENGTH IS NOT EXCEEDED.
000104	4720 1104	0011C	RC	2,0,0	LOAD R2 WITH DISCRETE VALUE.
000100	5P21 4060	00000	99	2,0,1,111	THE NUMBER OF TIMES CORRESPONDING
000104	1A79	100	SR	7,6	TO THE DIGIT IN RA.
000126	422C 4632	00644	101	7,A121	CONTINUF. TO LMP
00012A	52C6 AF04	00106	102	12,6,0-4	UNTIL FIRST 4 DIGITS OF
000105	1A74	103	SR	0,4	FACH P111 HAVE BEEN SHIFTED.
000110	5714 40CE	000F0	104	1,4,17	THIS SECTION OF CODE
000114	1A63	105	SR	5	RECS M111
000114	1A54	106	CR	4,16	
000114	4741 AFCC	000DE	107	5	
00011C	5A30 A616	00028	108	3,1012	
000123	5A30 AF22	00A3A	109	3,M6	
000124	1A3C	110	CR	1,12	

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LOC	OBJECT CODE	ADDR1	ADDR2	STAT	SOURCE STATEMENT
000126	47B0 4122	00134	111	AC	4.0UT
00012A	181C	112		SR	1.12
00012C	1A10 4672	00634	113	A	3.H+A
000130	5731 4622	00634	114	ST	3.Y+8
000134	00114 540N A1C2	00134	115 OUT	DS	0W
000138	00118 9EFC 00UC	00134	116	L	13.ARE+4
00013C	07FF	00134	117	LM	15.12.12(13)
000140	C7	00014	118	BRK	15.14
000144	9003 0014	00014	119	DRUP	10
00014A	1801	00146	120	ENTRY RANDIS	
00015C	47FD F070	00146	121	USING RANDIS15	
00015C	1C27 F4F2	00146	122	DC	X11'07'
000154	5032 F4F6	00146	123	DC	CLT ¹ RANDIS ⁰
00015A	1822	00146	124 RANDIS	SIW	0.1.2(13)
00015C	9C23 0004	00146	125 BRANCH	SFR	0.0
000154	5420 4606	00636	126	N	2.RPAT
00015A	1801	00636	127	ST	3.LA
00015C	9C23 0004	00636	128	SR	2.P
000154	5420 4606	00636	129	SR	2.P
00015A	1822	00636	130	SR	0.0
00015C	9C23 0004	00636	131	SLOL	2.4
000154	5420 4606	00636	132	AC	2.S
00015A	1822 0054	00636	133	AC	4.FX
00015C	5420 4606	00636	134	SLOL	2.4
000154	5420 4606	00636	135	AC	2.S+4
000170	47B0 F016	00636	136	AC	11.N2
000174	5A23 54F6	0562C	137	S	7.N4
000178	47F0 F05A	00146	138	B	EX
00017C	9022 0006	000146	139 N2	SPL	2.4
00018A	1920 F40F	00626	140	C	2.S+8.
000184	47A0 F04A	00190	141	AC	11.N4
00018A	5A23 54FA	00619	142	S	7.N4
00018C	47F3 F05A	00146	143	R	EX
0001A0	RU21 0006	000146	144 N3	SLOL	2.4
000194	5420 4606	00628	145	C	2.S+12
00019A	47A1 F016	00146	146	BC	11.NRANCH
00019C	5420 46EF	00634	147	S	2.N+8
0001A4	47A0 F04A	00646	148 EX	IC	0A1?1
0001A4	5D03. F002	00216	149	ST	0A1?1
0001A8	9246 F002	00216	150	RES	X+46.
0001AC	7A01 F002	00216	151	AE	0.RES
0001A0	9803 0114	00C14	152	LR	0.3.2(13)
000144	17FF	153		ACR	15.14
000186	5A30 F4F6	0063C	154 TEMP	L	1.LA
0001A4	0203 F006 FC70 C014C	00146	155	MVC	3.0(1.1)
0001C4	5A33 3UN0	00000	156	L	3.0(1.1)
0001CA	5A33 3UN0	00000	157	ST	3.LA
0001CC	47F0 F004	0063C	158	S	3.LA
0001C4	5A33 3UN0	00150	159	BRANCH	1SF
000140	1900	160 AREA	DS	1SF	
000218		161 RES	DS	1F	
00021C		162 P	DS	25F	
00061C		163 S	DS	4F	
00062C		164 N	DS	3F	
00063E	48C27195	165 BPAT	DC	XL4.4AC27395!	

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LCC	OBJECT CODE	ADDR1	ADDR2	STMT	SOURCE STATEMENT
00063C	0CU640 40000000			166	LA DS 1F
300644				167	ZERO DC X'40000000'
000E18	0.0000001			168	A DS 2000XL1
000F1C	0.00000000			169	END
000F20	000000700			170	
				171	=F'1'
				172	=F'0'
					=F'2000'

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